

A boundary approach for set inversion

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Virtual Seminars on Interval Methods in Control

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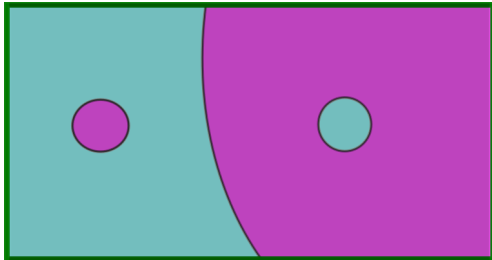
2020, December 18



Motivation

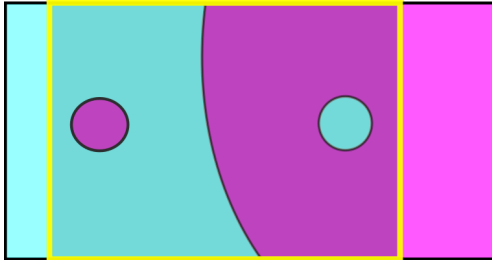
Motivation

Forward-backward sequence
Directed contractors
Boundary approach
Test-cases



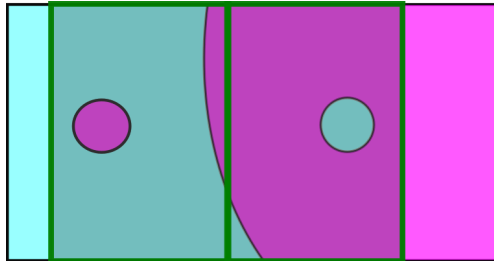
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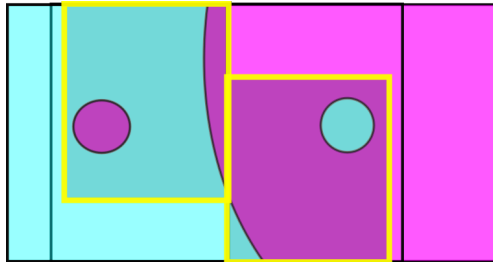
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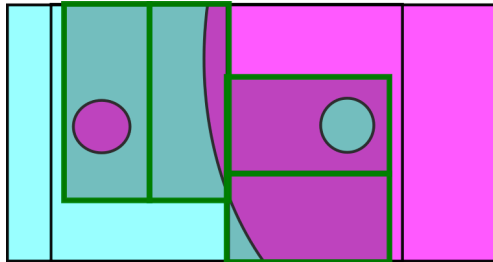
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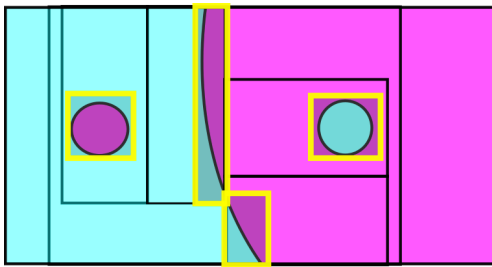
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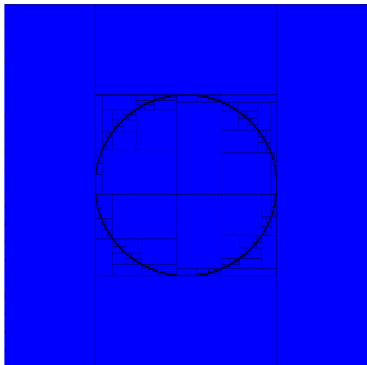
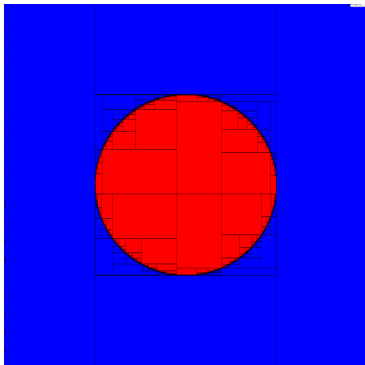
Motivation

Forward-backward sequence

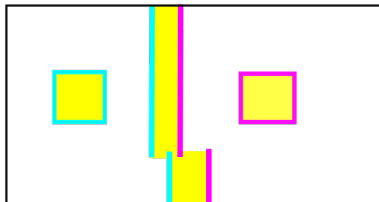
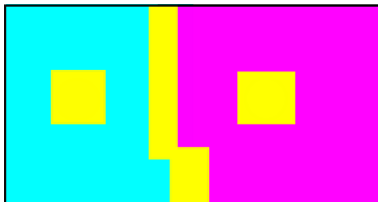
Directed contractors

Boundary approach

Test-cases



- A lot of computation are performed twice with classical approaches.
- With the boundary approach : keep the color.



