



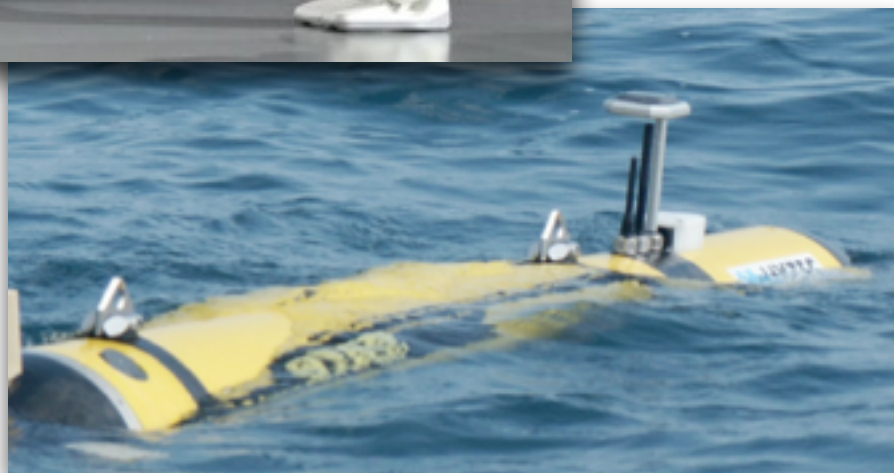
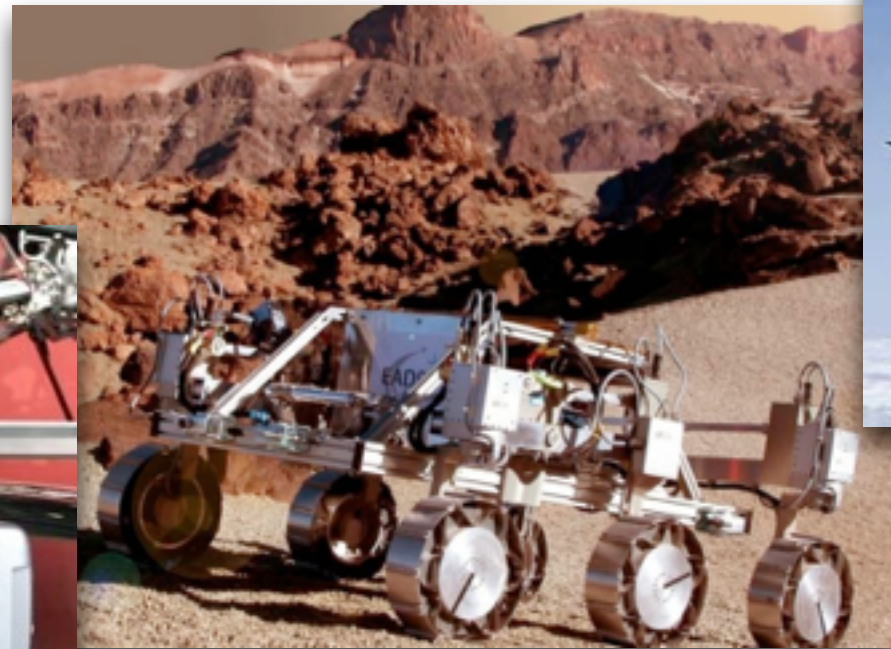
# **Nonlinear and hybrid reachability analysis in presence of uncertainty.**

**Nacim RAMDANI,  
Univ. Orléans, EA 4229 PRISME à Bourges.**

**GT MEA - CPNL, 3 Avril 2014, Paris.**

- Hybrid Reachability : Motivations
  - Analysis of complex dynamical systems
  - Reachability-based methods
  
- Nonlinear reachability
  - Interval Taylor methods
  - Bracketing enclosures
  - Software implementation

# Hybrid Cyber-Physical Systems



- **Interaction discrete  
+ continuous dynamics**
- **Safety-critical  
embedded systems**
- **Networked  
autonomous systems**



# Hybrid Cyber-Physical Systems



## ■ Verification

- Numerical proof
- Falsification via counter-example



# Hybrid Cyber-Physical Systems

## ■ Modelling → **hybrid automaton** (Alur, et al. 1995)

- Non-linear continuous dynamics
- Bounded uncertainty

$$H = (\mathcal{Q}, \mathcal{D}, \mathcal{P}, \Sigma, \mathcal{A}, \text{Inv}, \mathcal{F}),$$

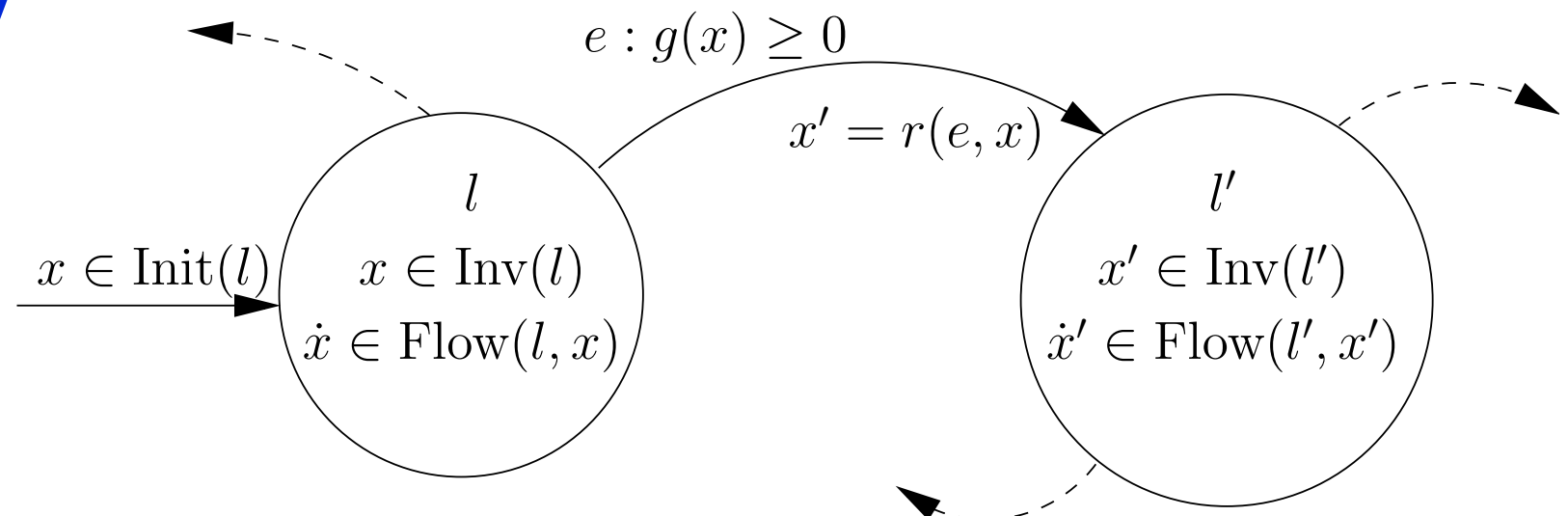
*Continuous dynamics*

$$\begin{aligned} \text{flow}(q) : \quad & \dot{\mathbf{x}}(t) = f_q(\mathbf{x}, \mathbf{p}, t), \\ \text{Inv}(q) : \quad & \nu_q(\mathbf{x}(t), \mathbf{p}, t) < 0, \end{aligned}$$

