

Cooperative Exploration and Mapping

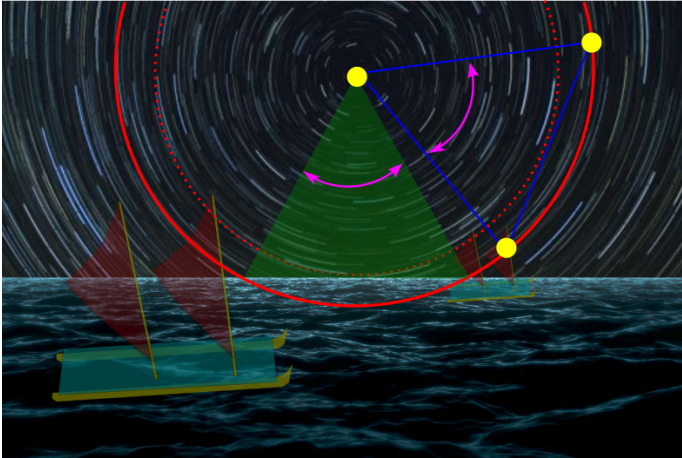
L. Jaulin
ENSTA Bretagne, Robex, LabSTICC
2019, June 10, Brest



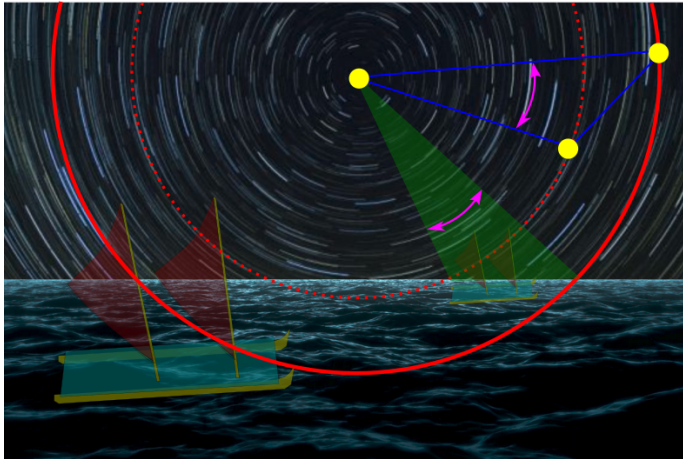
Polynesian navigation



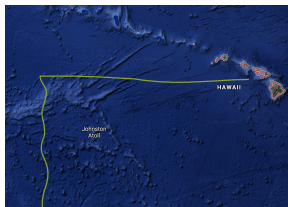
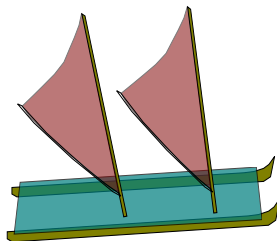
Find the route without GPS, compass and clocks with *wa'a kaulua*[2]

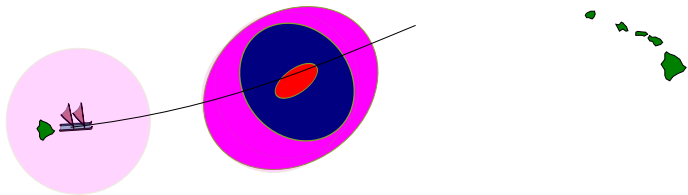


Pair of stars technique

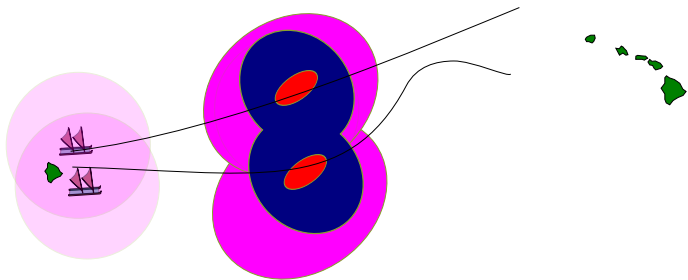


Pair of stars technique





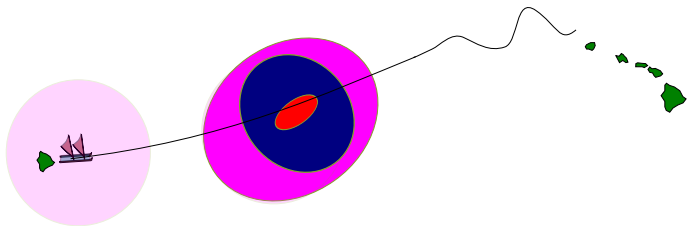
Prove that islands will be reached by one boat