What we usually compute. What we propose to compute[3]

Idea: A box is represented with colored faces

$$[\mathbf{x}] = [1, 2] \times [1, 4].$$

Forward-backward sequence

$$\mathbf{f}(\mathbf{x}) \in \mathbb{Y}, \mathbf{x} \in \mathbb{X}(0)$$

$$\mathbf{f} = \mathbf{f}_n \circ \cdots \circ \mathbf{f}_2 \circ \mathbf{f}_1$$

$$\mathbb{X} = \mathbb{X}(0) \cap \mathbf{f}^{-1}(\mathbb{Y})$$

```

Input:  $\mathbb{X}(0)$ 
1  For  $k = 1$  to  $n$ 
2       $\vec{\mathbb{X}}(k) = \mathbf{f}_k(\vec{\mathbb{X}}(k-1))$ 
3   $\overleftarrow{\mathbb{X}}(n) = \mathbb{Y} \cap \vec{\mathbb{X}}(n)$ 
4  For  $k = n$  to 1
5       $\overleftarrow{\mathbb{X}}(k-1) = \vec{\mathbb{X}}(k-1) \cap \mathbf{f}_k^{-1}(\overleftarrow{\mathbb{X}}(k))$ 
Return  $\overleftarrow{\mathbb{X}}(0)$ 
    
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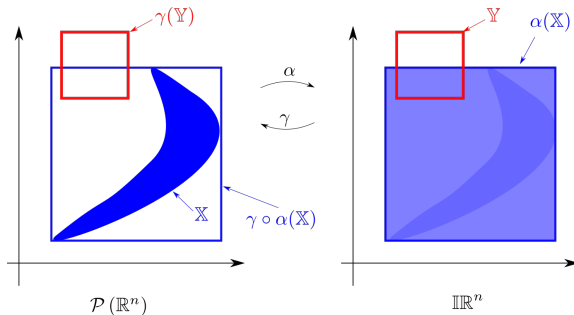
$\vec{\mathbb{X}}(k)$: set of states at time k consistent with the past
 $\overleftarrow{\mathbb{X}}(k)$: set of states consistent with both past and future

Let $(\mathbb{A}, \leq)$ and $(\mathbb{B}, \leq)$ be two partially ordered sets [1].
A Galois connection consists of two monotonic functions:
 $\alpha : \mathbb{A} \rightarrow \mathbb{B}$ and $\gamma : \mathbb{B} \rightarrow \mathbb{A}$, such that

$$\alpha(x) \leq y \Leftrightarrow x \leq \gamma(y).$$

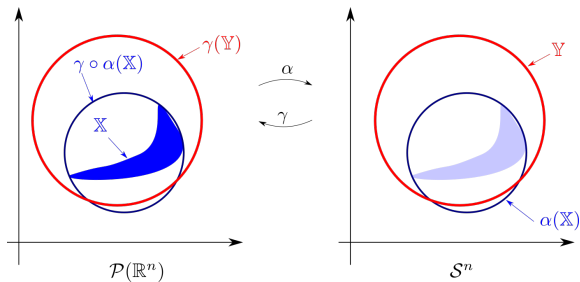
With boxes, we have

$$\alpha(\mathbb{X}) \subset \mathbb{Y} \Leftrightarrow \mathbb{X} \subset \gamma(\mathbb{Y})$$

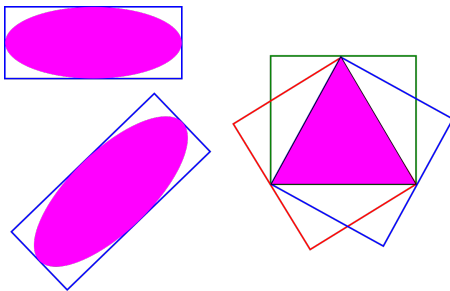


With sphere, we have

$$\alpha(\mathbb{X}) \subset Y \Rightarrow \mathbb{X} \subset \gamma(Y)$$