*Discrete dynamics*

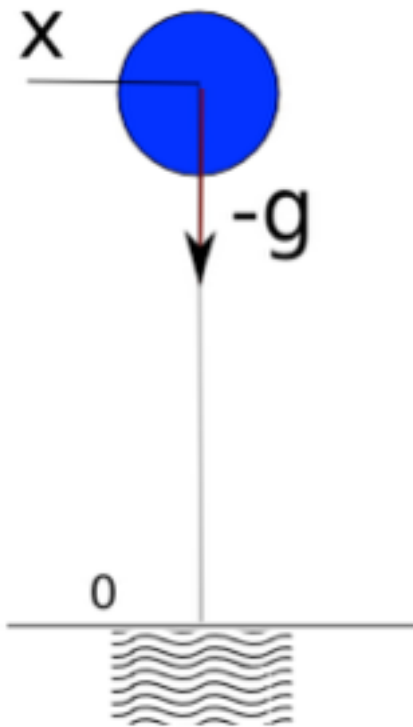
$$\begin{aligned} \mathcal{A} \ni e : \quad & (q \rightarrow q') = (q, \text{guard}, \sigma, \rho, q'), \\ \text{guard}(e) : \quad & \gamma_e(\mathbf{x}(t), \mathbf{p}, t) = 0, \end{aligned}$$

$$t_0 \leq t \leq t_N, \quad \mathbf{x}(t_0) \in \mathbb{X}_0 \subseteq \mathbb{R}^n, \quad \mathbf{p} \in \mathbb{P}$$



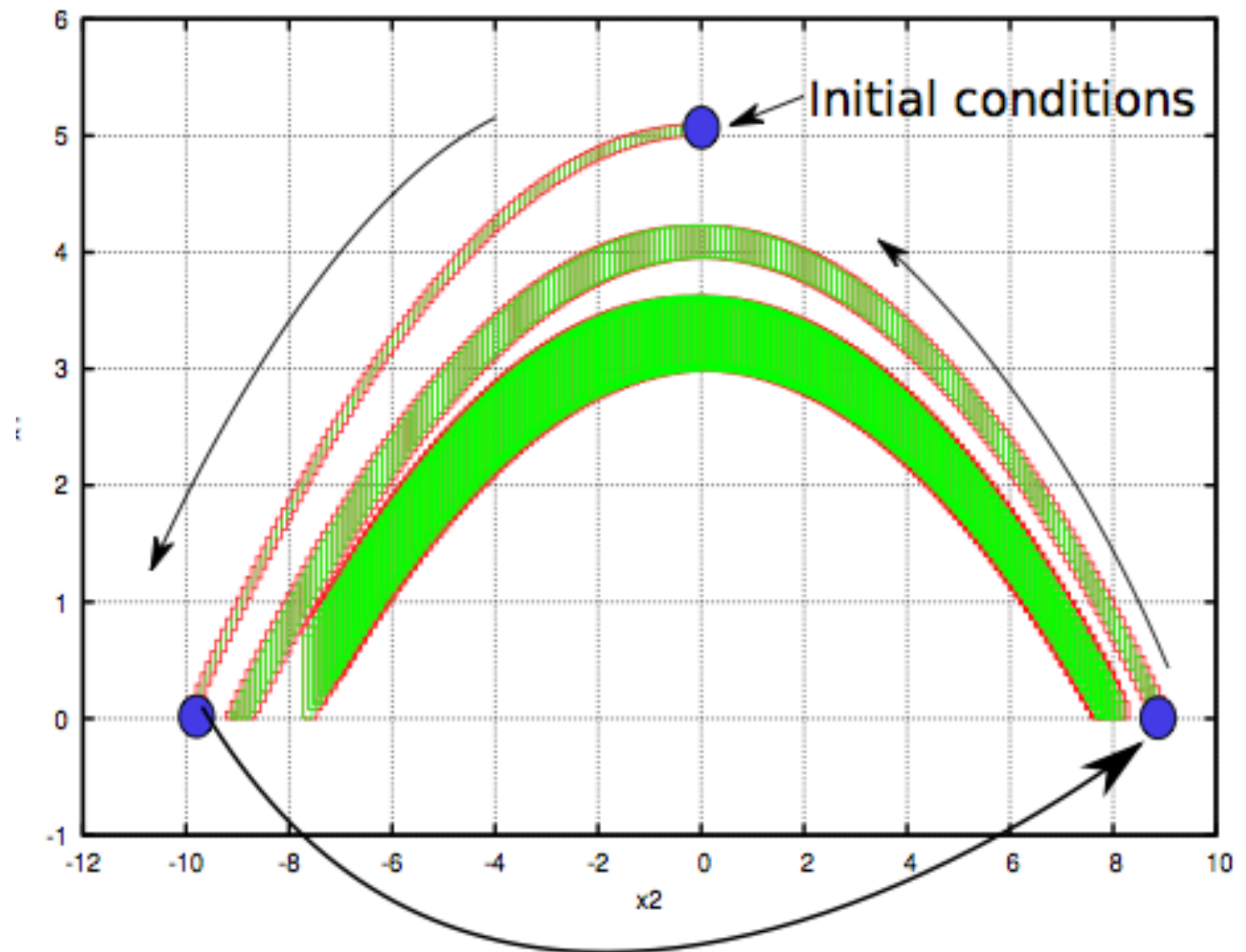
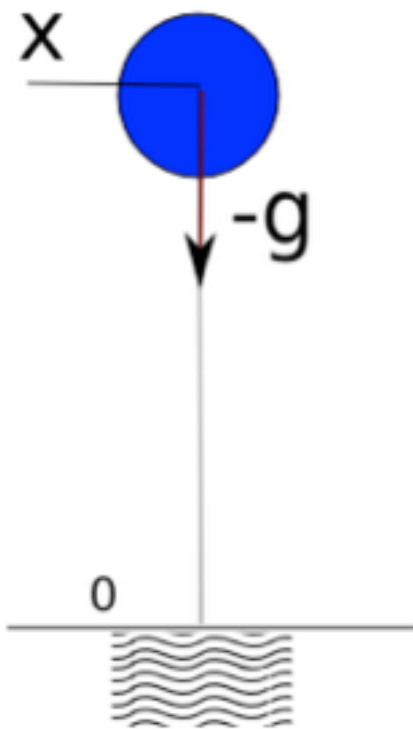
# Hybrid Cyber-Physical Systems