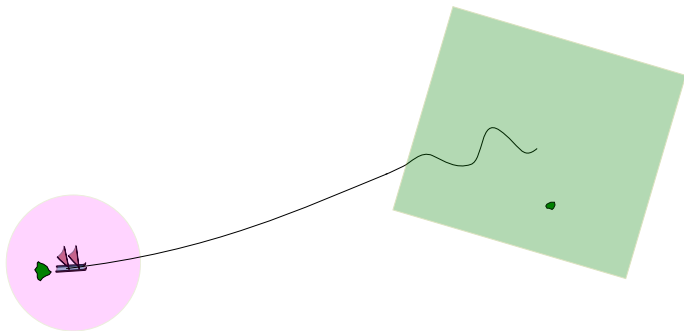
Prove that islands will be reached by the n boats



Alignment to keep the heading in case of clouds



Find a control to reach the geo-localized islands



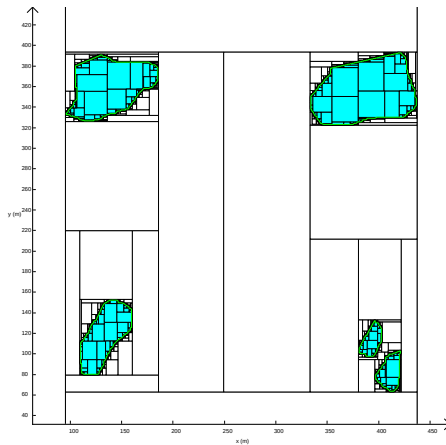
Explore a given area entirely to find new islands

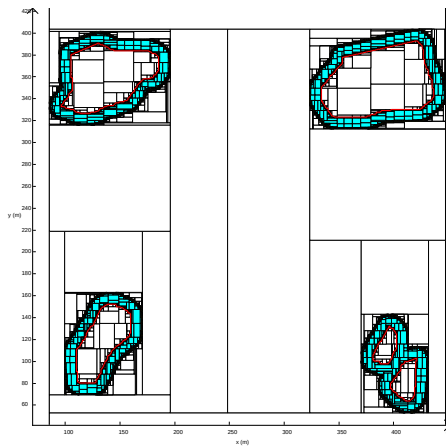
Polynesian exploration problem

- Given a set of unmapped islands $\mathbf{m}_i, i \geq 0$ with coastal area:

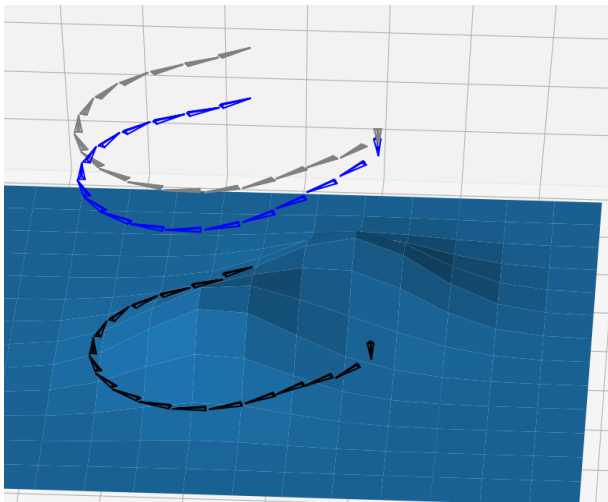
$$\mathbb{C}_i = \{\mathbf{x} \mid c_i(\mathbf{x}) \leq 0\}.$$

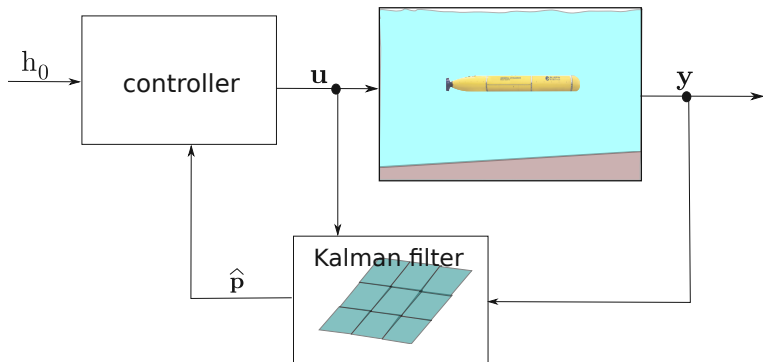
- A robot has to map this environment without being lost.