Oriented boxes do not yield a Galois connection



This excludes Lohner-type algorithms

Directed contractors

A *directed contractor* $\mathcal{C}$ for the constraint $\mathbf{y} = \mathbf{f}(\mathbf{x})$ is a pair of two operators

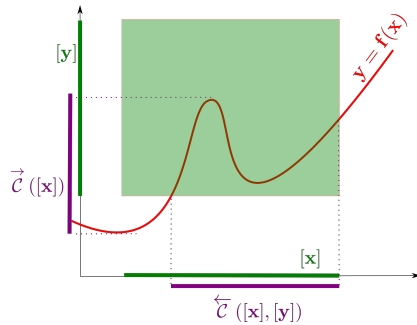
$$\mathcal{C} : ([\mathbf{x}], [\mathbf{y}]) \rightarrow \left(\overset{\rightarrow}{\mathcal{C}}([\mathbf{x}]), \overset{\leftarrow}{\mathcal{C}}([\mathbf{x}], [\mathbf{y}]) \right)$$

with

$$\begin{aligned} \mathbf{f}([\mathbf{x}]) &\subset \overset{\rightarrow}{\mathcal{C}}([\mathbf{x}]) \\ \mathbf{f}^{-1}([\mathbf{y}] \cap [\mathbf{x}]) &\subset \overset{\leftarrow}{\mathcal{C}}([\mathbf{x}], [\mathbf{y}]) \subset [\mathbf{x}] \end{aligned}$$

Moreover

$$\left\{ \begin{array}{l} [\mathbf{a}] \subset [\mathbf{x}] \\ [\mathbf{b}] \subset [\mathbf{y}] \end{array} \right\} \Rightarrow \left\{ \begin{array}{l} \overset{\rightarrow}{\mathcal{C}}([\mathbf{a}]) \subset \overset{\rightarrow}{\mathcal{C}}([\mathbf{x}]) \\ \overset{\leftarrow}{\mathcal{C}}([\mathbf{a}], [\mathbf{b}]) \subset \overset{\leftarrow}{\mathcal{C}}([\mathbf{x}], [\mathbf{y}]) \end{array} \right.$$

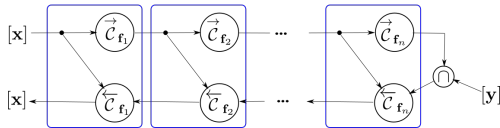
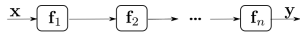


Contractor chain [2]

$$\mathbf{f}(\mathbf{x}) \in [\mathbf{y}], \mathbf{x} \in \mathbb{X}(0)$$

$$\mathbf{f} = \mathbf{f}_n \circ \dots \circ \mathbf{f}_2 \circ \mathbf{f}_1$$

$$\mathbb{X} = \mathbb{X}(0) \cap \mathbf{f}^{-1}([\mathbf{y}])$$



Composition

