## Internship report

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#### Abstract

This report describes my work during my internship in the Åland Sailing Robots research team of the Åland University of Applied Science.

The main goal of this internship was to design a deliberative collision avoidance algorithm. This algorithm should avoid the collision with the other boats, and avoid our boat to cross the borders of a given sailing zone. It uses an interval analysis method. It is divided in two main steps. The detection step check if there is any risk of collision or risk of crossing a border of the sailing zone. If there is, the path replanning step start. It aims to choose a new path, which avoid every obstacle, and in the same time assure that the boat doesn't cross the border of the sailing zone.

The end of my internship matched with the participation of the team to the World Robotic Sailing Championship. I also took part in the preparation of this competition.


## Context

My internship took place in the Åland Sailing Robots research team of the Åland University of Applied Science. The goal of this research team is to develop an autonomous sailboat, in order to develop Åland islands' competency and innovation regarding unmanned and autonomous shipping and green technology.

They are currently working on the ASPire project, which aims to develop a marine research platform. This platform takes the shape of a 4 meters sailboat, with a free rotating wingsail, which consists of a mainwing and a tailwing, instead of a classic soft sail. The goal of this platform is to carry out some long survey missions in the Baltic sea, in order to gather data on the water composition. This project has currently funding for three years, and at the time of my internship the project was at the half of its development.


Figure 1: ASPire on her christening on July 2017

## Goals \& concerns

The main goal of my internship was to design a deliberative collision avoidance algorithm. Indeed, during her missions, ASPire could meet some other boats, such as ferries, containers ships, or pleasure boats. In order to assure her own safety and the safety of the
other boats, she has to be able to avoid them somehow. She also must stay in a defined sailing zone, in order to avoid coasts and shallow waters.

A reactive collision avoidance algorithm has already been implemented. It is based on a voter's system, which decide at a given frequency toward which direction the boat should aim at. But this is most adapted for short range collision avoidance, because at long range this algorithm can't assure to reach the destination point in an efficient way.

For long range collision avoidance, the idea is to use a deliberative algorithm. It must take into account all the information about every obstacle and allowed sailing zones, and decide if it can follow the current path or if it has to change it. Currently, it is still at a prototype and simulation stage, and it isn't linked to the boat control system yet.

The end of my internship matched with our participation in the World Robotics Sailing Championship (WRSC) in Horten, Norway. In order to prepare this competition, I contributed to the development of the control system of the vessel and of its simulator. I was also involved in the on-sea tests of this control system, before and during this competition.

## Internship realisation

## I) Deliberative collision avoidance

A) Introduction

The control system of ASPire is organized on several levels. First, the lower level controller has for inputs the desired course and the current course, and sends to the rudder actuator and the tailwing actuator a command. Above that, we have currently two possibilities available. The first option is a line following algorithm that takes for inputs two waypoints, and sends to the lower level controller the desired course in order to follow the line between the two waypoints. It considers the wind direction in order to allow the boat to tack when needed. The other option is the reactive collision avoidance based on a voter's system. It considers the wind, the detected obstacles, and the next waypoint to determine a target course. The list of every waypoints is stored in a database as a list of coordinates. A waypoint manager checks when the next waypoint is reached, and if so, sends the new waypoint to the line following algorithm.


Figure 2: simplified diagram of the control system of ASPire
The idea of the deliberative collision avoidance is to read the list of waypoints, and check if there is any collision with another boat or if the boat crosses the border of a given sailing zone. If so, it should modify this list, that is, modify the position of one or several waypoints, or even add a new waypoint, in order to avoid this collision. We have some information about the surrounding vessels thanks to the AIS (Automatic Identification System). This device receives some information about the other boats, such as their GPS position, speed and course.

## B) Collision detection

The main idea is extracted from the last part of the article A Simple Controller for Line Following of Sailboats [2], which propose a collision detection algorithm based on interval analysis.