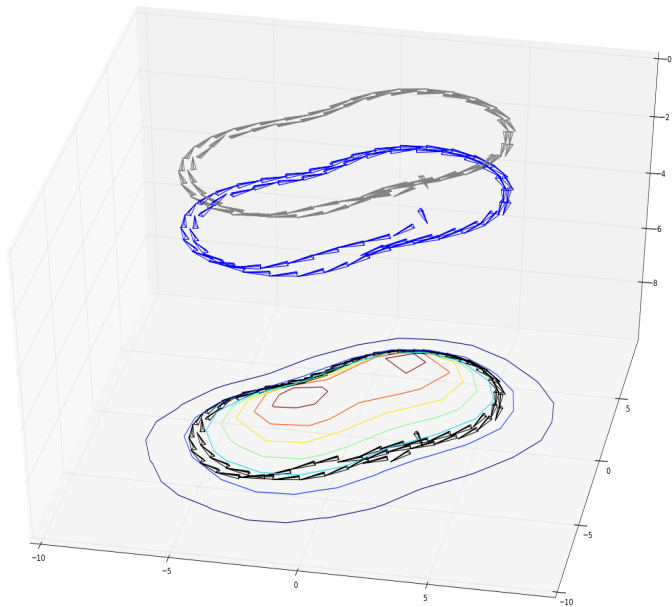




Follow isobaths







Bind Exploration

Visible area

The robot has a state $\mathbf{x}$. The visible area is $\mathbb{V}(\mathbf{x})$

Example. The robot is able to see all up to 3 meters

$$\mathbb{V}(\mathbf{x}) = \left\{ (z_1, z_2) \mid (z_1 - x_1)^2 + (z_2 - x_2)^2 \leq 9 \right\}.$$

Blind exploration

The explored zone $\mathbb{Z}$ is defined by [1]

$$\left\{ \begin{array}{l} \dot{\mathbf{x}} = \mathbf{f}(\mathbf{x}, \mathbf{u}), \mathbf{u}(t) \in [\mathbf{u}](t) \\ \mathbb{Z} = \bigcup_{t \geq 0} \mathbb{V}(\mathbf{x}(t)) \\ \mathbf{x}(0) = \mathbf{x}_0 \end{array} \right.$$

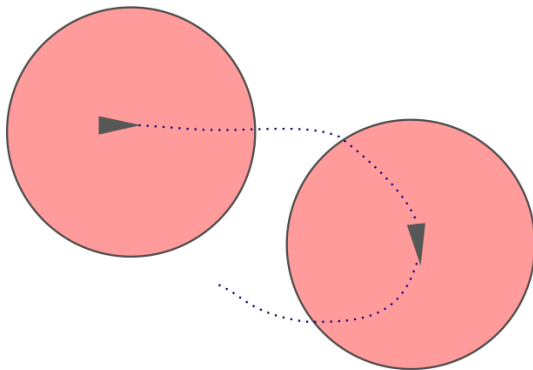
We have

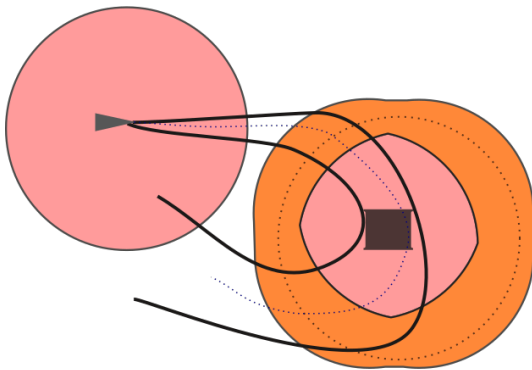
$$\underbrace{\bigcap_{\mathbf{x}(\cdot) \in \mathcal{X}(\cdot)} \bigcup_{t \geq 0} \mathbb{V}(\mathbf{x}(t))}_{\mathbb{Z}^-} \subset \mathbb{Z} \subset \underbrace{\bigcup_{\mathbf{x}(\cdot) \in \mathcal{X}(\cdot)} \bigcup_{t \geq 0} \mathbb{V}(\mathbf{x}(t))}_{\mathbb{Z}^+}.$$

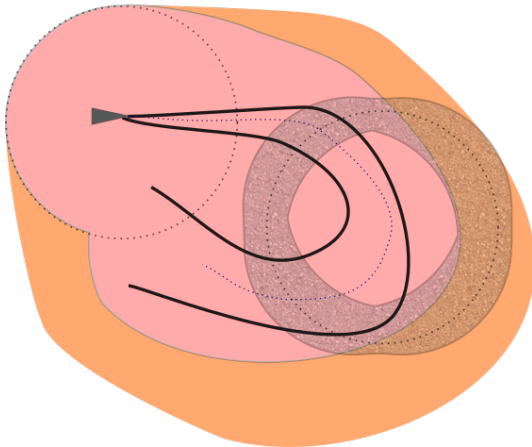
$\mathbb{Z}^-$ is the *certainly explored zone*.