## ■ Example : bouncing ball



# Hybrid Cyber-Physical Systems

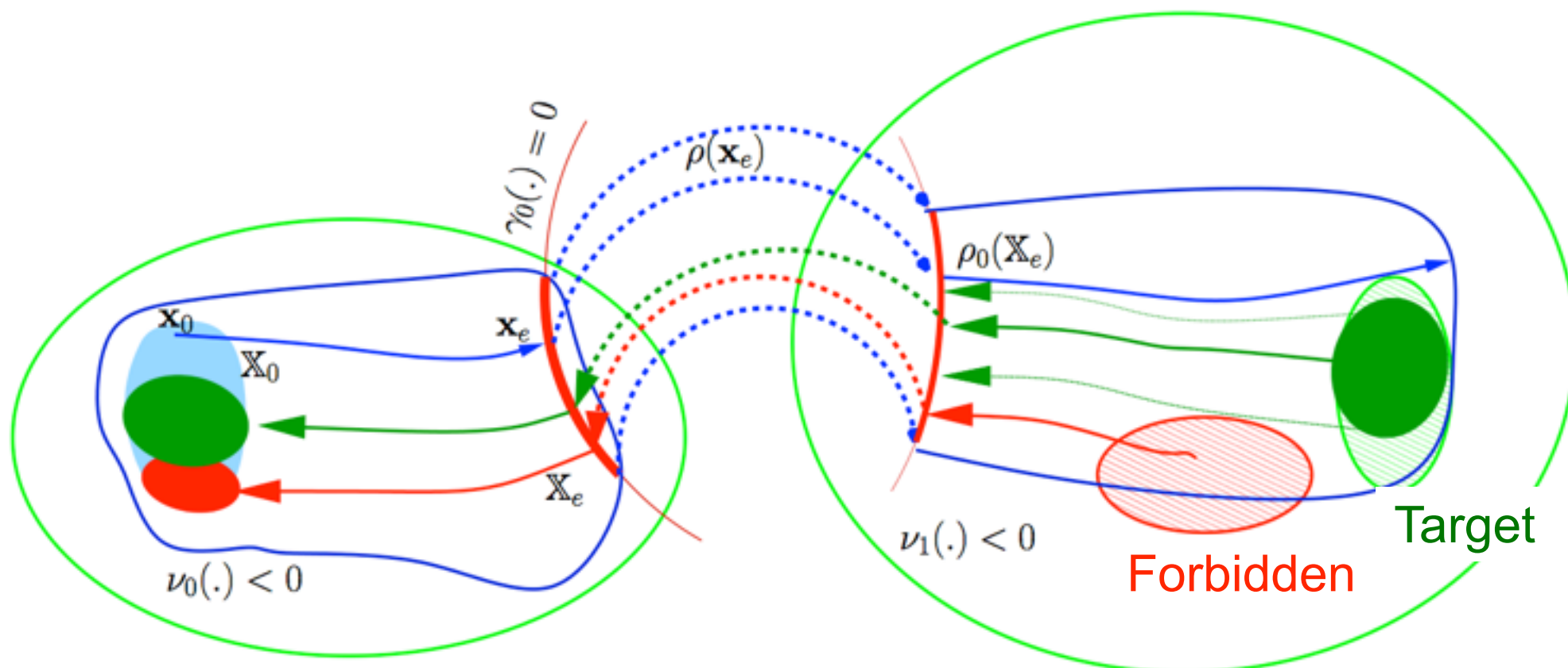
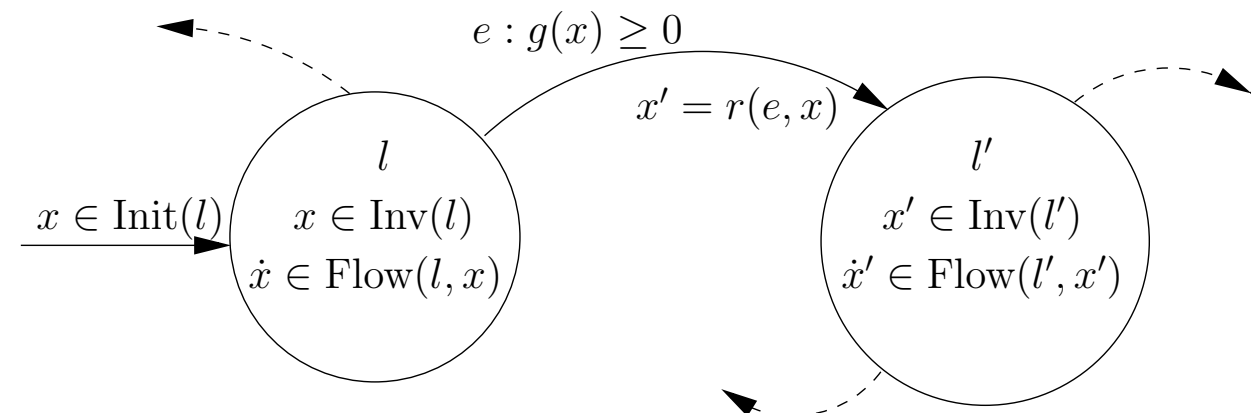
## ■ Example : bouncing ball



# Hybrid Cyber-Physical Systems

## ■ Verification

- *Modelling :*
- *Property specification :*
- Verification algorithm :
  - Hybrid / Continuous reachability





# Hybrid Cyber-Physical Systems

## ■ Continuous reachability

$$\mathbb{R}([t_0, t]; \mathbb{X}_0) = \left\{ \begin{array}{l} \mathbf{x}(\tau), \quad t_0 \leq \tau \leq t \mid \\ \dot{\mathbf{x}}(\tau) = f(\mathbf{x}, \mathbf{p}, \tau) \wedge \mathbf{x}(t_0) \in \mathbb{X}_0 \wedge \mathbf{p} \in \mathbb{P} \end{array} \right\}$$

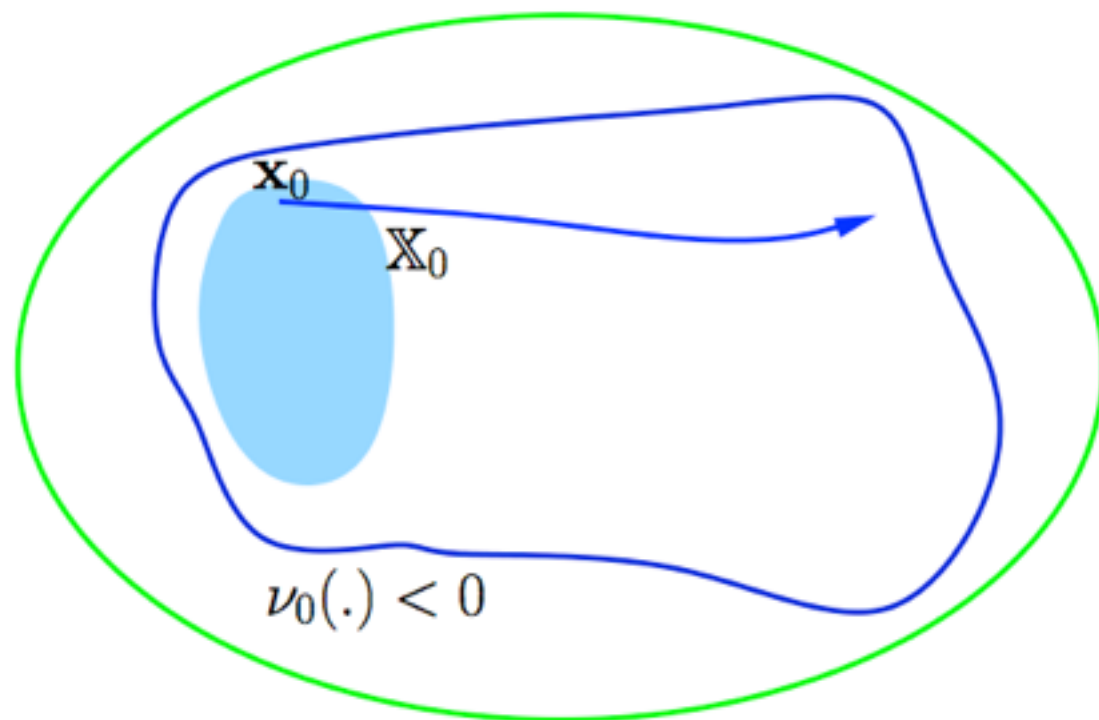
- Set integration
  - Interval Taylor methods
  - Bracketing enclosures