The assumptions of this algorithm are the following. We assume that the earth is approximated by a plane, with a fixed cartesian frame Oxy. We consider that our boat detects $m$ other boats at the time $t=0$. We assume that we measure the speed, the course and the position of these boats with a known accuracy thanks to the AIS. The speed is considered as constant. The position $\boldsymbol{m}^{i}(t)$ of the $i^{\text {th }}$ detected boat is assumed to satisfy:

$$
\boldsymbol{m}^{i}(t)=\boldsymbol{a}^{i} . t+\boldsymbol{b}^{i}, \quad i \in\{1, \ldots m\}
$$

Where $\boldsymbol{a}^{\boldsymbol{i}}$ and $\boldsymbol{b}^{\boldsymbol{i}}$ are vectors of $\mathbb{R}^{2}$ which correspond to the speed vector and the initial position of the $i^{\text {th }}$ boat. The AIS doesn't return the speed vector, but a projection allows us to switch from the speed and course to the speed vector. As we know those two quantities with a known accuracy, we have two boxes $\left[\boldsymbol{a}^{\boldsymbol{i}}\right]$ and $\left[\boldsymbol{b}^{\boldsymbol{i}}\right]$ which enclose $\boldsymbol{a}^{\boldsymbol{i}}$ and $\boldsymbol{b}^{\boldsymbol{i}}$ respectively. The information about all the obstacles are stored in a list.

Concerning our boat, we assume that its trajectory is the following:

$$
\boldsymbol{m}^{\mathbf{0}}(t)=\boldsymbol{a}^{\mathbf{0}} \cdot t+\boldsymbol{b}^{\mathbf{0}}
$$

The two vectors $\boldsymbol{a}^{\mathbf{0}}$ and $\boldsymbol{b}^{\mathbf{0}}$ depend on the line to be followed by our boat, so on the position of the two waypoints that we want to link, on the speed of our boat, and on its initial position. In a first approximation, we will assume that the initial position of the boat is on the $1^{\text {st }}$ waypoint. We also assume that the waypoints coordinates are expressed in our fixed cartesian frame. We assume that we know the progression speed of the boat on the followed line. In many case, it will correspond to the GPS speed projected on the course, but when she tacks it won't be the case. however, it will be possible to make an estimation. Again, we can easily switch from the course and speed to the speed vector with a projection. Here again, we will consider that we have two boxes $\left[\boldsymbol{a}^{\mathbf{0}}\right]$ and $\left[\boldsymbol{b}^{\mathbf{0}}\right]$ which enclose $\boldsymbol{a}^{\mathbf{0}}$ and $\boldsymbol{b}^{\mathbf{0}}$.

In reality, the initial position of the boat may not be on the $1^{\text {st }}$ waypoint. So, some adjustment should be made in the current algorithm, for example create a waypoint on the current position of the boat in the initialization of the algorithm. In any case, for the good working of the method, we need to make sure that our boat is always somewhere on the line between the two waypoints.

Then the article proposes a method based on Interval Analysis. It aims to prove that the robot's trajectory is safe. The idea is to check in a time interval [ $0, t_{\text {max }}$ ] if there is a collision. There is a collision in our trajectory if:

$$
\exists i \in\{1, \ldots m\}, \exists t \in\left[0, t_{\max }\right], \boldsymbol{m}^{i}(t)=\boldsymbol{m}^{0}(t)
$$

That is if the following system has a solution:

$$
\left\{\begin{array}{c}
\left(a^{0}-a^{i}\right) \cdot t+b^{0}-b^{i}=0, \\
t \in\left[0, t_{\max }\right], i \in[1, \ldots m] \\
a^{0} \in\left[a^{0}\right], a^{i} \in\left[a^{i}\right], b^{0} \in\left[b^{0}\right], b^{i} \in\left[b^{i}\right]
\end{array}\right.
$$

By projecting this equation on x and y , and by combining the two equations obtained in a $3^{\text {rd }}$ one, we get:

$$
\left\{\begin{array}{c}
\left(a_{x}^{0}-a_{x}^{i}\right) \cdot t+b_{x}^{0}-b_{x}^{i}=0 \\
\left(a_{y}^{0}-a_{y}^{i}\right) \cdot t+b_{y}^{0}-b_{y}^{i}=0 \\
\left(a_{y}^{0}-a_{y}^{i}\right)\left(b_{x}^{0}-b_{x}^{i}\right)-\left(a_{x}^{0}-a_{x}^{i}\right)\left(b_{y}^{0}-b_{y}^{i}\right)=0 \\
t \in\left[0, t_{\max }\right], i \in[1, \ldots m] \\
a_{x}^{0} \in\left[a_{x}^{0}\right], a_{x}^{i} \in\left[a_{x}^{i}\right], b_{x}^{0} \in\left[b_{x}^{0}\right], b_{x}^{i} \in\left[b_{x}^{i}\right] \\
a_{y}^{0} \in\left[a_{y}^{0}\right], a_{y}^{i} \in\left[a_{y}^{i}\right], b_{y}^{0} \in\left[b_{y}^{0}\right], b_{y}^{i} \in\left[b_{y}^{i}\right]
\end{array}\right.
$$