$\mathbb{Z}^+$ is the *maybe explored zone*.

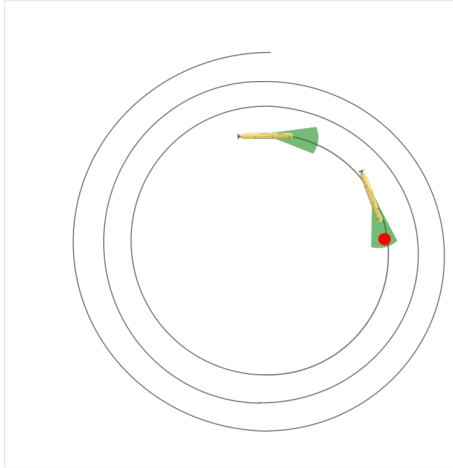
$\mathbb{Z}^+ \setminus \mathbb{Z}^-$ is the *penumbra*.

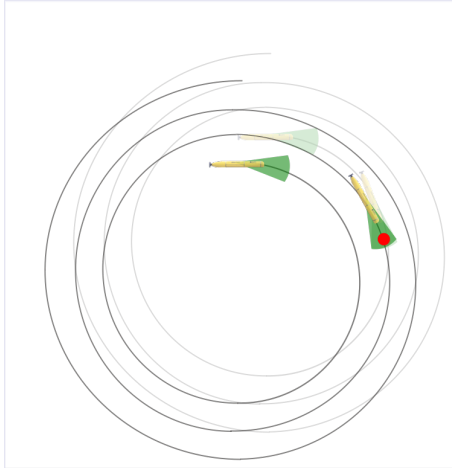


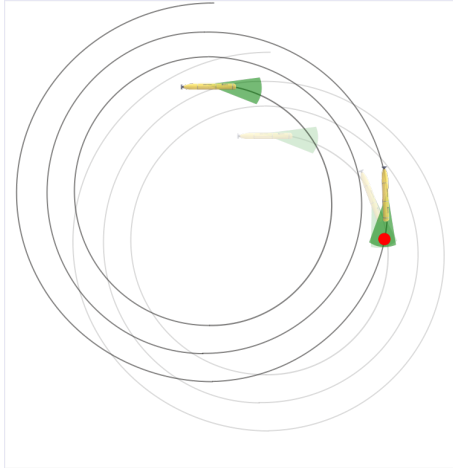


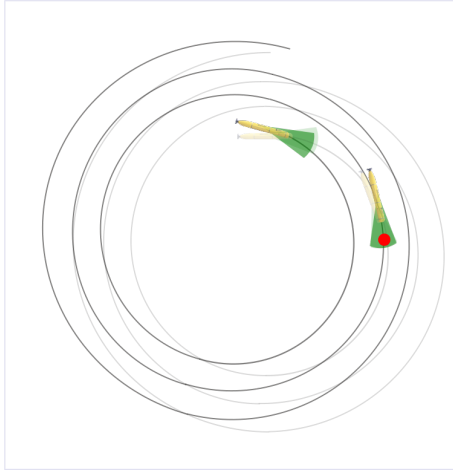


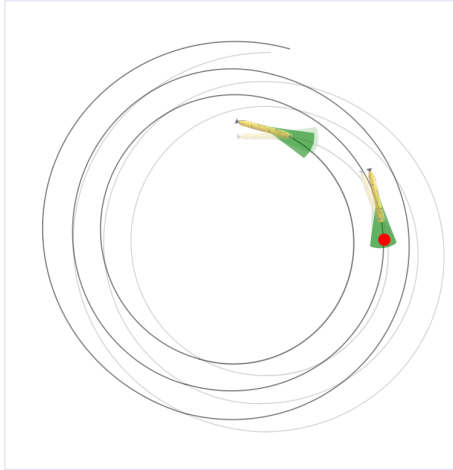
Spiral scan

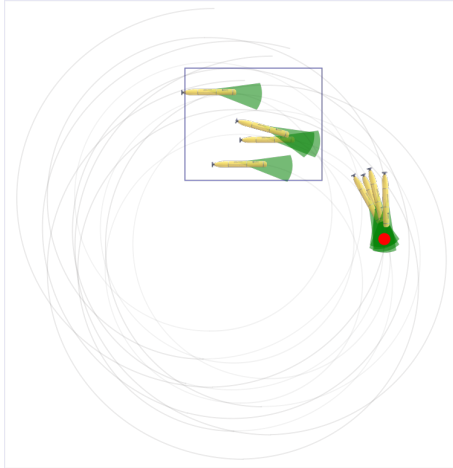






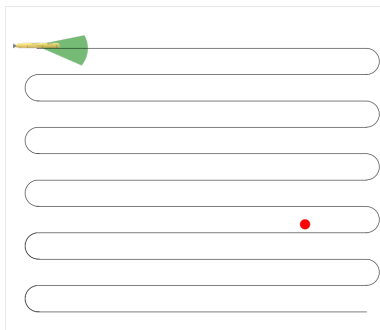


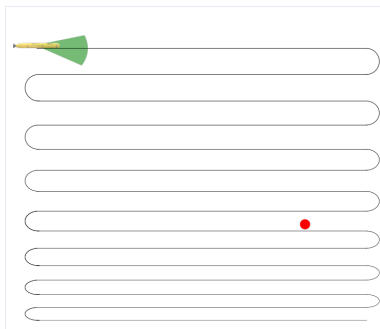




Boustrophedon

Which pattern is the best for exploration?





Reach an island

Robot

The robot is described by

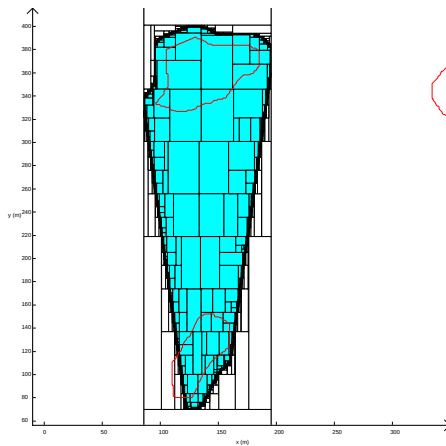
$$\begin{cases} \dot{\mathbf{x}} = \mathbf{f}(\mathbf{x}, \mathbf{u}), & \mathbf{u}(\cdot) \in [\mathbf{u}](t) \\ \mathbf{x}(0) = \mathbf{x}_0 \end{cases}$$

