Interval analysis for proving properties of dynamical systems; application to sailboat robotics Luc Jaulin, Labsticc, IHSEV, OSM, ENSTA-Bretagne http://www.ensta-bretagne.fr/jaulin/ March 5, 2014, Cardiff. March 5, 2014. Mathematics Colloquium



1 Interval analysis

Problem. Given $f : \mathbb{R}^n \to \mathbb{R}$ and a box $[\mathbf{x}] \subset \mathbb{R}^n$, prove that

$$\forall \mathbf{x} \in [\mathbf{x}], f(\mathbf{x}) \geq \mathbf{0}.$$

Interval arithmetic can solve efficiently this problem.

Example. Is the function

 $f(\mathbf{x}) = x_1 x_2 - (x_1 + x_2) \cos x_2 + \sin x_1 \cdot \sin x_2 + 2$ always positive for $x_1, x_2 \in [-1, 1]$? Interval arithmetic

$$\begin{array}{ll} [-1,3]+[2,5] &= [1,8], \\ [-1,3]\cdot [2,5] &= [-5,15], \\ & {\rm abs}\left([-7,1]\right) &= [0,7] \end{array}$$

The interval extension of

$$f(x_1, x_2) = x_1 \cdot x_2 - (x_1 + x_2) \cdot \cos x_2 + \sin x_1 \cdot \sin x_2 + 2$$

is

$$[f]([x_1], [x_2]) = [x_1] \cdot [x_2] - ([x_1] + [x_2]) \cdot \cos [x_2] + \sin [x_1] \cdot \sin [x_2] + 2.$$

Theorem (Moore, 1970)

 $[f]([\mathbf{x}]) \subset \mathbb{R}^+ \Rightarrow \forall \mathbf{x} \in [\mathbf{x}], f(\mathbf{x}) \ge \mathbf{0}.$

2 Computing with sets

Subsets $\mathbb{X} \subset \mathbb{R}^n$ can be bracketted by subpavings :



which can be obtained using interval calculus

Example.

 $\mathbb{X} = \{ (x_1, x_2) \mid x_1^2 + x_2^2 + \sin(x_1 + x_2) \in [4, 9] \}.$



Contractors

The operator \mathcal{C} : $\mathbb{IR}^n \to \mathbb{IR}^n$ is a *contractor* for the equation $f(\mathbf{x}) = 0$, if

 $\left\{ \begin{array}{ll} \mathcal{C}([\mathbf{x}]) \subset [\mathbf{x}] & (\text{contractance}) \\ \mathbf{x} \in [\mathbf{x}] \text{ and } f(\mathbf{x}) = \mathbf{0} \Rightarrow \mathbf{x} \in \mathcal{C}([\mathbf{x}]) & (\text{consistence}) \end{array} \right.$

Example. Consider the primitive equation:

$$x_2 = \sin x_1.$$









$\begin{array}{lll} \mathcal{C} \text{ is monotonic if } & [\mathbf{x}] \subset [\mathbf{y}] \Rightarrow \mathcal{C}([\mathbf{x}]) \subset \mathcal{C}([\mathbf{y}]) \\ \mathcal{C} \text{ is idempotent if } & \mathcal{C}\left(\mathcal{C}([\mathbf{x}])\right) = \mathcal{C}([\mathbf{x}]) \end{array}$

Contractor algebra

intersection	$\left(\mathcal{C}_{1}\cap\mathcal{C}_{2} ight)\left(\left[\mathbf{x} ight] ight)\overset{def}{=}\mathcal{C}_{1}\left(\left[\mathbf{x} ight] ight)\cap\mathcal{C}_{2}\left(\left[\mathbf{x} ight] ight)$
union	$\left(\mathcal{C}_{1}\cup\mathcal{C}_{2} ight)\left(\left[\mathbf{x} ight] ight)\stackrel{def}{=}\left[\mathcal{C}_{1}\left(\left[\mathbf{x} ight] ight)\cup\mathcal{C}_{2}\left(\left[\mathbf{x} ight] ight) ight]$
composition	$(\mathcal{C}_1 \circ \mathcal{C}_2)([\mathbf{x}]) \stackrel{def}{=} \mathcal{C}_1(\mathcal{C}_2([\mathbf{x}]))$
reiteration	$\mathcal{C}^{\infty} \stackrel{def}{=} \mathcal{C} \circ \mathcal{C} \circ \mathcal{C} \circ \ldots$

Contractor associated with a database

The robot with coordinates (x_1, x_2) is in the water.