Then, if we express this system using intervals, we get that if:

$$
\exists i \in\{1, \ldots m\},\left\{\begin{array}{c}
0 \in\left(\left[a_{x}^{0}\right]-\left[a_{x}^{i}\right]\right) *\left[0, t_{\max }\right]+\left[b_{x}^{0}\right]-\left[b_{x}^{i}\right] \text { and } \\
0 \in\left(\left[a_{y}^{0}\right]-\left[a_{y}^{i}\right]\right) *\left[0, t_{\max }\right]+\left[b_{y}^{0}\right]-\left[b_{y}^{i}\right] \text { and } \\
0 \in\left(\left[a_{y}^{0}\right]-\left[a_{y}^{i}\right]\right) *\left(\left[b_{x}^{0}\right]-\left[b_{x}^{i}\right]\right)-\left(\left[a_{x}^{0}\right]-\left[a_{x}^{i}\right]\right) *\left(\left[b_{y}^{0}\right]-\left[b_{y}^{i}\right]\right)
\end{array}\right.
$$

Then our boat will have a collision in the time interval $\left[0, t_{\max }\right.$ ].
In order to check for collision between the two waypoints, we should take $t_{\max }$ so that our boat reach the $2^{\text {nd }}$ waypoint at $t_{\max }$. We can estimate $t_{\max }$, as we assume we know the progression speed of the boat on the line between the two waypoints.

In case we have several waypoints, we can repeat this collision detection for each segment of the trajectory. We assume that all the waypoints are stored in a list, and that the last waypoint in the list is a checkpoint, that is we will have to reach it somehow and we cannot modify it. It will be important for the path replanning step. In case we have several waypoints, we must update the variables which represent the boat and obstacles initial positions, so that they correspond to the value they should have at the beginning of this new trajectory segment. We can calculate them because we consider that the speed is constant, and that each trajectory is a line.

In case we detect a collision, we start the path replanning step, which will be described later in the chapter $D$.

```
( int i = 1; i< waypoints. size(); i++){
updateBoatAndObstaclesData(waypoints, boatState, boatSpeed, timeInterval, boatHead, obstacles, i);
bool collisionDetected = 0;
int j = 0;
/check for each obstacles if there is a collision, if yes, call the pathReplanning method
while (j < obstacles.size() and collisionDetected == 0){
    if (collisionCondition(boatSpeed[i-1], boatState[0], boatState[1], boatHead, obstacles[j][0], obstacles[j][1],
        obstacles[j][2], obstacles[j][3], timeInterval*timeSafetyMargin))
    {
        cout << "collision detected in the " << i <<" segment, with obstacles " << j << endl;
        pathReplanning(boatHead, boatSpeed[i-1], boatState, timeInterval*timeSafetyMargin, obstacles, borderList,
            useProjectionForBuildingObstacleSeparator);
        waypointManagement(boatHead, boatSpeed, boatState, timeInterval.ub(), waypoints, i);
        collisionDetected = 1;
    }
```

Figure 3: processing of the waypoint list and collision detection

## C) Detection of the crossing of the sailing zone borders

We decided to model the sailing zone as a list of polygons of which the boat shouldn't cross the sides. Each polygon is modelled by the list of the coordinates of its vertices. We assume that those coordinates are expressed in a fixed cartesian frame. The crossing conditions are inspired by the example presented in the article Path planning using intervals and graphs [1].