Offshore, the robot is blind and an open loop strategy.

Backward reach set

Given the set $\mathbb{A}$, the backward reach set is defined by

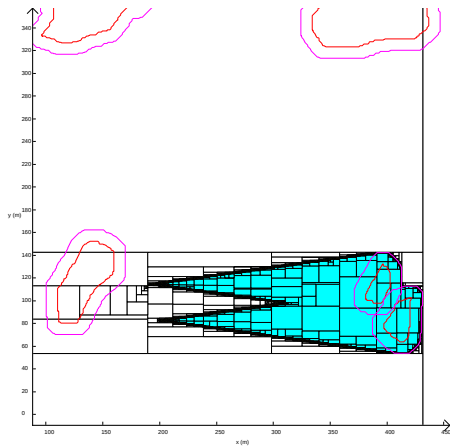
$$\text{Back}(\mathbb{A}) = \{\mathbf{x} \mid \forall \varphi \in \Phi, \exists t \geq 0, \varphi(t, \mathbf{x}) \in \mathbb{A}\}.$$

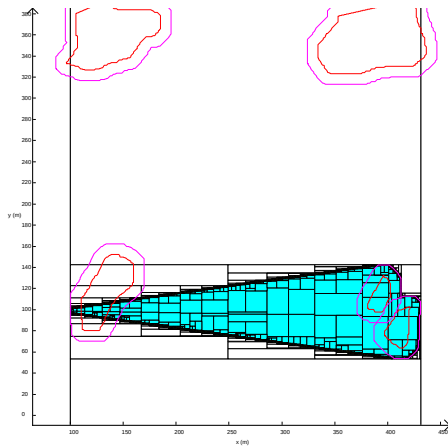


Archipelagic effect

We have

$$\text{Back}(\mathbb{A} \cup \mathbb{B}) \supset \text{Back}(\mathbb{A}) \cup \text{Back}(\mathbb{B})$$





No-lost zone

Moving between coastal zones

- We have m coastal sets $\mathbb{C}_1, \mathbb{C}_2, \dots, i \in \{1, 2, \dots\}$
- We have open loop control strategies $\mathbf{u}_j, j \in \{1, 2, \dots\}$,
- Equivalently, we have set flows $\Phi_j(t, \mathbf{x}_0)$.
- The control strategy cannot change offshore.

