## Intervals for state estimation

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# 1 Interval analysis

# 1.1 Basic notions on set theory

### If f is defined as follows



$$f(A) = \{2,3,4\} = \operatorname{Im}(f).$$
  

$$f^{-1}(B) = \{a,b,c,e\} = \operatorname{dom}(f).$$
  

$$f^{-1}(f(A)) = \{a,b,c,e\} \subset A$$
  

$$f^{-1}(f(\{b,c\})) = \{a,b,c\}.$$

If 
$$f(x) = x^2$$
, then

$$f([2,3]) = [4,9]$$
  
$$f^{-1}([4,9]) = [-3,-2] \cup [2,3].$$

This is consistent with the property

$$f\left(f^{-1}\left(\mathbb{Y}\right)
ight)\subset\mathbb{Y}.$$

### 1.2 Interval arithmetic

$$\begin{aligned} \mathsf{If} \diamond \in \{+,-,\cdot,/,\mathsf{max},\mathsf{min}\} \\ [x] \diamond [y] = \left[\{x \diamond y \mid x \in [x], y \in [y]\}\right]. \end{aligned}$$

For instance,

$$egin{array}{rl} [-1,3]+[2,5]&=[1,8]\ [-1,3]\cdot [2,5]&=[-5,15]\ [-1,3]/[2,5]&=[-rac{1}{2},rac{3}{2}]. \end{array}$$

$$\begin{aligned} [x^-, x^+] + [y^-, y^+] &= & [x^- + y^-, x^+ + y^+]. \\ [x^-, x^+] \cdot [y^-, y^+] &= & [x^- y^- \wedge x^+ y^- \wedge x^- y^+ \wedge x^+ y^+, \\ & & x^- y^- \vee x^+ y^- \vee x^- y^+ \vee x^+ y^+]. \end{aligned}$$

If  $f \in \{\cos, \sin, \operatorname{sqrt}, \log, \exp, \dots\}$  $f([x]) = [\{f(x) \mid x \in [x]\}].$ 

For instance,

$$\begin{array}{rcl} \sin\left([0,\pi]\right) &=& [0,1],\\ \operatorname{sqr}\left([-1,3]\right) &=& [-1,3]^2 = [0,9],\\ \operatorname{abs}\left([-7,1]\right) &=& [0,7],\\ \operatorname{sqrt}\left([-10,4]\right) &=& \sqrt{[-10,4]} = [0,2],\\ \log\left([-2,-1]\right) &=& \emptyset. \end{array}$$

### 1.3 Boxes

A box, or interval vector  $[\mathbf{x}]$  of  $\mathbb{R}^n$  is

 $[\mathbf{x}] = [x_1^-, x_1^+] \times \cdots \times [x_n^-, x_n^+] = [x_1] \times \cdots \times [x_n].$ 

The set of all boxes of  $\mathbb{R}^n$  will be denoted by  $\mathbb{IR}^n$ .

The *principal plane* of [x] is the symmetric plane [x] perpendicular to its largest side.



## 1.4 Inclusion function

The interval function [f] from  $\mathbb{IR}^n$  to  $\mathbb{IR}^m$ , is an *inclusion function* of f if

 $\forall [\mathbf{x}] \in \mathbb{IR}^n, \ \mathbf{f}([\mathbf{x}]) \subset [\mathbf{f}]([\mathbf{x}]).$ 



Inclusion functions [f] and  $[f]^*$ ; here,  $[f]^*$  is minimal.

The natural inclusion function for  $f(x) = x^2 + 2x + 4$  is

$$[f]([x]) = [x]^2 + 2[x] + 4.$$

If [x] = [-3, 4], we have

$$[f]([-3,4]) = [-3,4]^2 + 2[-3,4] + 4$$
  
= [0,16] + [-6,8] + 4  
= [-2,28].

Note that  $f([-3, 4]) = [3, 28] \subset [f]([-3, 4]) = [-2, 28]$ .

#### If ${\bf f}$ is given by the algorithm

Algorithm f(in:  $\mathbf{x} = (x_1, x_2, x_3)$ , out:  $\mathbf{y} = (y_1, y_2)$ ) 1  $z := x_1$ ; 2 for k := 0 to 100 3  $z := x_2(z + kx_3)$ ; 4 next; 5  $y_1 := z$ ; 6  $y_2 := \sin(z \cdot x_1)$ ; Its natural inclusion function is

Algorithm [f](in: [x], out: [y])

 1
 
$$[z] := [x_1];$$

 2
 for  $k := 0$  to 100

 3
  $[z] := [x_2] * ([z] + k * [x_3]);$ 

 4
 next;

 5
  $[y_1] := [z];$ 

 6
  $[y_2] := sin([z] \cdot [x_1]);$ 

Here,  $[{\bf f}]$  is a convergent, thin and monotonic inclusion function for  ${\bf f}.$ 

## 1.5 Subpavings

A subpaving of  $\mathbb{R}^n$  is a set of non-overlapping boxes of  $\mathbb{R}^n$ .