# Hybrid Cyber-Physical Systems

## Continuous reachability

$$\mathbb{R}([t_0, t]; \mathbb{X}_0) = \left\{ \mathbf{x}(\tau), t_0 \leq \tau \leq t \mid \dot{\mathbf{x}}(\tau) = f(\mathbf{x}, \mathbf{p}, \tau) \wedge \mathbf{x}(t_0) \in \mathbb{X}_0 \wedge \mathbf{p} \in \mathbb{P} \right\}$$

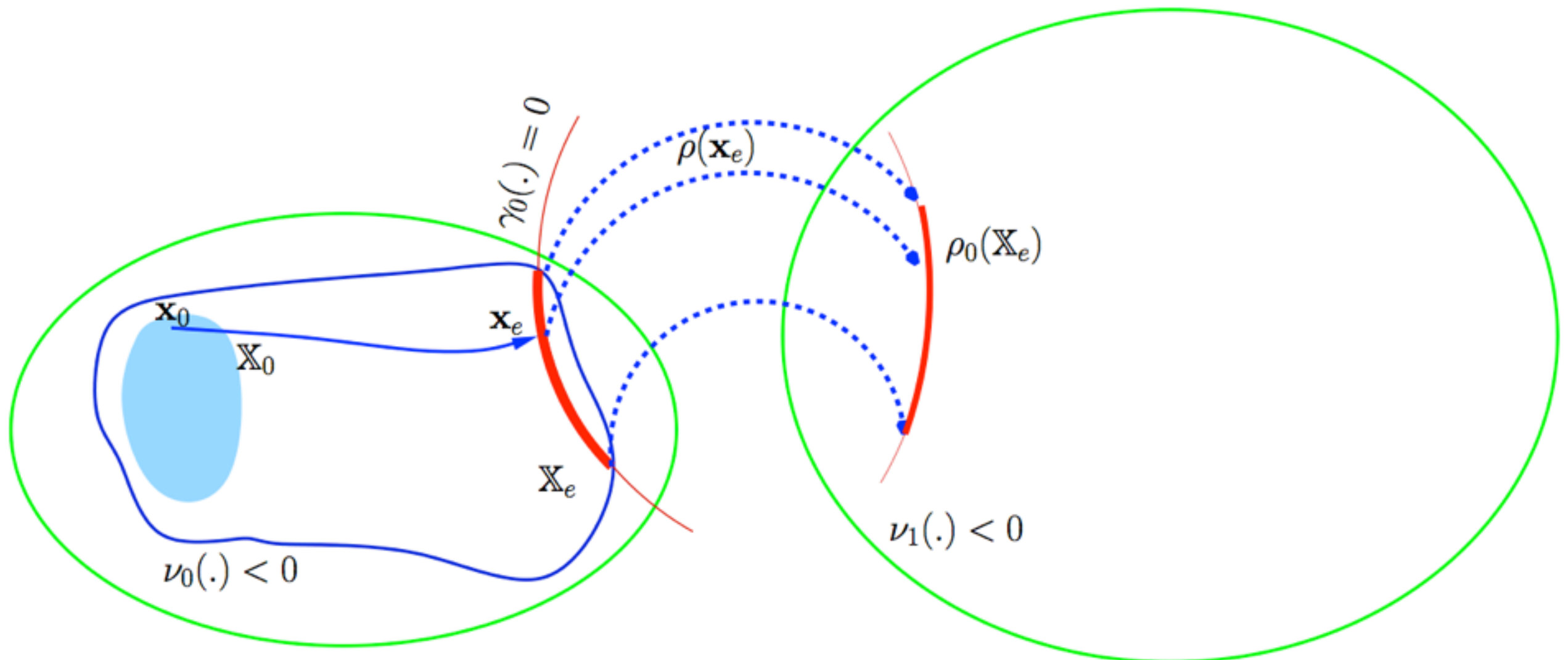
- Set integration
  - Interval Taylor methods
  - Bracketing enclosures