$$\mathcal{C}_2 \circ \mathcal{C}_1 \left(\begin{bmatrix} \mathbf{x} \\ \mathbf{y} \end{bmatrix} \right) = \left(\vec{\mathcal{C}}_2 \circ \vec{\mathcal{C}}_1([\mathbf{x}]), \overleftarrow{\mathcal{C}}_1 \left(\overleftarrow{\mathcal{C}}_2 \left(\begin{bmatrix} \mathbf{x} \\ \vec{\mathcal{C}}_1([\mathbf{x}]) \\ \vec{\mathcal{C}}_2 \circ \vec{\mathcal{C}}_1([\mathbf{x}]) \cap [\mathbf{y}] \end{bmatrix} \right) \right) \right)$$

The minimal directed contractor for the constraint
 $y = f(\mathbf{x}) = x_1 + x_2$ is:

$$\begin{aligned}\overrightarrow{\mathcal{C}}([\mathbf{x}]) &= [x_1] + [x_2] \\ \overleftarrow{\mathcal{C}}([\mathbf{x}], [y]) &= \begin{pmatrix} [x_1] \cap ([y] - [x_2]) \\ [x_2] \cap ([y] - [x_1]) \end{pmatrix}\end{aligned}$$

A function $\mathbf{f}$ for which a minimal directed contractor is available is said to be *contractible*.

Since the minimal contractor for

$$\mathbf{y} = \mathbf{A} \cdot \mathbf{x}$$

is

$$\mathcal{C} \left(\begin{array}{c} [\mathbf{x}] \\ [\mathbf{y}] \end{array} \right) = \left(\begin{array}{c} \vec{\mathcal{C}}([\mathbf{x}]) \\ \overleftarrow{\mathcal{C}}([\mathbf{x}], [\mathbf{y}]) \end{array} \right) = \left(\begin{array}{c} \mathbf{A} \cdot [\mathbf{x}] \\ \mathbf{A}^{-1} \cdot [\mathbf{y}] \cap [\mathbf{x}] \end{array} \right)$$

the function $\mathbf{f}(\mathbf{x}) = \mathbf{A} \cdot \mathbf{x}$ is contractible.

Contractible decomposition

A *contractible decomposition* of a function $\mathbf{f}$ has the form

$$\mathbf{f} = \mathbf{f}_n \circ \cdots \circ \mathbf{f}_2 \circ \mathbf{f}_1 = \mathbf{f}_{1:n}$$

where each $\mathbf{f}_i$ is contractible.

Counterexample. The Fresnel integral