We are in the same case than for the collision detection, that is the trajectory of our boat is a segment between two waypoints. We can handle a case with several waypoints by proceeding in the same way than in the collision detection. We consider that we have two boxes $A$ and $B$. The box $A$ is the initial position of the boat, and the box $B$ is the final position of the boat. We denote by $V_{i}$ the $i^{\text {th }}$ vertex of a polygon. We assume that $[A, B]$ denote the set of all the segments with an endpoint in $A$ and an endpoint in $B$. For each polygon of the sailing zone, we need to check if:

$$
\forall i, \forall S \in[A, B], \quad\left[V_{i}, V_{i+1}\right] \cap S=\varnothing
$$

If it's true for all sides of all polygons, the boat never crosses the sailing zone border.
But it's difficult to implement this relation. To avoid this issue, we will consider the contrapositive and use these equivalent relations:

$$
\exists i, \quad\left\{\begin{array}{cc}
\exists S \in[A, B], & \left(V_{i}, V_{i+1}\right) \cap S \neq \emptyset \text { and } \\
\exists D \in(A, B), & {\left[V_{i}, V_{i+1}\right] \cap D \neq \emptyset \text { and }} \\
{\left[V_{i} \cup V_{i+1}\right] \cap[A \cup B] \neq \emptyset}
\end{array}\right.
$$

We assume that $(A, B)$ denote the set of all the lines supported by a point in A and a point in $\mathrm{B},\left(V_{i}, V_{i+1}\right)$ the line supported by $V_{i}$ and $V_{i+1},\left[V_{i}, V_{i+1}\right]$ the segment linking $V_{i}$ and $V_{i+1}$, and [ $A \cup B$ ] and $\left[V_{i} \cup V_{i+1}\right.$ ] the smallest box including A and B or $V_{i}$ and $V_{i+1}$ respectively.

The last condition is important for a special case, where all the points are aligned. In that case, the two first conditions will be true event if the trajectory doesn't intersect the polygon side.

We can't implement directly the two first conditions, so we will use the following equivalence in the code instead:

$$
\exists i, \quad\left\{\begin{array}{c}
\operatorname{det}\left(\overrightarrow{\left(\overrightarrow{V_{l} B}, \overrightarrow{V_{l} A}\right) * \operatorname{det}\left(\overrightarrow{V_{l+1} B}, \overrightarrow{V_{l+1} A}\right) \leq 0 \text { and }}\right. \\
\operatorname{det}\left(\overrightarrow{V_{l+1} V_{l}}, \overrightarrow{V_{l} A}\right) * \operatorname{det}\left(\overrightarrow{V_{l+1} \vec{V}_{l}}, \overrightarrow{V_{l} B}\right) \leq 0 \text { and } \\
{\left[V_{i} \cup V_{i+1}\right] \cap[A \cup B] \neq \varnothing}
\end{array}\right.
$$

If it's true for one of the polygon of the sailing zone, then the boat trajectory is crossing this polygon.

As $A$ and $B$ are boxes, the coordinates of the vectors are some intervals, and the result of the determinants are intervals too. So, in fact we manipulate a set of vectors and a set of determinants. The mathematical notations may be inappropriate, but in practice it gives some good results. We make this test just after the collision detection. If we detect that the boat will cross the sailing zone border, we call the path replanning method.


Figure 4: detection of the crossing of the sailing zone borders
D) Path replanning

The second part, that is the modification of the list of waypoints, is a path planning problem with dynamic obstacles. Some path planning problems, with static obstacles, can be solved easily and efficiently with interval analysis method and graphs [1]. But the dynamic obstacles make this approach much more difficult, because it also depends on time, and it's harder to define the configuration space.

The issue of dynamic path planning is not trivial, and the solution often chosen is to consider that the obstacles are static, and make the calculation at a high frequency in order to take into account the fact that the obstacles are moving. This solution is simple on the theoretical level, and easy to implement, but it is not efficient, and it can be unsafe in some situations. Some solution for the dynamic path planning had been developed recently. But often, the papers describing those solutions aren't freely available. And when they are, those solutions are complex on the theoretical level, and we were not sure to have the skill and the time required to adapt it to our situation and to implement it.

As the detection step uses some interval analysis, we choose to use the same approach for the path replanning step. The general idea is to compute the set of all the speed vectors that allow our boat to avoid a collision and that allow it to stay in the sailing zone.

For that, we use the separator algebra. We build a separator corresponding to the collision condition of each obstacles, and to the crossing condition of each side of each polygon of the sailing zone. Then, we compute the union of all these separators. Finally, we can compute with a paving method the set of all the feasible speed vectors, that is the set which avoid collisions and which avoid crossing the sailing zone borders. Then, we choose the speed vector which is the closest to our initial speed vector.