# Hybrid Cyber-Physical Systems

## ■ Hybrid reachability

- Continuous reachability
- Guard conditions, jumps & resets



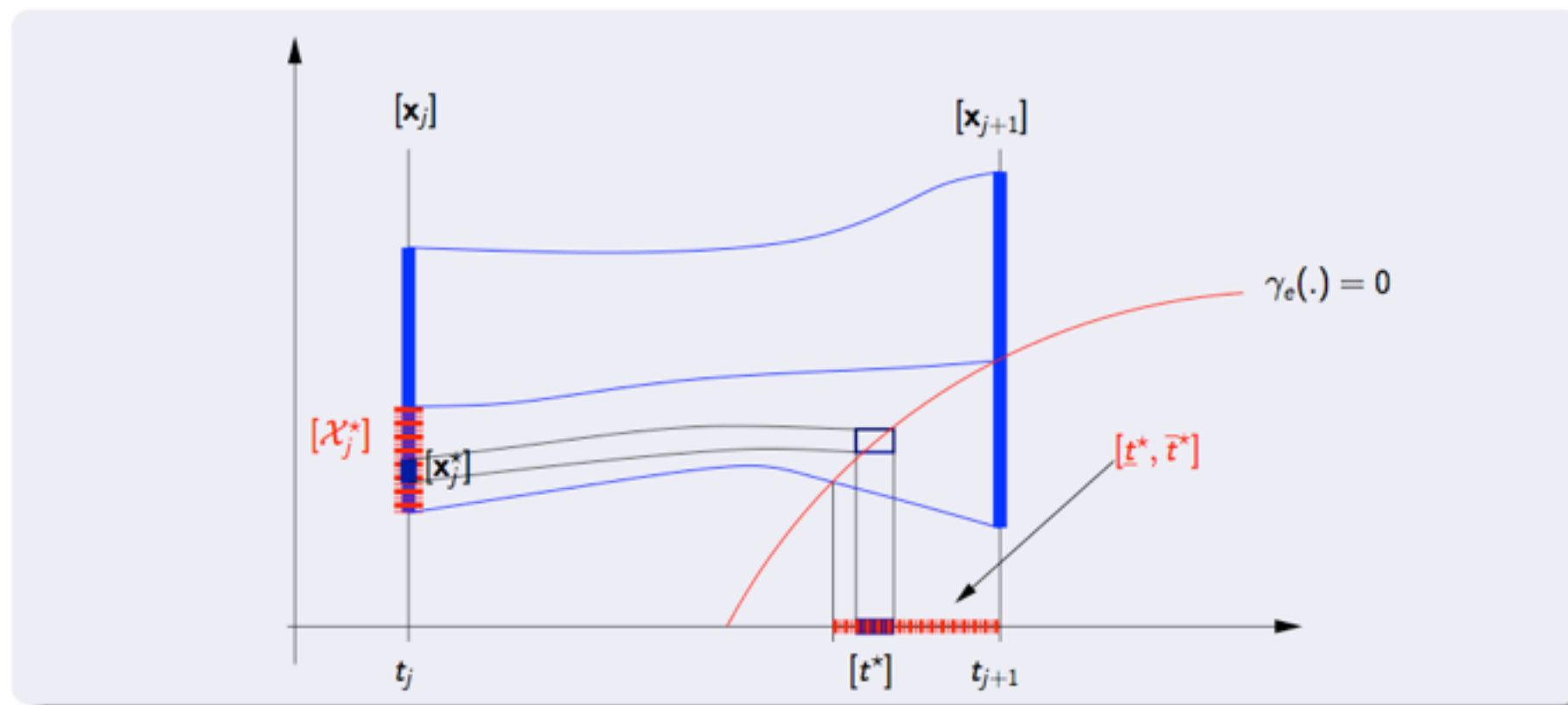
# Hybrid Reachability Analysis

## ■ Guaranteed event detection & localization

### ● An interval constraint propagation approach

● (Ramdani, et al., Nonlinear Analysis Hybrid Systems 2011)

Time grid  $\rightarrow t_0 < t_1 < t_2 < \dots < t_N$



Compute  $[t^*, \bar{t}^*] \times [x_j^*]$



# Hybrid Reachability Analysis

## ■ Guaranteed event detection & localization

### ● An interval constraint propagation approach

● (Ramdani, et al., Nonlinear Analysis Hybrid Systems 2011)

Time grid  $\rightarrow t_0 < t_1 < t_2 < \dots < t_N$

- $[\mathbf{x}](t) = \text{Interval Taylor Series (ITS)}(t, [\mathbf{x}_j], [\tilde{\mathbf{x}}_j])$
- $\gamma([\mathbf{x}](t)) = 0$

$$\Rightarrow \gamma \circ \text{ITS}(t, \mathbf{x}_j, [\tilde{\mathbf{x}}_j]) \rightarrow \psi(t, \mathbf{x}_j)$$

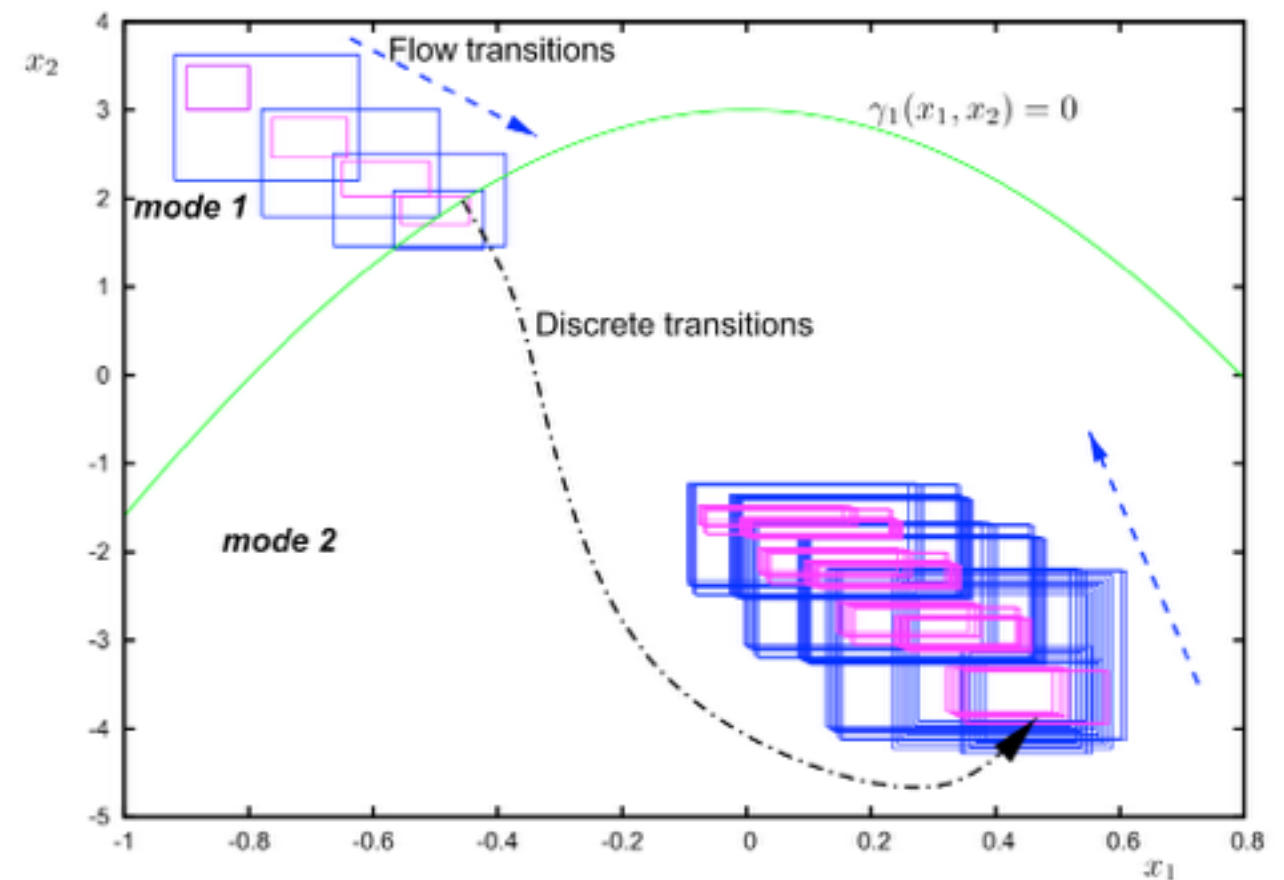
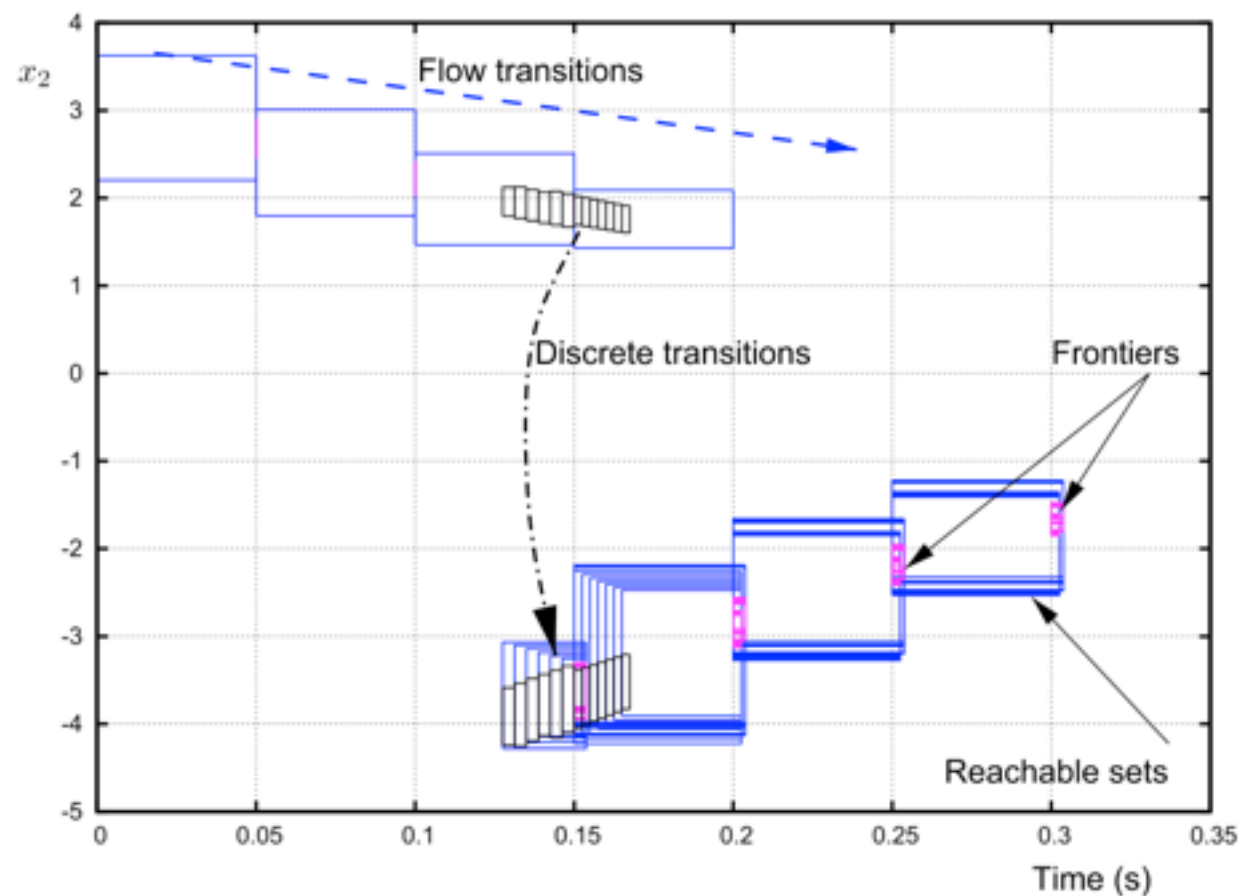
Solve CSP  $([t_j, t_{j+1}] \times [\mathbf{x}_j], \psi(.,.) \ni 0)$