Building contractors for equations

Consider the primitive equation

 $x_1 + x_2 = x_3$

with $x_1 \in [x_1]$, $x_2 \in [x_2]$, $x_3 \in [x_3]$.

We have

 The contractor associated with $x_1 + x_2 = x_3$ is thus

$$\mathcal{C}\begin{pmatrix} [x_1]\\ [x_2]\\ [x_3] \end{pmatrix} = \begin{pmatrix} [x_1] \cap ([x_3] - [x_2])\\ [x_2] \cap ([x_3] - [x_1])\\ [x_3] \cap ([x_1] + [x_2]) \end{pmatrix}$$

4 Solver

Example. Solve the system

$$y = x^2$$
$$y = \sqrt{x}.$$

We build two contractors

$$\mathcal{C}_{1}: \begin{cases} [y] = [y] \cap [x]^{2} \\ [x] = [x] \cap \sqrt{[y]} \end{cases} \text{ associated with } y = x^{2} \\ \mathcal{C}_{2}: \begin{cases} [y] = [y] \cap \sqrt{[x]} \\ [x] = [x] \cap [y]^{2} \end{cases} \text{ associated with } y = \sqrt{x} \end{cases}$$


















Sailboat robotics











6 Vaimos

 $\label{eq:collaboration} Collaboration \ ENSTA/IFREMER$



Vaimos à la WRSC (ENSTA-IFREMER-Ecole Navale).



The robot satisfies a state equation

$$\dot{\mathbf{x}} = \mathbf{f}(\mathbf{x}, \mathbf{u})$$
.

With the controller $\mathbf{u}=\mathbf{g}\left(\mathbf{x}\right)$, the robot satisfies

$$\dot{\mathbf{x}} = \mathbf{f}(\mathbf{x})$$
.

With all uncertainties, the robot satisfies.

 $\dot{\mathbf{x}} \in \mathbf{F}\left(\mathbf{x}
ight)$

which is a differential inclusion.

7 Line following



Controller of a sailboat robot



Heading controller

$$\begin{cases} \delta_r &= \frac{\delta_r^{\max}}{\pi} \operatorname{.atan}(\tan \frac{\theta - \overline{\theta}}{2}) \\ \delta_s^{\max} &= \frac{\pi}{2} \cdot \left(\frac{\cos(\psi - \overline{\theta}) + 1}{2} \right). \end{cases}$$

Rudder

$$\left\{ \delta_r = \frac{\delta_r^{\max}}{\pi} \operatorname{.atan}(\tan \frac{\theta - \overline{\theta}}{2}) \right\}$$



Sail



7.1 Vector field





A course θ^* may be unfeasible



Keep close hauled strategy.

7.2 Controller

Controlleur : in:
$$\mathbf{m}, \theta, \psi, \mathbf{a}, \mathbf{b}$$
; out: $\delta_r, \delta_s^{\max}$; inout: q
1 $e = \frac{\det(\mathbf{b}-\mathbf{a},\mathbf{m}-\mathbf{a})}{\|\mathbf{b}-\mathbf{a}\|}$
2 if $|e| > \frac{r}{2}$ then $q = \operatorname{sign}(e)$
3 $\overline{\theta} = \operatorname{atan2}(\mathbf{b}-\mathbf{a}) - \frac{1}{2} \cdot \operatorname{atan}\left(\frac{e}{r}\right)$
4 if $\cos\left(\psi - \overline{\theta}\right) + \cos\zeta < 0$ then $\overline{\theta} = \pi + \psi - q.\zeta$.
5 $\delta_r = \frac{\delta_r^{\max}}{\pi} \cdot \operatorname{atan}(\tan\frac{\theta-\overline{\theta}}{2})$
6 $\delta_s^{\max} = \frac{\pi}{2} \cdot \left(\frac{\cos(\psi-\overline{\theta})+1}{2}\right)$.

8 Validation by simulation



9 Theoretical validation

*Jaulin, Le Bars (2012). An interval approach for stability analysis; Application to sailboat robotics. IEEE TRO.