Graph

From $\mathbb{C}_1$ we can reach $\mathbb{C}_2$ with the j th control strategy if

$$\mathbb{C}_1 \cap \text{Back}(j, \mathbb{C}_2) \neq \emptyset.$$

From $\mathbb{C}_1$ we can reach $\mathbb{C}_2$ with at least one control strategy if

$$\mathbb{C}_1 \cap \bigcup_j \text{Back}(j, \mathbb{C}_2) \neq \emptyset.$$

From $\mathbb{C}_1$ we can reach $\mathbb{C}_2 \cup \mathbb{C}_3$ with at least one control strategy if

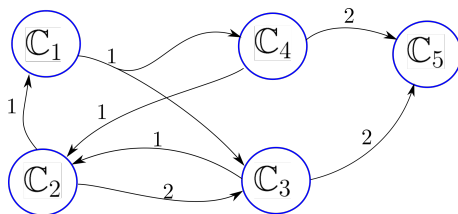
$$\mathbb{C}_1 \cap \bigcup_j \text{Back}(j, \mathbb{C}_2 \cup \mathbb{C}_3) \neq \emptyset.$$

We define $\hookrightarrow$ as:

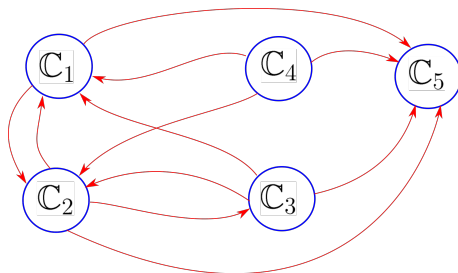
- $\mathbb{C}_a \hookrightarrow \mathbb{C}_b$ if from $\mathbb{C}_a$ we can reach $\mathbb{C}_b$ with at least one control strategy j .
- $\hookrightarrow$ is the smallest transitive relation such that

$$\left\{ \begin{array}{l} \forall k \in \mathbb{K}, \mathbb{C}_{i_k} \hookrightarrow \mathbb{C}_b \\ \exists j, \mathbb{C}_a \cap \text{Back}(j, \bigcup_{k \in \mathbb{K}} \mathbb{C}_{i_k}) \neq \emptyset \end{array} \right. \Rightarrow \mathbb{C}_a \hookrightarrow \mathbb{C}_b$$

Graph



$$\mathbb{C}_1 \cap \text{Back}(1, \mathbb{C}_3 \cup \mathbb{C}_4) \neq \emptyset \Rightarrow \mathbb{C}_1 \rightarrow (\mathbb{C}_3, \mathbb{C}_4)$$



Graph of $\hookrightarrow$

Secure a zone

Secure a zone

INFO OBS. Un sous-marin nucléaire russe repéré dans le Golfe de Gascogne



Le navire a été repéré en janvier. Ce serait la première fois depuis la fin de la Guerre Froide qu'un tel sous-marin, doté de missiles nucléaires, se serait aventuré dans cette zone au large des côtes françaises.



Bay of Biscay 220 000 km²



An intruder

- Several robots $\mathcal{R}_1, \dots, \mathcal{R}_n$ at positions $\mathbf{a}_1, \dots, \mathbf{a}_n$ are moving in the ocean.
- If the intruder is in the visibility zone of one robot, it is detected.[3]

Complementary approach

- We assume that a virtual intruder exists inside $\mathbb{G}$.
- We localize it with a set-membership observer inside $\mathbb{X}(t)$.
- The secure zone corresponds to the complementary of $\mathbb{X}(t)$.

Assumptions

- The intruder satisfies

$$\dot{\mathbf{x}} \in \mathbb{F}(\mathbf{x}(t)).$$

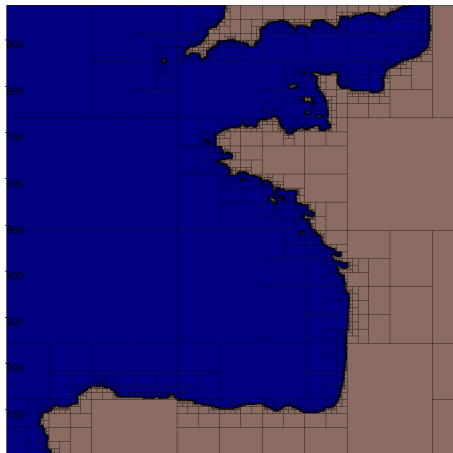
- Each robot $\mathcal{R}_i$ has the visibility zone $g_{\mathbf{a}_i}^{-1}([0, d_i])$ where d_i is the scope.