Compact sets  $\mathbb X$  can be bracketed between inner and outer subpavings:

 $\mathbb{X}^{-}\subset\mathbb{X}\subset\mathbb{X}^{+}.$ 

#### Example.

 $\mathbb{X} = \{(x_1, x_2) \mid x_1^2 + x_2^2 \in [1, 2]\}.$ 



Set operations such as  $\mathbb{Z} := \mathbb{X} + \mathbb{Y}, \ \mathbb{X} := \mathbf{f}^{-1}(\mathbb{Y}), \mathbb{Z} := \mathbb{X} \cap \mathbb{Y} \dots$  can be approximated by subpaving operations.

### **1.6 Set inversion**

Let  $\mathbf{f} : \mathbb{R}^n \to \mathbb{R}^m$  and let  $\mathbb{Y}$  be a subset of  $\mathbb{R}^m$ . Set inversion is the characterization of

$$\mathbb{X} = \{\mathbf{x} \in \mathbb{R}^n \mid \mathbf{f}(\mathbf{x}) \in \mathbb{Y}\} = \mathbf{f}^{-1}(\mathbb{Y}).$$

We shall use the following tests.

$$\begin{array}{lll} (\mathsf{i}) & [\mathbf{f}]([\mathbf{x}]) \subset \mathbb{Y} & \Rightarrow & [\mathbf{x}] \subset \mathbb{X} \\ (\mathsf{ii}) & [\mathbf{f}]([\mathbf{x}]) \cap \mathbb{Y} = \emptyset & \Rightarrow & [\mathbf{x}] \cap \mathbb{X} = \emptyset. \end{array}$$

Boxes for which these tests failed, will be bisected, except if they are too small.

Algorithm Sivia(in: [x](0), f, Y) 1  $\mathcal{L} := \{[x](0)\};$ 2 pull [x] from  $\mathcal{L};$ 3 if  $[f]([x]) \subset Y$ , draw([x], 'red'); 4 elseif  $[f]([x]) \cap Y = \emptyset$ , draw([x], 'blue'); 5 elseif  $w([x]) < \varepsilon$ , {draw ([x], 'yellow')}; 6 else bisect [x] and push into  $\mathcal{L};$ 7 if  $\mathcal{L} \neq \emptyset$ , go to 2 If  $\Delta \mathbb{X}$  denotes the union of yellow boxes and if  $\mathbb{X}^-$  is the union of red boxes then :

$$\mathbb{X}^- \subset \mathbb{X} \subset \mathbb{X}^- \cup \Delta \mathbb{X}.$$

# 2 Contractors

To characterize  $\mathbb{X} \subset \mathbb{R}^n$ , bisection algorithms bisect all boxes in all directions and become inefficient. Interval methods can still be useful if

- the solution set X is small (optimization problem, solving equations),
- contraction procedures are used as much as possible,
- bisections are used only as a last resort.

## 2.1 Definition

The operator  $\mathcal{C}_{\mathbb{X}}:\mathbb{IR}^n\to\mathbb{IR}^n$  is a *contractor* for  $\mathbb{X}\subset\mathbb{R}^n$  if

 $\forall [\mathbf{x}] \in \mathbb{IR}^n, \begin{cases} \mathcal{C}_{\mathbb{X}}([\mathbf{x}]) \subset [\mathbf{x}] & (\text{contractance}), \\ \mathcal{C}_{\mathbb{X}}([\mathbf{x}]) \cap \mathbb{X} = [\mathbf{x}] \cap \mathbb{X} & (\text{completeness}). \end{cases}$ 





 $\mathcal{C}_{\mathbb{X}}$  is said to be  $\mathit{convergent}$  if

 $[\mathbf{x}](k) \to \mathbf{x} \quad \Rightarrow \quad \mathcal{C}_{\mathbb{X}}([\mathbf{x}](k)) \to \{\mathbf{x}\} \cap \mathbb{X}.$ 

## 2.2 **Projection of constraints**

Let x, y, z be 3 variables such that

$$egin{array}{rcl} x &\in & [-\infty, 5], \ y &\in & [-\infty, 4], \ z &\in & [6, \infty], \ z &= & x+y. \end{array}$$

Which values for x, y, z are consistent.