We keep this speed vector during the same amount of time as it would take to our boat to reach the waypoint with its previous speed vector. This assumption allows us to compute a new waypoint, and to replace the former one by this one. For the last waypoint of the list, we assume it is a checkpoint, so we shouldn't modify it. In that case, we add a new waypoint before the last one, instead of directly modifying it. However, if we are close enough to the last waypoint, that is if the last waypoint is inside the final position uncertainty box of the boat, we replace the last waypoint. It allows us to avoid some bugs link to the size of the uncertainty box. If we compute a very long path, at the end the uncertainty box is so huge that it could be impossible for the center of the box to reach the waypoint without a part of the box collide with something or cross a border.

## Collision separators

To build the collision separators, we use the same conditions than in the detection part. We build one intermediary separator for each condition, then compute the intersection.

Then this separator, which correspond to one obstacle, goes in a list. This list aims to store all the separators in order to make the union when we have built them all.

However, there is an issue with these conditions. Indeed, the inside of the set they produce is empty. It produces a lot of boxes during the paving process and it can lead to a very long execution time. In order to bypass this issue, we needed to rephrase these constraints. We also needed to use the separators associated to the projection of the set defined by the separator built with these constraints. Indeed, the dimension of the separators is too high to allow them to have a non-empty inside. The projection allows us to decrease this dimension. We rephrased the constraints as follow:

$$
\exists i \in\{1, \ldots m\},\left\{\begin{array}{c}
{\left[b_{x}^{i}\right]-\left[b_{x}^{0}\right] \in\left(\left[a_{x}^{0}\right]-\left[a_{x}^{i}\right]\right) *\left[0, t_{\max }\right] \text { and }} \\
{\left[b_{y}^{i}\right]-\left[b_{y}^{0}\right] \in\left(\left[a_{y}^{0}\right]-\left[a_{y}^{i}\right]\right) *\left[0, t_{\max }\right] \text { and }} \\
\left(\left[b_{y}^{i}\right]-\left[b_{y}^{0}\right]\right) \in \frac{\left(\left[a_{y}^{0}\right]-\left[a_{y}^{i}\right]\right)}{\left(\left[a_{x}^{0}\right]-\left[a_{x}^{i}\right]\right) *\left(\left[b_{x}^{i}\right]-\left[b_{x}^{0}\right]\right)}
\end{array}\right.
$$

After building the separators linked with these constraints, we build the separators of their projection. Then, we make the intersection of these three separators of projection. Finally, we can add them in the separator list, as previously.

But the fact is that the projection operation is quite slow, so the execution time doesn't improve as expected. Depending on the resolution of the paving and the number of obstacles, in some cases it's faster to use projection, and in other cases, it's faster without using it. So, we decided to let the two version in the code, and to allow the user to choose which version he wants to use. That way, it will be easy to make more tests, and eventually decide later which version we should use.

## Example

On figure $5 \& 6$, you can see two paving made in the same conditions. There are three obstacles, and the sailing zone is defined by a square around the evolution area of our boat.

The red boxes are inside the set of speed vectors that lead to a collision or to cross a border of the sailing zone. The Yellow boxes are on the border of this set. Finally, the light blue boxes are outside this set, so they represent the set of the feasible speed vectors. The green box is the speed vector chosen by the algorithm.


Figure 5: result of the paving process without use of projection
The red boxes we can see on the edge of the image correspond to the speeds that lead to cross the sailing zone border. At the center, the two triangle shaped yellow set (surrounded in black here for more readability) correspond to the speeds that lead to a collision with one of the obstacles.


Figure 6: result of the paving process with use of projection
Here we can see that the use of projection doesn't modify the feasible speed set, and that there are less yellow boxes in the set corresponding to the collision with an obstacle.


Figure 7: trajectories of the boats
On figure 7 we can see the trajectories of the different boats. In red is displayed the sailing zone borders. In yellow we can see the initially planned path of our boat. In black, there are the paths of the other boats that we need to avoid. Finally, in blue is displayed the path
of our boat after the collision avoidance. We are not able to see it here, but at run time we display the different uncertainty boxes moving on their paths.