Theorem. An (undetected) intruder has a state vector $\mathbf{x}(t)$ inside the set

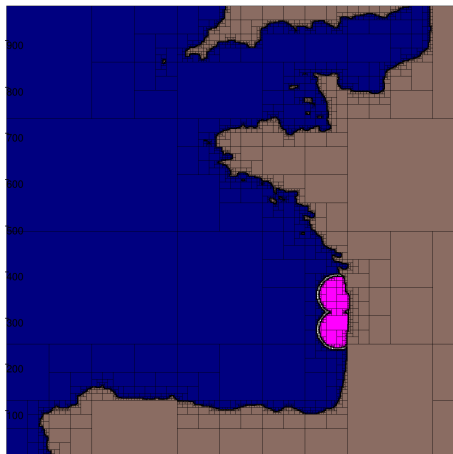
$$\mathbb{X}(t) = \mathbb{G} \cap (\mathbb{X}(t - dt) + dt \cdot \mathbb{F}(\mathbb{X}(t - dt))) \cap \bigcap_i g_{\mathbf{a}_i(t)}^{-1}([d_i(t), \infty]),$$

where $\mathbb{X}(0) = \mathbb{G}$. The secure zone is

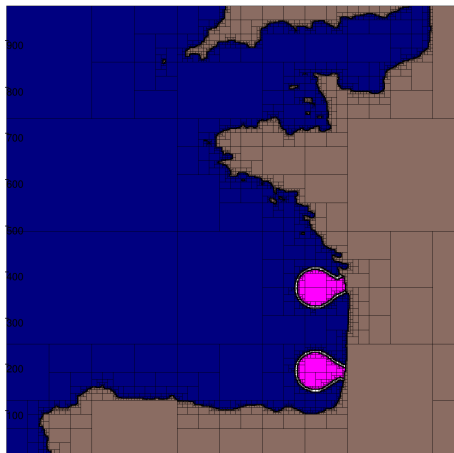
$$\mathbb{S}(t) = \overline{\mathbb{X}(t)}.$$



Set $\mathbb{G}$ in blue

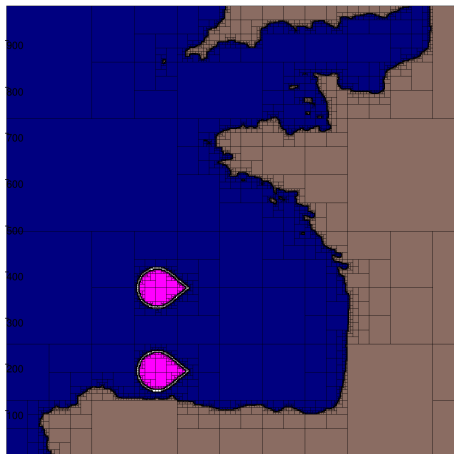


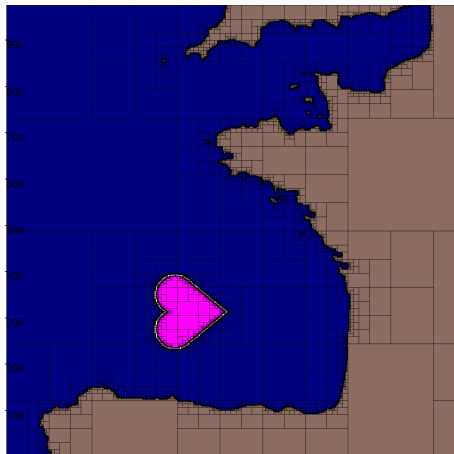
Magenta: $\mathbb{G} \cap \bigcup_i g_{\mathbf{a}_i(t)}^{-1}([0, d_i(t)])$ Blue: $\mathbb{G} \cap \bigcap_i g_{\mathbf{a}_i(t)}^{-1}([d_i(t), \infty])$