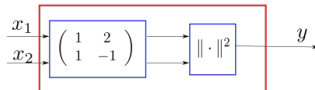
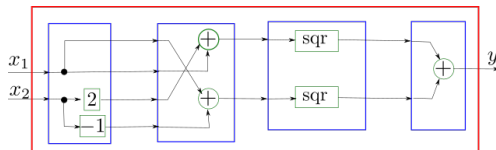
$$f(x) = \int_0^x \sin \tau^2 \cdot d\tau$$

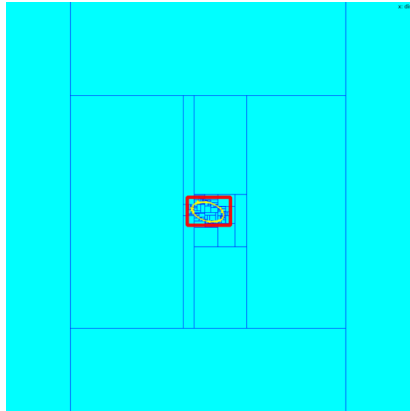
has no contractible decomposition.

The function

$$f(x_1, x_2) = (x_1 + 2x_2)^2 + (x_1 - x_2)^2$$

has scalar and vector decompositions:



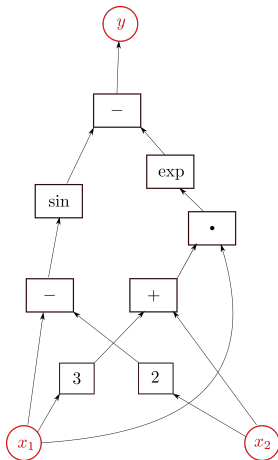


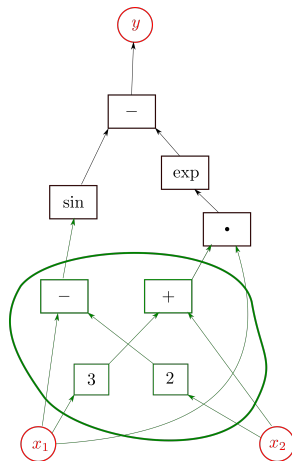
$$(x_1 + 2x_2)^2 + (x_1 - x_2)^2 = 1$$

Forward and cut algorithm

We want the contractible decomposition of

$$y = \sin(x_1 - 2x_2) - \exp(x_1 \cdot (3x_1 + x_2))$$

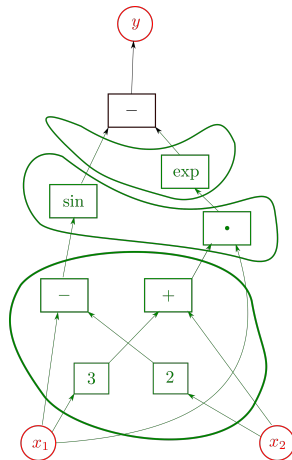




$$\begin{pmatrix} 1 & -2 \\ 1 & 0 \\ 3 & 1 \end{pmatrix}$$



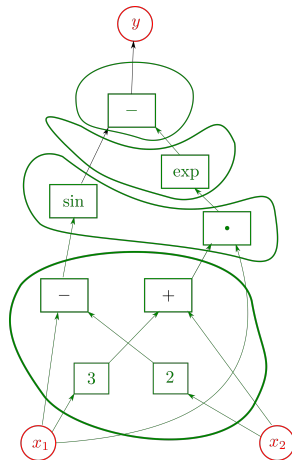
$$\begin{pmatrix} 1 & -2 \\ 1 & 0 \\ 3 & 1 \end{pmatrix}$$



$$\begin{pmatrix} 1 \\ \exp \end{pmatrix}$$

$$\begin{pmatrix} \sin \\ \cdot \end{pmatrix}$$

$$\begin{pmatrix} 1 & -2 \\ 1 & 0 \\ 3 & 1 \end{pmatrix}$$



$$\begin{pmatrix} 1 & -1 \end{pmatrix}$$

$$\begin{pmatrix} 1 \\ \exp \end{pmatrix}$$

$$\begin{pmatrix} \sin \\ \cdot \end{pmatrix}$$

$$\begin{pmatrix} 1 & -2 \\ 1 & 0 \\ 3 & 1 \end{pmatrix}$$

$$\begin{pmatrix} 1 & -2 \\ 1 & 0 \\ 3 & 1 \end{pmatrix} \rightarrow \begin{pmatrix} \sin \\ \cdot \end{pmatrix} \rightarrow \begin{pmatrix} 1 \\ \exp \end{pmatrix} \rightarrow (1 \quad -1)$$

For

$$\mathbf{f}_n \circ \dots \circ \mathbf{f}_2 \circ \mathbf{f}_1(\mathbf{x}) \in \mathbb{Y}, \mathbf{x} \in [\mathbf{x}](0)$$

the algorithm returns a box $[\mathbf{a}](0) \supset \{\mathbf{x} \in [\mathbf{x}](0) \mid \mathbf{f}(\mathbf{x}) \in \mathbb{Y}\}$.

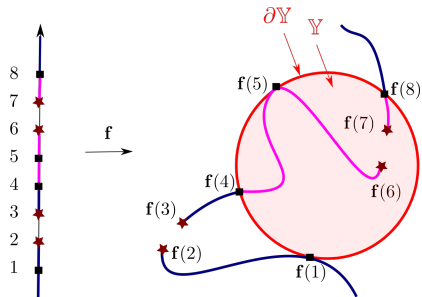
```