## Sailing zone border crossing separators

For the sailing zone border crossing separators, we were not able to use the same conditions than for the detection step. The two first conditions are usable, but the $3^{\text {rd }}$ one cause some troubles. As a reminder, the $3^{\text {rd }}$ condition is the following:

$$
\exists i, \quad\left[V_{i} \cup V_{i+1}\right] \cap[A \cup B] \neq \varnothing
$$

See chapter C for more information. This condition is true if the two boxes are overlapping.
Indeed, with the library we are using for interval computation, Ibex, we can't use a union or an intersection operator in the expression of a constraint defining a separator. This is not allowed be the library, and we couldn't find an easy way to bypass this issue. We were forced to rephrase this constraint. The issue was to find a way to detect when two boxes are overlapping, without using any union nor intersection operator. We will show the main idea with a figure, because it's easy to understand visually, but we won't write the underlying explicit mathematical expression in term upper and lower bounds of boxes, as it is in the code. Indeed, it is very long and quite impenetrable, so it may be confusing. On figure 8, we show an example of two overlapping boxes. Let say the blue one is [ $V_{i} \cup V_{i+1}$ ], and the red one is $[A \cup B]$ for example.


Figure 8: example of two overlapping boxes
The two boxes are overlapping if:

$$
\left\{\begin{array}{l}
\max \left(d_{1}, D_{1}\right)-\max \left(L_{1}, l_{2}\right) \leq 0 \text { and } \\
\max \left(d_{2}, D_{2}\right)-\max \left(l_{1}, L_{2}\right) \leq 0
\end{array}\right.
$$

These conditions are used to build the separator, but the real expression is much more complex. Indeed, all the variables depend on the relative position of the upper bounds or lower bounds of the boxes A or B, and on the relative position of the two vertices.

Then, the separator corresponding to one side of a polygon is added in the separator list, in the same way as for the collision condition separators.
E) Perspectives of improvement

The various tests we made allowed us to find some potential issues with the actual algorithm that may cause some troubles when we will use this algorithm in real conditions.

First, for now the algorithm assume that we can control the speed of our boat when it chooses a new speed vector. Theoretically, it depends on the projection on the boat course of the force applied on the wingsail, so on the wind speed and direction, on the heading of the boat, and on the angle of attack of the wingsail. In theory, the angle of attack is controlled thanks to the angle of the tailwing. So, it seemed to be possible to control the speed of the boat. In practice, the on-sea tests we made show that it is very difficult to control the speed of the boat because it depends on too many parameters. A simple way to bypass this issue could be to choose a new speed vector which has the same norm as the actual speed vector. Indeed, it will require more tests in real conditions, but it might be possible to assume that our speed can remain constant, as soon as we process the collision avoidance on a reasonable time interval, for example about few tens of minutes.

There may be some issue when we choose a direction where the boat will have to tack. Indeed, the progression speed of the boat on the segment will be slightly slower than when it doesn't need to tack. This case need to be checked, because fixing this issue will require to complexify the algorithm. Indeed, if we can't choose the same speed norm than the actual speed for some directions, we will need to know somehow which are these directions, and which speed norm we can have in these directions. For that, the choose of the speed vector will have to take into account the wind direction.

Another issue is linked to the uncertainty box corresponding to the position of our boat. Actually, it doesn't match the real uncertainty we have about the position of the boat. Indeed, for each trajectory segment, we know that the boat will follow a line between the two waypoints, and will stay in a given corridor around this line. So, the box shape doesn't fit with this knowledge. It causes some trouble when we process the collision avoidance on a
long amount of time. Indeed, the uncertainty box become huge, but it's not consistent with the reality, because we know that our boat will stay in a corridor around the line. The only thing we don't know is where exactly our boat is in this corridor, mainly because of the speed uncertainty. It's an issue, because if the uncertainty box is too large, it will detect some collision whereas in reality there is no risk. In some cases when the box is very large, the set of feasible speed vector can even become empty, and the algorithm will stop without finding any path. So, we must choose carefully the amount of time on which we want to process the collision avoidance, and make sure it isn't too long. To avoid this issue, a simple solution would be to consider that the speed isn't an interval, but a fixed value. In that case the uncertainty box corresponding to the position of our boat would have a fixed size. Of course, we would lose partly the interest of the interval analysis, that allow us to take into account uncertainties. It will require some tests in real conditions, because if in reality the speed of our boat doesn't vary a lot between two waypoints, the use of an interval to model it isn't necessary.

## II) Preparation for the WRSC

The end of my internship matched with the participation of the ASPire project team to the WRSC (World Robotic Sailing Championship). In order to prepare this competition, we needed to have a working simulator, to test the control system before the on-sea tests.