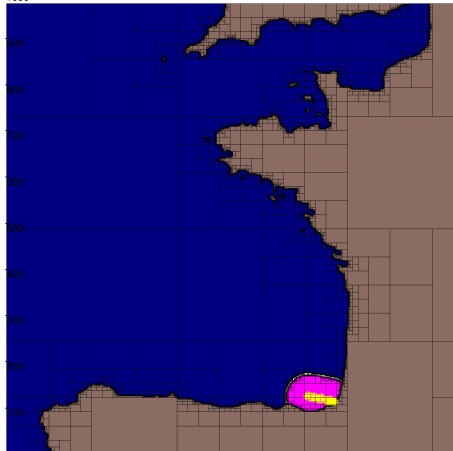


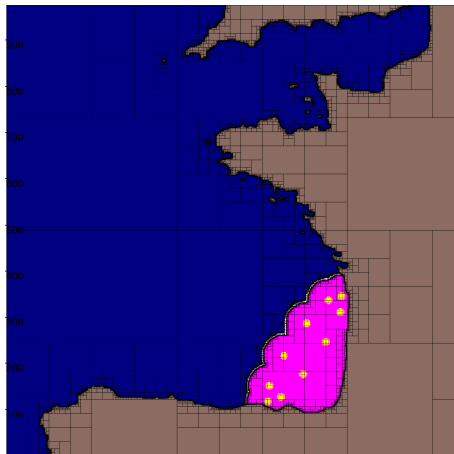
Blue:

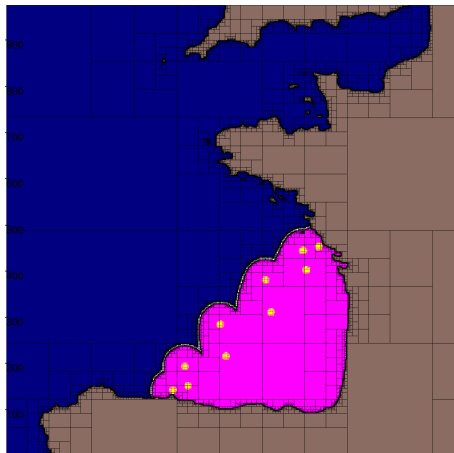
$$\mathbb{X}(t) = \mathbb{G} \cap (\mathbb{X}(t - dt) + dt \cdot \mathbb{F}(\mathbb{X}(t - dt))) \cap \bigcap_i g_{a_i(t)}^{-1}([d_i(t), \infty)).$$

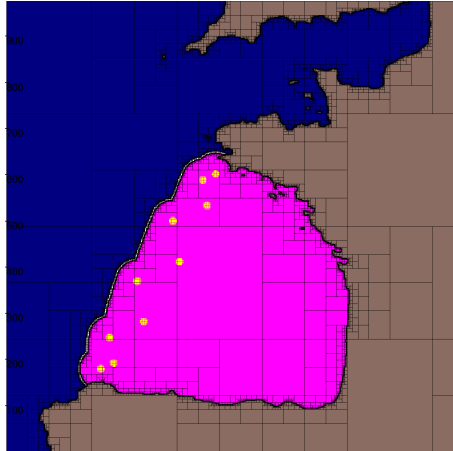












Video : <https://youtu.be/rNcDW6npLfE>

Smoother

Idea: Take into account the future.

The feasible set can be obtained by the following contractions

$$\begin{aligned}\overrightarrow{\mathbb{X}}(t) &= \overrightarrow{\mathbb{X}}(t) \cap (\mathbb{X}(t-dt) + dt \cdot \mathbb{F}(\mathbb{X}(t-dt))) \\ \overleftarrow{\mathbb{X}}(t) &= \overleftarrow{\mathbb{X}}(t) \cap (\mathbb{X}(t+dt) - dt \cdot \mathbb{F}(\mathbb{X}(t+dt))) \\ \mathbb{X}(t) &= \overrightarrow{\mathbb{X}}(t) \cap \overleftarrow{\mathbb{X}}(t)\end{aligned}$$

with the initialization

$$\mathbb{X}(t) = \overrightarrow{\mathbb{X}}(t) = \overleftarrow{\mathbb{X}}(t) = \mathbb{G}.$$



B. Desrochers and L. Jaulin.

Computing a guaranteed approximation the zone explored by a robot.

IEEE Transaction on Automatic Control, 62(1):425–430, 2017.



T. Nico, L. Jaulin, and B. Zerr.

Guaranteed polynesian navigation.

In *SWIM'19, Paris, France*, 2019.



K. Vencatasamy, L. Jaulin, and B. Zerr.

Secure the zone from intruders with a group robots.

Special Issue on Ocean Engineering and Oceanography,
Springer, 2018.