Input:  $[\mathbf{x}](0)$ 
1  For  $k = 1$  to  $n$ 
2       $[\mathbf{x}](k) = \vec{\mathcal{C}}_k([\mathbf{x}](k-1))$ 
3   $[\mathbf{a}](n) = [\mathbb{Y} \cap [\mathbf{x}](n)]$ 
4  For  $k = n$  to  $1$ 
5       $[\mathbf{a}](k-1) = \overleftarrow{\mathcal{C}}_k([\mathbf{x}](k-1), [\mathbf{a}](k))$ 
Return  $[\mathbf{a}](0)$ 
    
```


Boundary approach

Consider a continuous function $\mathbf{f}: \mathbb{R}^n \rightarrow \mathbb{R}^p$ defined everywhere. If $\mathbb{X} = \mathbf{f}^{-1}(\mathbb{Y})$, we have

$$\partial \mathbb{X} \subset \mathbf{f}^{-1}(\partial \mathbb{Y}).$$



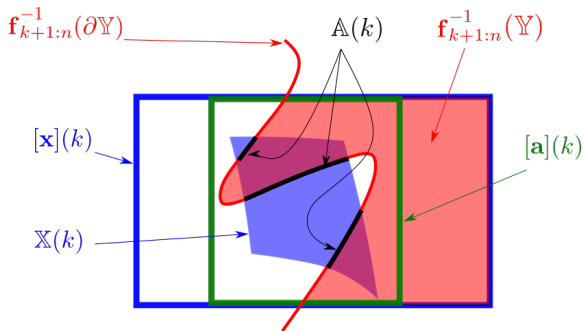
$$\begin{aligned}\mathbb{X} &= \mathbf{f}^{-1}(\mathbb{Y}) = \{1\} \cup [4, 6] \cup [7, 8] \\ \text{Dom}(\mathbf{f}) &=]-\infty, 2] \cup [3, 6] \cup [7, \infty] \\ \partial\mathbb{X} &= \{1, 4, \mathbf{6}, \mathbf{7}, 8\}, \mathbf{f}^{-1}(\partial\mathbb{Y}) = \{1, 4, \mathbf{5}, 8\}\end{aligned}$$

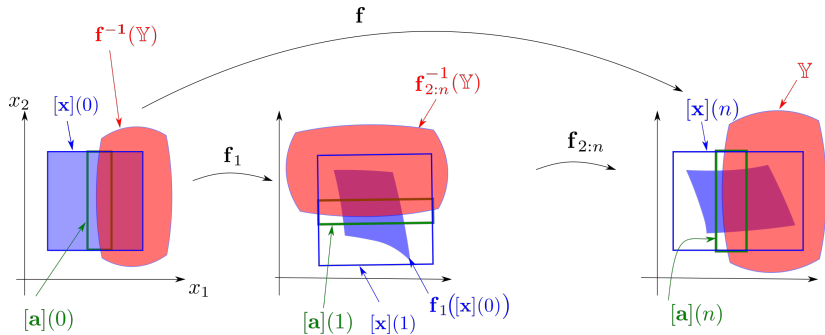
Input: $\mathbb{X}(0)$

- 1 For $k = 1$ to n
- 2 $\mathbb{X}(k) = \mathbf{f}_k(\mathbb{X}(k-1))$
- 3 $\mathbb{A}(n) = \partial \mathbb{Y} \cap \mathbb{X}(n)$
- 4 For $k = n$ to 1
- 5 $\mathbb{A}(k-1) = \mathbb{X}(k-1) \cap \mathbf{f}_k^{-1}(\mathbb{A}(k))$

Return $\mathbb{A}(k-1)$

Boundary forward-backward sequence





Cardinal directions

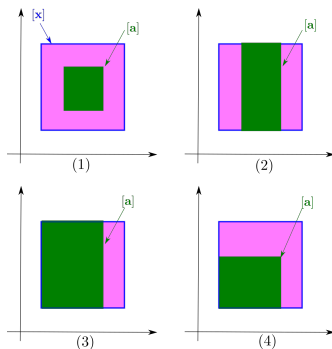
For $\lambda = (2, -)$ and

$$[\mathbf{x}] = [1, 2] \times [3, 4] \times [5, 6].$$

We have $x^\lambda = 3$, the associated face is $[1, 2] \times [3, 3] \times [5, 6]$ and $\mathcal{H}_\lambda([\mathbf{x}]) = \{\mathbf{x} \mid x_2 < 3\}$.

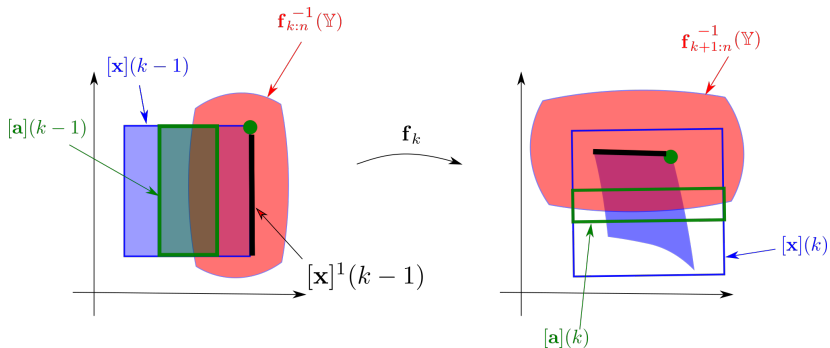
Win boxes. Take $[a] \subset [x]$ and $\lambda \in \mathcal{D}$, the λ th win box, is

$$[x] \setminus [a]_{|\lambda} = [x] \cap \mathcal{H}_{\lambda}([a])$$