# Hybrid Reachability Analysis

## ■ Guaranteed event detection & localization

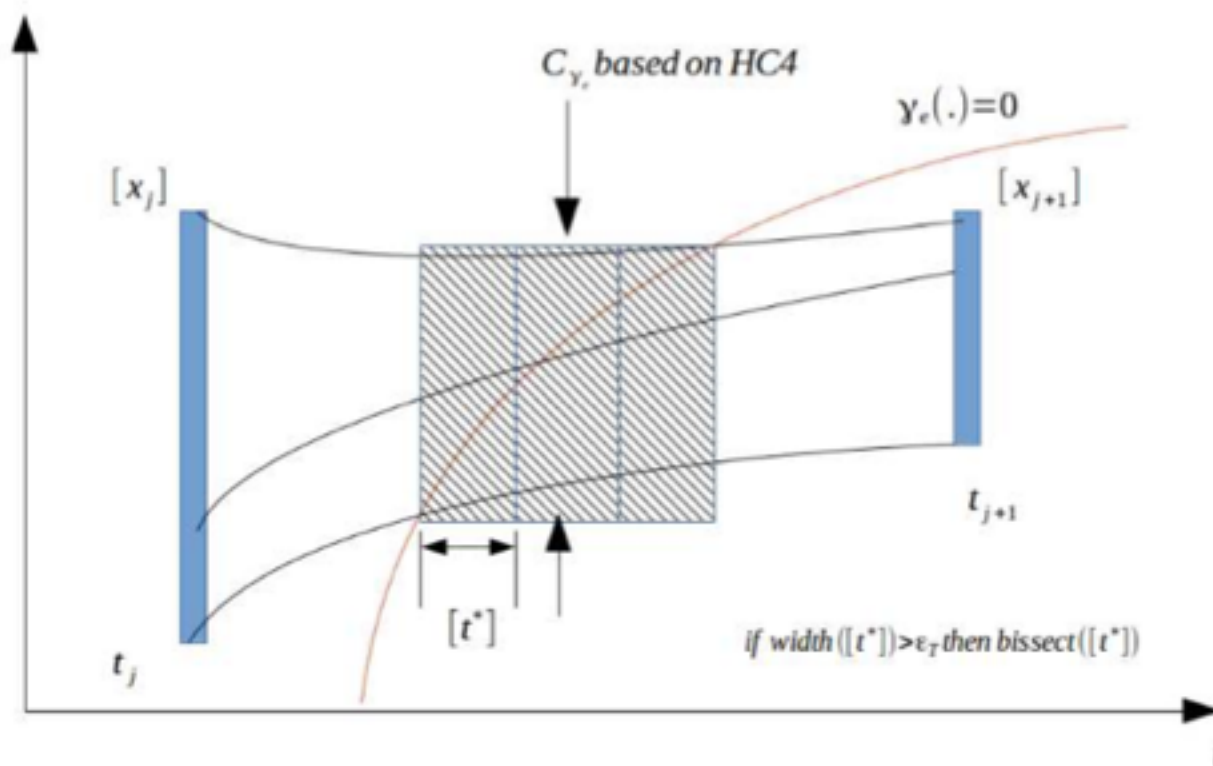
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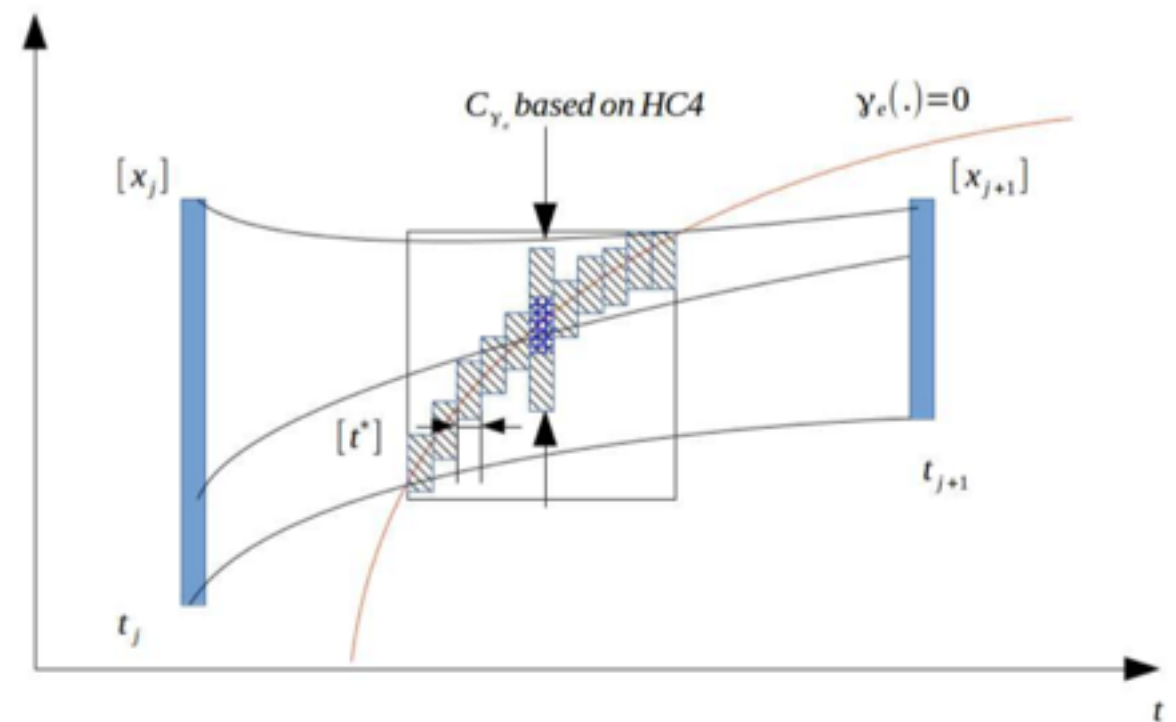
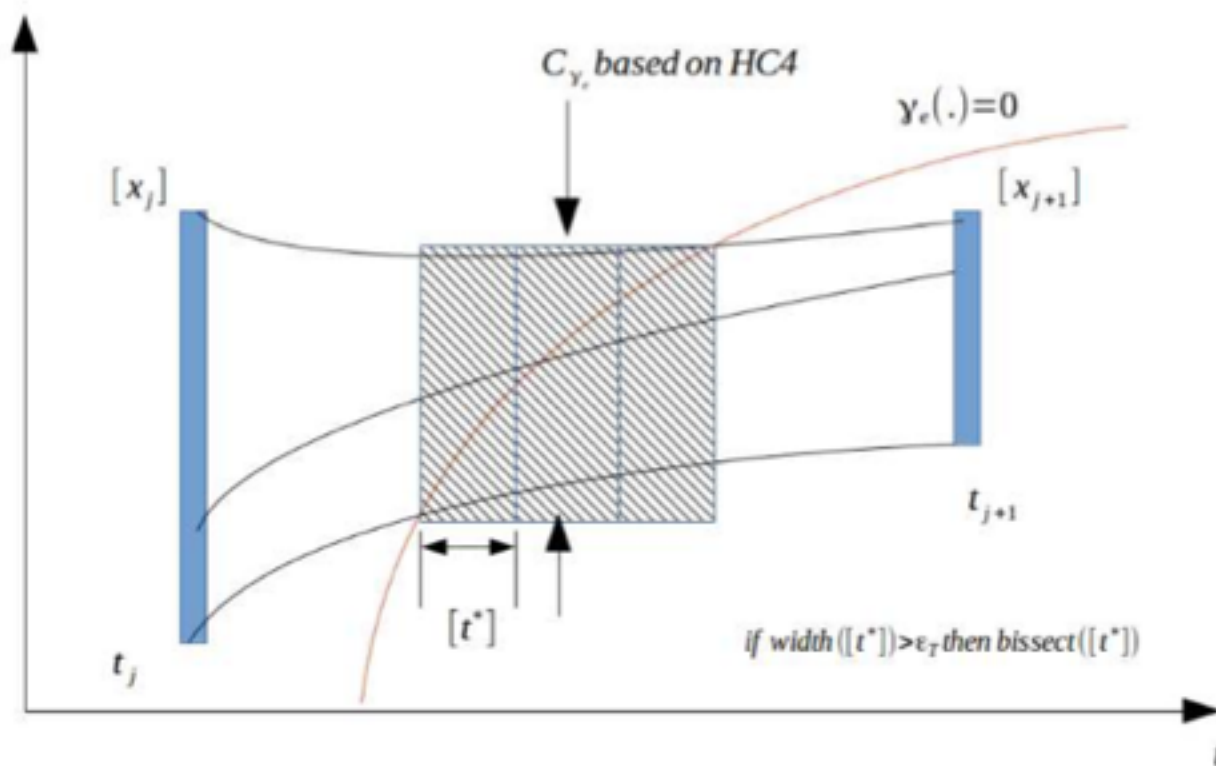