The simulator was reworked this summer and it wasn't working properly yet. I was assigned with fixing the remaining issue with this simulator. It is organised on a client-server architecture. The server is the control system, and the client is the simulator and the display. A lot of issues was linked to some conversion errors between two conventions. Indeed, the control system use a north east down convention, and the simulator an east north up convention. Moreover, in the simulator the angles are in the range [-180, 180], whereas in the control system they are in the range [ 0,360 ]. So, my main task was to check for each variable that was sent from one side to another if it was in the right convention on both sides.

But when all the convention issues were fixed, the simulated boat still had some illogical behaviours. I was assigned with investigating these remaining issues. I finally found out that there was something wrong with the calculation of the force on the wingsail. I couldn't find exactly what was wrong, but the value and the sign of the force seems
incoherent with the heading of the boat and the wind direction. So, some further investigation should be made on this part of the code, and maybe the physics equations should be checked too.

I was also charged to develop a station keeping controller for the station keeping task of the WRSC. The idea was to take advantage of the ability of our boat to go backward thanks to the wingsail, to adjust our position as close as possible to the waypoint. Indeed, we could stop on the waypoint, and if we start to drift forward, we could go back to the waypoint. Likewise, if we drift backward, we could reach back to it. However, we haven't any means to know the direction of our speed at low speed. Indeed, we use the GPS speed, but at low speed its direction is irrelevant. Without this knowledge, we can't control correctly the rudder and the tailwing, because obviously it won't be the same control if we are going backward as if we are going forward. Finally, we used the controller developed last year, which try to go back to the waypoint when the boat passed it, using the line following controller. In practice, the boat keeps turning around the waypoint.

Finally, during the WRSC, I was involved in the on-sea tests, and in the trials preparation.

## Economic analysis

The ASPire project is financed for three years by the European Union, through the European Regional Development Fund (50\%) and the Åland government, through the Åland's program for structural funds "Entrepreneurship and competency" (50\%). They contributed up to around $350000 €$.

Around $80 \%$ of this budget is dedicated to personnel costs. This include the project leader wages for 3 years, 11.5 months of wages for a software developer, and 6.5 months of wages for an electronic and mechanical engineer. A part of this budget is also dedicated to some internships. Nevertheless, my internship wasn't on this budget. Indeed, as it was an Erasmus internship, I was paid by Erasmus.

About 9\%, that is around $30000 €$, of the total budget is allocated to the purchase of services. For example, the seaworthiness calculation, some part of the construction of the
boat, the dimensioning and the construction of the wingsail, weren't realised by the project team, but were subcontracted.

Approximately the same part of the budget is dedicated to the purchase of materials (electronic and mechanical) and tools, and the renting of machines and devices. However, this budget is already almost entirely used whereas the project should still last for one year and a half. Indeed, a great part of the construction wasn't subcontracted, and some of the building work that needed to be done was hardly predictable.

Finally, the remaining budget is allocated to travel expenses, such as travel allowances, tickets, food \& lodging, car costs... For example, this budget was used for attendance at HIPER 2017 and OCEANS 2017 conferences. Our trip to Norway for the WRSC and the expenses made during our stay wasn't on this budget, but on the budget of the Finnish Maritime foundation.

## Conclusion

First, this internship was for me the opportunity to gain a lot of technical skill, in some various fields. First, I was able to increase my knowledge about interval analysis. Then, As the whole project is in $\mathrm{C}++$, it allowed me to practice this language a lot, and more particularly the interval library, Ibex. Moreover, I'm now familiar with the git workflow, as the project is using it. It was also the occasion to work in an international team, and to practice the scrum, a project management method. So, on the technical field this internship was very rewarding.

It also allowed me to have an insight of the research world, as I was integrated in a research team in the Åland University of Applied Science, and as my internship subject was research oriented. This experience confirmed my wish to work in this sector.

So, this internship was very interesting on many aspects, and I hope my work will be useful for the future developments of the project.

## Acknowledgment

I thank Anna Friebe, my project manager, for offering me the opportunity to make this internship, and for being as available as she could to answer my questions, on technical or practical concerns.

I also thank Luc Jaulin, my school internship supervisor and one of my teachers, for giving me Anna's contact for this internship, and for his advices and explanations about interval analysis. Finally, I thank Maël Le Gallic, my internship supervisor, for making sure that I had all the information to understand well the goal of my internship subject and the actual status of the project in general. I also thank him for revising this report.