Since  $x \in [-\infty, \mathbf{5}], y \in [-\infty, \mathbf{4}], z \in [\mathbf{6}, \infty]$  and z = x + y , we have

$$\begin{array}{rcl} z = x + y \Rightarrow & z \in & [6, \infty] \cap ([-\infty, 5] + [-\infty, 4]) \\ & = [6, \infty] \cap [-\infty, 9] = [6, 9]. \\ x = z - y \Rightarrow & x \in & [-\infty, 5] \cap ([6, \infty] - [-\infty, 4]) \\ & = [-\infty, 5] \cap [2, \infty] = [2, 5]. \\ y = z - x \Rightarrow & y \in & [-\infty, 4] \cap ([6, \infty] - [-\infty, 5]) \\ & = [-\infty, 4] \cap [1, \infty] = [1, 4]. \end{array}$$

The contractor associated with z = x + y is.

<b>Algorithm</b> pplus(inout: $[z], [x], [y]$ )	
1	$[z]:=[z]\cap \left( \left[ x ight] +\left[ y ight]  ight)$ ;
2	$[x]:=[x]\cap \left( \left[ z ight] -\left[ y ight]  ight)$ ;
3	$[y] := [y] \cap ([z] - [x]).$
The projection procedure developed for plus can be extended to other ternary constraints such as mult:  $z = x \cdot y$ , or equivalently

$$\mathsf{mult} \triangleq \left\{ (x, y, z) \in \mathbb{R}^3 \mid z = x \cdot y \right\}.$$

The resulting projection procedure becomes

Algorithm pmult(inout: 
$$[z], [x], [y]$$
)

 1
  $[z] := [z] \cap ([x] \cdot [y]);$ 

 2
  $[x] := [x] \cap ([z] \cdot 1/[y]);$ 

 3
  $[y] := [y] \cap ([z] \cdot 1/[x]).$ 

Consider the binary constraint

$$\exp \triangleq \{(x, y) \in \mathbb{R}^n | y = \exp(x)\}.$$

The associated contractor is

<b>Algorithm</b> pexp(inout: $[y], [x]$ )	
1	$[y]:=[y]\cap \exp\left( \left[ x ight]  ight)$ ;
2	$[x] := [x] \cap \log([y]).$

Any constraint for which such a projection procedure is available will be called a *primitive constraint*.



Projection of the sine constraint

#### 2.3 Solvers

A CSP (Constraint Satisfaction Problem) is composed of 1) a set of variables  $\mathcal{V} = \{x_1, \dots, x_n\}$ ,

- 2) a set of constraints  $C = \{c_1, \ldots, c_m\}$  and
- 3) a set of interval domains  $\{[x_1], \ldots, [x_n]\}$ .

Principle of propagation techniques: contract  $[\mathbf{x}] = [x_1] \times \cdots \times [x_n]$  as follows:

 $(((((([\mathbf{x}] \square c_1) \square c_2) \square \dots) \square c_m) \square c_1) \square c_2) \dots,$ until a steady box is reached. **Example.** Consider the system.

$$y = x^2$$
$$y = \sqrt{x}.$$

We build two contractors

$$\mathcal{C}_{1}: \begin{cases} [y] = [y] \cap [x]^{2} \\ [x] = [x] \cap \sqrt{[y]} \end{cases} \text{ associated to } y = x^{2} \\ \mathcal{C}_{2}: \begin{cases} [y] = [y] \cap \sqrt{[x]} \\ [x] = [x] \cap [y]^{2} \end{cases} \text{ associated to } y = \sqrt{x} \end{cases}$$



















#### **Exemple**. Consider the system

$$\begin{cases} y = 3\sin(x) \\ y = x \end{cases} \quad x \in \mathbb{R}, \ y \in \mathbb{R}.$$





















## 2.4 Decomposition into primitive constraints

$$egin{array}{l} x+\sin(xy)\leq { extsf{0}},\ x\in [-1,1], y\in [-1,1] \end{array}$$

can be decomposed into

$$\left\{ egin{array}{ll} a=xy & x\in [-1,1] & a\in [-\infty,\infty] \ b= \sin(a) &, y\in [-1,1] & b\in [-\infty,\infty] \ c=x+b & c\in [-\infty,0] \end{array} 
ight.$$

# 3 Redermor



The Redermor, GESMA



The *Redermor* at the surface

# Why choosing an interval constraint approach for SLAM ?

- 1) A reliable method is needed.
- 2) The model is nonlinear.
- 3) The pdf of the noises are unknown.
- 4) Reliable error bounds are provided by the sensors.
- 5) A huge number of redundant data are available.