Given the pair $[\mathbf{x}](k), [\mathbf{a}](k)$ in the algorithm. We associate to each bound $a^\lambda(k)$ of $[\mathbf{a}](k)$, the quantity $c(a^\lambda(k)) \in \{0, 1, ?\}$ such that

$$\begin{aligned} c(a^\lambda(k)) = 1 &\Rightarrow \mathbf{f}_{k+1:n}([\mathbf{x}]^\lambda) \subset \mathbb{Y} \\ c(a^\lambda(k)) = 0 &\Rightarrow \mathbf{f}_{k+1:n}([\mathbf{x}]^\lambda) \cap \mathbb{Y} = \emptyset \end{aligned}$$



Backward propagation of the bound colors

Sum. $f(\mathbf{x}) = x_1 + x_2$.

Forward step

$$[y] = [x_1] + [x_2]$$

Backward step, $\overleftarrow{\mathcal{C}}([x], [y])$

$$[a] = [x] \cap \begin{pmatrix} [y] - [x_2] \\ [y] - [x_1] \end{pmatrix}$$

and

if $x_1^- < a_1^-$, then $c(a_1^-) = c(y^-)$

if $a_1^+ < x_1^+$, then $c(a_1^+) = c(y^+)$

if $x_2^- < a_2^-$, then $c(a_2^-) = c(y^-)$

if $a_2^+ < x_2^+$, then $c(a_2^+) = c(y^+)$

Square. $f(x) = x^2$

Forward step

$$[y] = [x]^2$$

Backward step

$$[a] = [\{x \in [x], x^2 \in [y]\}]$$

if $x^- < a^-$,

if $x^{-2} < y^-$, then $c(a^-) = c(y^-)$
else $c(a^-) = c(y^+)$

if $a^+ < x^+$,

if $x^{+2} < y^+$, then $c(a^+) = c(y^+)$
else $c(a^+) = c(y^-)$

Test-cases

Test-case 1. Consider the set inversion problem

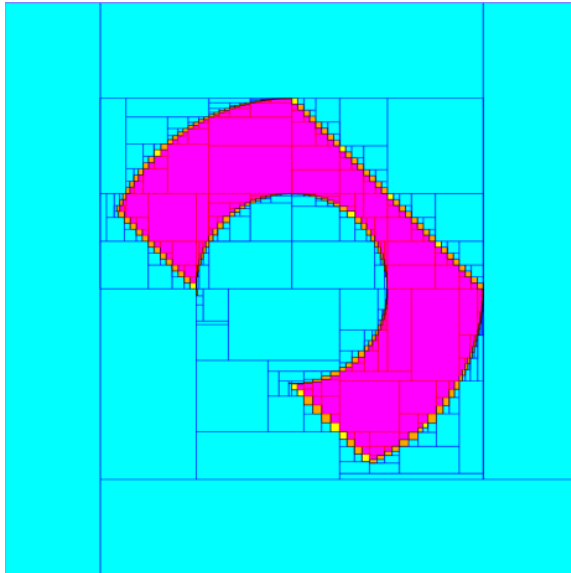
$$\mathbb{X} = \left(\begin{array}{c} [-3,3] \\ [-3,3] \end{array} \right) \cap \mathbf{f}^{-1} \left(\begin{array}{c} [-1,2] \\ [1,4] \end{array} \right)$$

where

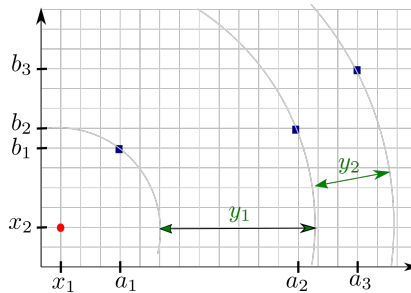
$$\mathbf{f} \left(\begin{array}{c} x_1 \\ x_2 \end{array} \right) = \left(\begin{array}{c} x_1 + x_2 \\ (x_1 + x_2)^2 \end{array} \right).$$

The function $\mathbf{f}$ has the following contractible decomposition:

$$\left(\begin{array}{c} x_1 \\ x_2 \end{array} \right) \rightarrow (s = x_1 + x_2) \rightarrow \left(\begin{array}{c} s \\ s^2 \end{array} \right)$$



Test-case 2. Pseudorange multilateration



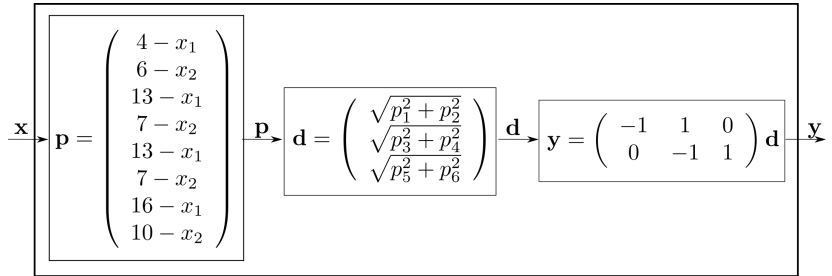
The robot measures the pseudo distances $y_1 = 8, y_2 = 4$ to the stations