## References

[1] Luc Jaulin, Path planning using intervals and graphs, Reliable Computing Journal, 2001, 7(1), pp.1-15
[2] Luc Jaulin, Fabrice Le Bars, A Simple Controller for Line Following of Sailboats, Robotic Sailing 2012, p 117-129

Appendix

RAPPORT D'EVALUATION ASSESSMENT REPORT

## ENSTA

Bretagne

Marci de returner ce rapport en fin du stage ad :
Please return this report at the end of the internship to :
ENSTA Bretagne - Bureau des stages - 2 rue François Verny - 29806 BREST codex 9 - FRANCE
面 00.33 (0) 2.98.34.87.70

- Fax 00.33 (0) 2.98.38.87.90- stages(o)ensta-bretagne.fr

I - ORGANISME / HOST ORGANISATION
nom/Name Alland University of Applied Sciences
Adresse /Address Neplunigtan 17 PB 1010, $\theta x-22111$
Mariehamn, Bland, Finland
Tél / Phone (including country and area code) +358 18 537000
Fax / Fax (including country and area code) $\qquad$
Nom du superviseur / Name of placement supervisor


Adresse e-mail / E-mail address maul. leg alice. oo ha. ax
Nom du stagiaire accueilli / Name of trainee ROLINAT Elément

## II - EVALUATION / ASSESSMENT

Veuillez attribuer une note, en encerclant la lettre appropriée, pour chacune does caractéristiques suivantes. Cette note devra se situer entre $\mathbf{A}$ (trees bien) et $\mathbf{F}$ (très faible)
Please attribute a mark from $A$ (very good) to $F$ (very weak).

## MISSION / TASK

* La mission de départ a-t-elle été remplie ?
$A(B C D E F$
Was the initial contract carried out to your satisfaction?
* Manquait-il au stagiaire de connaissances ?

Was the trainee lacking skills?
Si oui, lesquelles ?/ If so, which skills? $\qquad$

## ESPRIT D'EQUIPE / TEAM SPIRIT

* Le stagiaire s'est-il been intégré dan l'organisme d'accueil (disponible, sérieûx, s'est adapté au travail en groupe) / Did the trainee easily integrate the host organisation? (flexible, conscientious, adapted to team work)


Souhaitez-vous nous fare part d'observations on suggestions?/ If you wish to comment or make a suggestion, please do so here $\qquad$

COMPORTEMENT AU TRAVAIL / BEHAVIOUR TOWARDS WORK
Le comportement du stagiaire était-il conforme à vos attentes (Ponctuel, ordonné, respectueux, soucieux de participer et d'acquérir de nouvelles connaissances) ?
Did the trainee live up to expectations? (Punctual, methodical, responsive to management instructions, attentive to quality, concerned with acquiring new skills)?
(A)BCDEF

Souhaitez-vous nous faire part d'observations on suggestions? / If you wish to comment or make a suggestion, please do so here $\qquad$

INITIATIVE - AUTONOMIE / INITIATIVE - AUTONOMY
Le stagiaire s'est -il rapidement adapté à de nouvelles situations? (Proposition de solutions aux problèmes rencontrés, autonomic dan le travail, etc.)

Did the trainee adapt well to new situations?
$A(B) C D E F$ (eg. suggested solutions to problems encountered, demonstrated autonomy in his/her job, etc.)

Souhaitez-vous nous fair part d'observations on suggestions? / If you wish to comment or make a suggestion, please do so here $\qquad$

CULTUREL - COMMUNICATION / CULTURAL - COMMUNICATION
Le stagiaire était-il ouvert, d'une manière générale, à la communication?
$A B C D E F$
Was the trainee open to listening and expressing himself /herself?
Souhaitez-vous nous fare part d'observations on suggestions? / If you wish to comment or make a suggestion, please do so here $\qquad$

OPINION GLOBALE / OVERALL ASSESSMENT

* La valeur technique du stagiaire était :

ABCDEF
Evaluate the technical skills of the trainee:

## III - PARTENARIAT FUTUR / FUTURE PARTNERSHIP

* Etes-vous prêt à accueillir un autre stagiaire l'an prochain ?

Would you be willing to host another trainee next year? $\quad \square$ oui/yes non/no