### 3.1 Sensors
A GPS (Global positioning system) at the surface only.

 $t_0 = 6000 \text{ s}, \quad \ell^0 = (-4.4582279^\circ, 48.2129206^\circ) \pm 2.5m$  $t_f = 12000 \text{ s}, \quad \ell^f = (-4.4546607^\circ, 48.2191297^\circ) \pm 2.5m$  **A sonar** (KLEIN 5400 side scan sonar). Gives the distance r between the robot to the detected object.







Screenshot of SonarPro



#### Detection of a mine using SonarPro

**A Loch-Doppler.** Returns the speed of the robot  $\mathbf{v}_r$  and the altitude a of the robot  $\pm 10$ cm.

**A Gyrocompass** (Octans III from IXSEA). Returns the roll  $\phi$ , the pitch  $\theta$  and the head  $\psi$ .

$$\left(egin{array}{c} \phi \ heta \ heta \ \psi \end{array}
ight)\in \left(egin{array}{c} ilde{\phi} \ ilde{ heta} \ ilde{\psi} \end{array}
ight)+\left(egin{array}{c} 1.75 imes10^{-4}.\ [-1,1] \ 1.75 imes10^{-4}.\ [-1,1] \ 5.27 imes10^{-3}.\ [-1,1] \end{array}
ight).$$

#### 3.2 Data

For each time  $t \in \{6000.0, 6000.1, 6000.2, \dots, 11999.4\}$ , we get intervals for

 $\phi(t), \theta(t), \psi(t), v_r^x(t), v_r^y(t), v_r^z(t), a(t).$ 

Six mines have been detected by the sonar:

37.90

36.71

	i		0	1	_	2		3		4		5	
7	$\overline{(i)}$	7054		7092		7374		7748		9038		9688	
c	$\sigma(i)$	1		2		1		0		1		5	
$\hat{i}$	$\tilde{i}(i)$	52.42		12.47		54.40		52.68		27.73		26.98	
_	6		7		0		0		10			11	
_	U		1		ð		9		TO			<b>1</b> 1	
	10024		10817		11172		11232		11279		1	11688	
	4		3		3		4			5		1	

37.37

15.05

33.51

31.03

# 3.3 Constraints satisfaction problem

$$t \in \{6000.0, 6000.1, 6000.2, \dots, 11999.4\},\$$

$$i \in \{0, 1, \dots, 11\},\$$

$$\begin{pmatrix} p_x(t) \\ p_y(t) \end{pmatrix} = 111120 \begin{pmatrix} 0 & 1 \\ \cos\left(\ell_y(t) * \frac{\pi}{180}\right) & 0 \end{pmatrix} \begin{pmatrix} \ell_x(t) - \ell_x^0 \\ \ell_y(t) - \ell_y^0 \end{pmatrix},\$$

$$\mathbf{p}(t) = (p_x(t), p_y(t), p_z(t)),\$$

$$\mathbf{R}_{\psi}(t) = \begin{pmatrix} \cos\psi(t) & -\sin\psi(t) & 0 \\ \sin\psi(t) & \cos\psi(t) & 0 \\ 0 & 0 & 1 \end{pmatrix},\$$

$$\mathbf{R}_{\theta}(t) = \begin{pmatrix} \cos\theta(t) & 0 & \sin\theta(t) \\ 0 & 1 & 0 \\ -\sin\theta(t) & 0 & \cos\theta(t) \end{pmatrix},\$$

$$egin{aligned} \mathbf{R}_arphi(t) &= egin{pmatrix} 1 & 0 & 0 \ 0 & \cos arphi(t) & -\sin arphi(t) \ 0 & \sin arphi(t) & -\sin arphi(t) \ 0 & \sin arphi(t) & \cos arphi(t) \end{pmatrix}, \ \mathbf{R}(t) &= \mathbf{R}_\psi(t).\mathbf{R}_ heta(t).\mathbf{R}_arphi(t), \ \dot{\mathbf{p}}(t) &= \mathbf{R}(t).\mathbf{v}_r(t) \ ||\mathbf{m}(\sigma(i)) - \mathbf{p}( au(i))|| &= r(i), \ \mathbf{R}^\mathsf{T}( au(i)) \left(\mathbf{m}(\sigma(i)) - \mathbf{p}( au(i))\right) \in [0] imes [0,\infty]^{ imes 2}, \ m_z(\sigma(i)) - p_z( au(i)) - a( au(i)) \in [-0.5, 0.5]. \end{aligned}$$