# Hybrid Reachability Analysis

- Detecting and localizing events
- Improved and enhanced version
  - (Maïga, et al., IEEE CDC 2013, ECC 2014)



# Hybrid Reachability Analysis

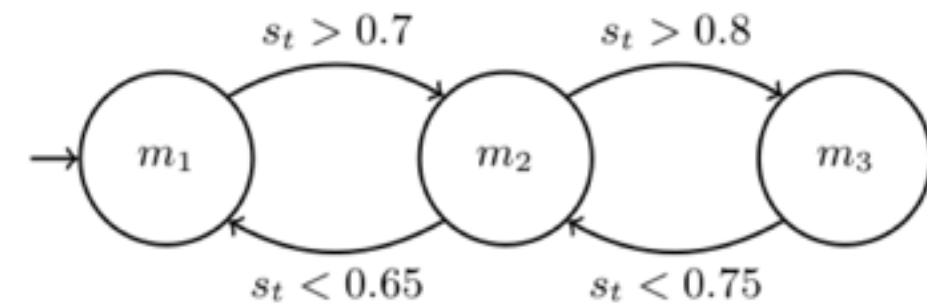
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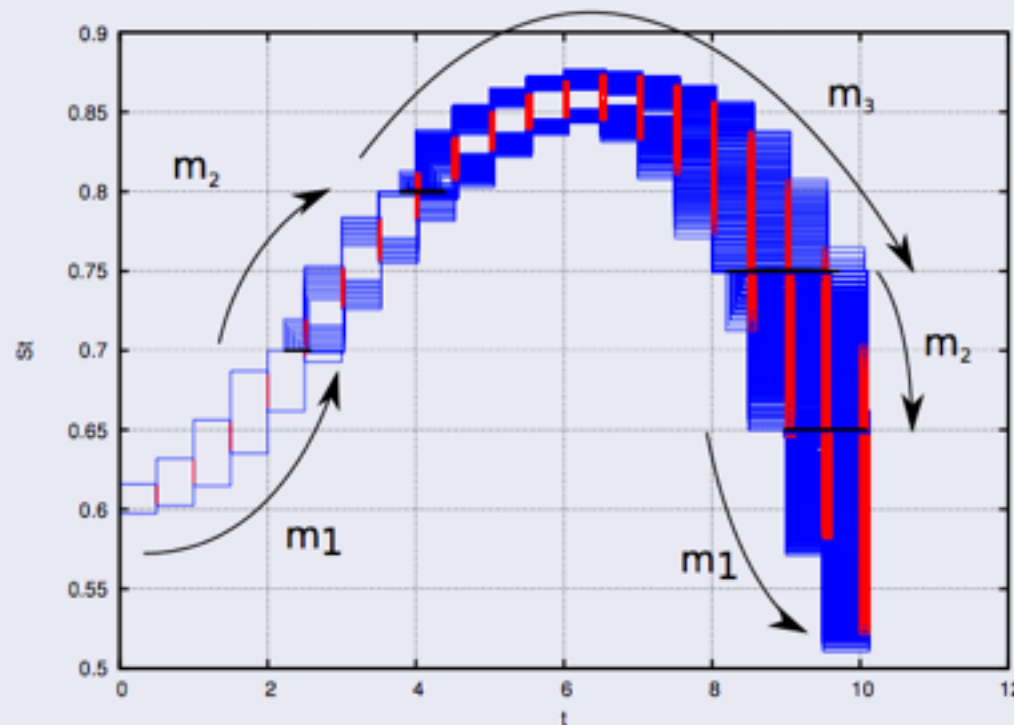