When the wind is known, the sailboat with the heading controller is described by

$$\dot{\mathbf{x}} = \mathbf{f}(\mathbf{x})$$
.

The system

$$\dot{\mathbf{x}} = \mathbf{f}(\mathbf{x})$$

is Lyapunov-stable (1892) is there exists $V\left(\mathbf{x}
ight)\geq$ 0 such that

$$\dot{V}(\mathbf{x}) < 0 \text{ if } \mathbf{x} \neq \mathbf{0},$$

 $V(\mathbf{x}) = 0 \text{ iff } \mathbf{x} = \mathbf{0}.$

Definition. Consider a differentiable function $V(\mathbf{x}) : \mathbb{R}^n \to \mathbb{R}$. The system $\dot{\mathbf{x}} = \mathbf{f}(\mathbf{x})$ is V-stable if

$$\left(V\left(\mathbf{x}\right) \geq \mathbf{0} \;\Rightarrow\; \dot{V}\left(\mathbf{x}\right) \leq \varepsilon < \mathbf{0}\right).$$



Theorem. If the system $\dot{\mathbf{x}} = \mathbf{f}(\mathbf{x})$ is *V*-stable then

(i) $\forall \mathbf{x}(0), \exists t \geq 0$ such that $V(\mathbf{x}(t)) < 0$ (ii) if $V(\mathbf{x}(t)) < 0$ then $\forall \tau > 0, V(\mathbf{x}(t+\tau)) < 0$. Now,

$$\begin{pmatrix} V(\mathbf{x}) \ge \mathbf{0} \implies \dot{V}(\mathbf{x}) < \mathbf{0} \\ \Leftrightarrow \quad \left(V(\mathbf{x}) \ge \mathbf{0} \Rightarrow \frac{\partial V}{\partial \mathbf{x}}(\mathbf{x}) \cdot \mathbf{f}(\mathbf{x}) < \mathbf{0} \right) \\ \Leftrightarrow \quad \forall \mathbf{x}, \frac{\partial V}{\partial \mathbf{x}}(\mathbf{x}) \cdot \mathbf{f}(\mathbf{x}) < \mathbf{0} \text{ or } V(\mathbf{x}) < \mathbf{0} \\ \Leftrightarrow \quad \neg \left(\exists \mathbf{x}, \frac{\partial V}{\partial \mathbf{x}}(\mathbf{x}) \cdot \mathbf{f}(\mathbf{x}) \ge \mathbf{0} \text{ and } V(\mathbf{x}) \ge \mathbf{0} \right)$$

Theorem. We have

$$\begin{cases} \frac{\partial V}{\partial \mathbf{x}}\left(\mathbf{x}\right).\mathbf{f}\left(\mathbf{x}\right) \geq \mathbf{0} \\ V(\mathbf{x}) \geq \mathbf{0} \end{cases} \text{ inconsistent } \Leftrightarrow \mathbf{\dot{x}} = \mathbf{f}\left(\mathbf{x}\right) \text{ is } V \text{-stable.} \end{cases}$$

Interval method could easily prove the $V\mbox{-stability}.$

Theorem. We have

 $\begin{cases} \frac{\partial V}{\partial \mathbf{x}}(\mathbf{x}) . \mathbf{a} \ge \mathbf{0} \\ \mathbf{a} \in \mathbf{F}(\mathbf{x}) \\ V(\mathbf{x}) \ge \mathbf{0} \end{cases} \text{ inconsistent } \Leftrightarrow \mathbf{\dot{x}} \in \mathbf{F}(\mathbf{x}) \text{ is } V \text{-stable} \end{cases}$



Differential inclusion $\dot{\mathbf{x}} \in \mathbf{F}(\mathbf{x})$ for the sailboat. $V(x) = x_2^2 - r_{\max}^2$.
10 Experimental validation

Collaboration ENSTA-Ifremer. Fabrice Le Bars, Olivier Ménage, Patrick Rousseau, ...



Rade de Brest

Brest-Douarnenez. January 17, 2012, 8am















Middle of Atlantic ocean, 350 km made by Vaimos in 53h, September 6-9, 2012.

Consequence.

It is possible for a sailboat robot to navigate inside a corridor.

Essential, to create circulation rules when robot swarms are considered.

Essential to determine who has to pay in case of accident.