# 3.4 GESMI



# GESMI (Guaranteed Estimation of Sea Mines with Intervals)









# Trajectory reconstructed by GESMI

# 4 SAUC'ISSE



#### Robot SAUC'ISSE



#### Portsmouth, July 12-15, 2007.











# 4.1 Localization with sonar



# 4.2 Set-membership approach

$$\left\{ egin{array}{ll} \mathbf{x}(k+1) &=& \mathbf{f}_k(\mathbf{x}(k),\mathbf{n}\left(k
ight)) \ \mathbf{y}(k) &=& \mathbf{g}_k(\mathbf{x}(k)), \end{array} 
ight.$$

with  $\mathbf{n}(k) \in \mathbb{N}(k)$  and  $\mathbf{y}(k) \in \mathbb{Y}(k)$ .

Without outliers

$$\mathbb{X}(k+1) = \mathbf{f}_k\left(\mathbb{X}(k) \cap \mathbf{g}_k^{-1}\left(\mathbb{Y}(k)\right), \mathbb{N}(k)\right).$$

# 4.3 Relaxed intersection



### 4.4 Robust localization

Define

$$\begin{cases} \mathbf{f}_{k:k} (\mathbb{X}) & \stackrel{\text{def}}{=} \mathbb{X} \\ \mathbf{f}_{k_1:k_2+1} (\mathbb{X}) & \stackrel{\text{def}}{=} \mathbf{f}_{k_2} (\mathbf{f}_{k_1:k_2} (\mathbb{X}), \mathbb{N} (k_2)), \ k_1 \leq k_2. \end{cases}$$
  
The set  $\mathbf{f}_{k_1:k_2} (\mathbb{X})$  represents the set of all  $\mathbf{x} (k_2)$ , consis-

tent with  $\mathbf{x}(k_1) \in \mathbb{X}$ .

Consider the set state estimator

$$\begin{cases} \mathbb{X}(k) = \mathbf{f}_{0:k}(\mathbb{X}(0)) & \text{if } k < m, \text{ (initialization step)} \\ \mathbb{X}(k) = \mathbf{f}_{k-m:k}(\mathbb{X}(k-m)) \cap \\ \{q\} \\ \bigcap_{i \in \{1,...,m\}} \mathbf{f}_{k-i:k} \circ \mathbf{g}_{k-i}^{-1}(\mathbb{Y}(k-i)) & \text{if } k \ge m \end{cases}$$



# 4.5 Application to localization


Sauc'isse robot inside a swimming pool

The robot evolution is described by

$$\begin{cases} \dot{x}_1 = x_4 \cos x_3 \\ \dot{x}_2 = x_4 \sin x_3 \\ \dot{x}_3 = u_2 - u_1 \\ \dot{x}_4 = u_1 + u_2 - x_4, \end{cases}$$

where  $x_1, x_2$  are the coordinates of the robot center,  $x_3$  is its orientation and  $x_4$  is its speed. The inputs  $u_1$  and  $u_2$ are the accelerations provided by the propellers. The system can be discretized by  $\mathbf{x}_{k+1} = \mathbf{f}_k\left(\mathbf{x}_k
ight)$  , where,

$$\mathbf{f}_{k}\begin{pmatrix}x_{1}\\x_{2}\\x_{3}\\x_{4}\end{pmatrix} = \begin{pmatrix}x_{1}+\delta.x_{4}.\cos(x_{3})\\x_{2}+\delta.x_{4}.\sin(x_{3})\\x_{3}+\delta.(u_{2}(k)-u_{1}(k))\\x_{4}+\delta.(u_{1}(k)+u_{2}(k)-x_{4})\end{pmatrix}$$



Underwater robot moving inside a pool





#### https://youtu.be/c-8ZW8nUh7U

# 5 Scout project



**Goal** : (i) coordination of underwater robots ; (ii) collaborative behavior.

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# 5.1 Simulator



#### MOOS, MORSE, BLENDER, IBEX, Vibes, GIT



MOOS architecture

# 5.2 Controller



#### 5.3 Localization

Range only Based on interval analysis Robust with respect to outliers Distributed computation Low rate communication



Presentation of the scout project

http://youtu.be/ATPabRHz0LA

#### 5.4 Tests



# www.ensta-bretagne.fr/jaulin/easibex.html