We assume that the accuracy of the pseudo distance measurements is $\varepsilon = 0.001$. The set $\mathbb{X}$ of all feasible location vectors is defined by

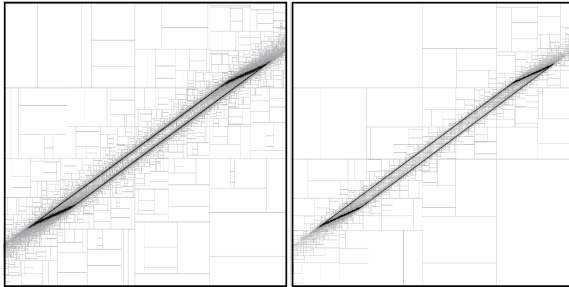
$$\mathbf{f}(\mathbf{x}) \in [8 - \varepsilon, 8 + \varepsilon] \times [4 - \varepsilon, 4 + \varepsilon]$$

where

$$\mathbf{f}(\mathbf{x}) = \begin{pmatrix} \sqrt{(13-x_1)^2 + (7-x_2)^2} - \sqrt{(4-x_1)^2 + (6-x_2)^2} \\ \sqrt{(16-x_1)^2 + (10-x_2)^2} - \sqrt{(13-x_1)^2 + (7-x_2)^2} \end{pmatrix}$$



Decomposition into contractible functions



Classical forward backward contractor generates 90841 boxes.

Our boundary based contractor with the contractible decomposition
generated 35586 boxes.

Perspectives

Directed non-monotonic contractor for the constraint $\mathbf{y} = \mathbf{f}(\mathbf{x})$
 is a pair of two operators

$$\mathcal{C} : ([\mathbf{x}], [\mathbf{y}]) \rightarrow \left(\overrightarrow{\mathcal{C}}([\mathbf{x}]), \overleftarrow{\mathcal{C}}([\mathbf{x}], [\mathbf{y}]) \right)$$

such that

$$\begin{aligned} \mathbf{f}([\mathbf{x}]) &\subset \overrightarrow{\mathcal{C}}([\mathbf{x}]) \\ \mathbf{f}^{-1}([\mathbf{y}] \cap [\mathbf{x}]) &\subset \overleftarrow{\mathcal{C}}([\mathbf{x}], [\mathbf{y}]) \end{aligned}$$

Where $[\mathbf{x}], [\mathbf{y}]$ could be oriented boxes, ellipsoids, etc.
 Lohner type contractors could thus be used.



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Abstract interpretation: A unified lattice model for static analysis of programs by construction or approximation of fixpoints.

In *Conference Record of the Fourth ACM Symposium on Principles of Programming Languages*, pages 238–252, Los Angeles, California, 1977.



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In *IFAC'2017, Toulouse, 2017*.



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