# Hybrid Reachability Analysis

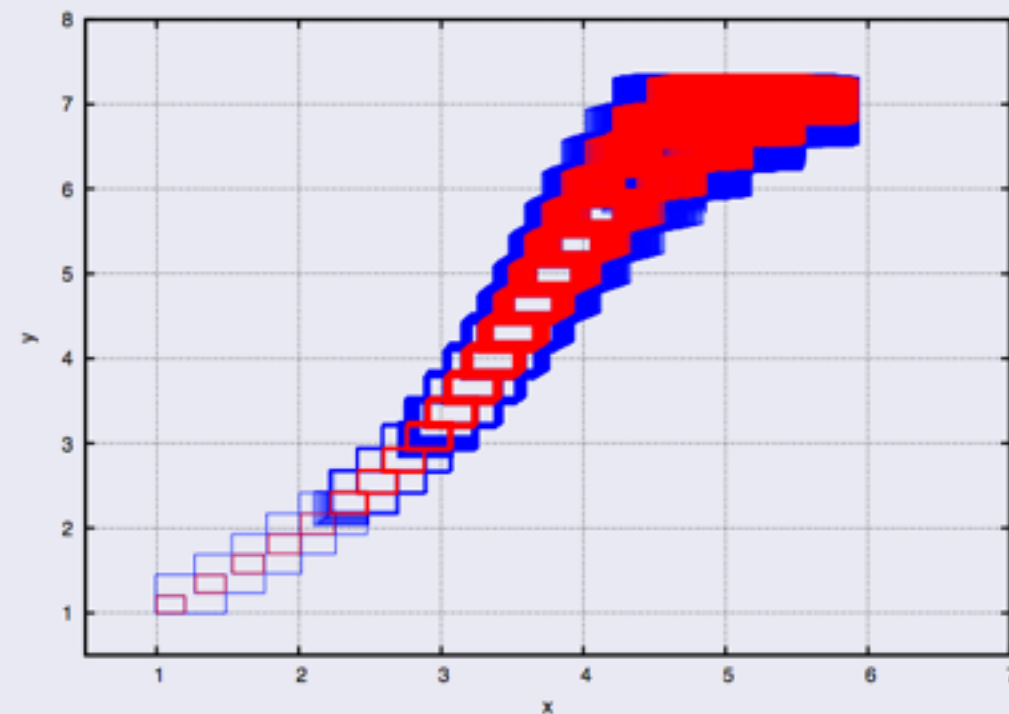
- Detecting and localizing events
- Improved and enhanced version
  - (Maïga, et al., IEEE CDC 2013, ECC 2014)



$\sigma = [0, 0.01]$  and  $h=0.5$



(e)  $S_t \times t$



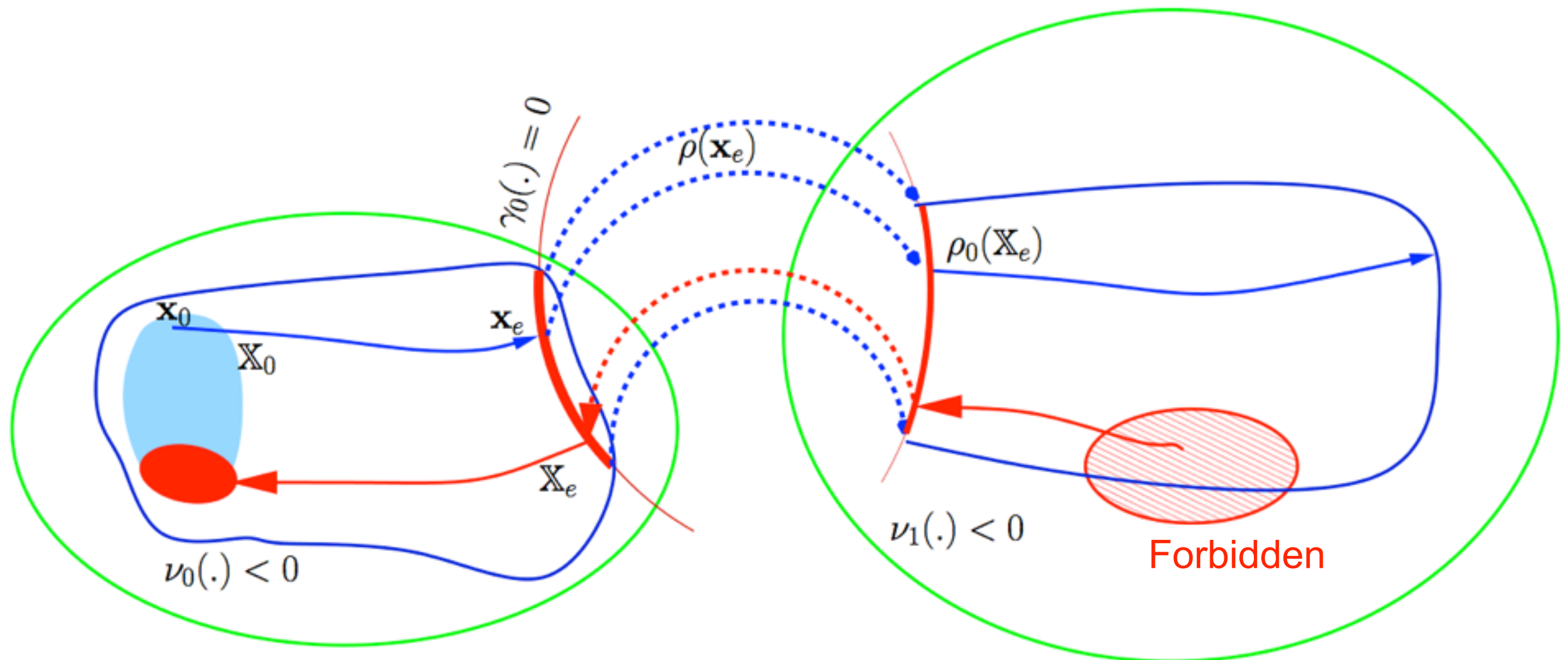
(f)  $Y \times X$  space

CPU times=87s with HC4 contractor  
 CPU times > 1h without HC4 contractor

# Verification of Hybrid Systems

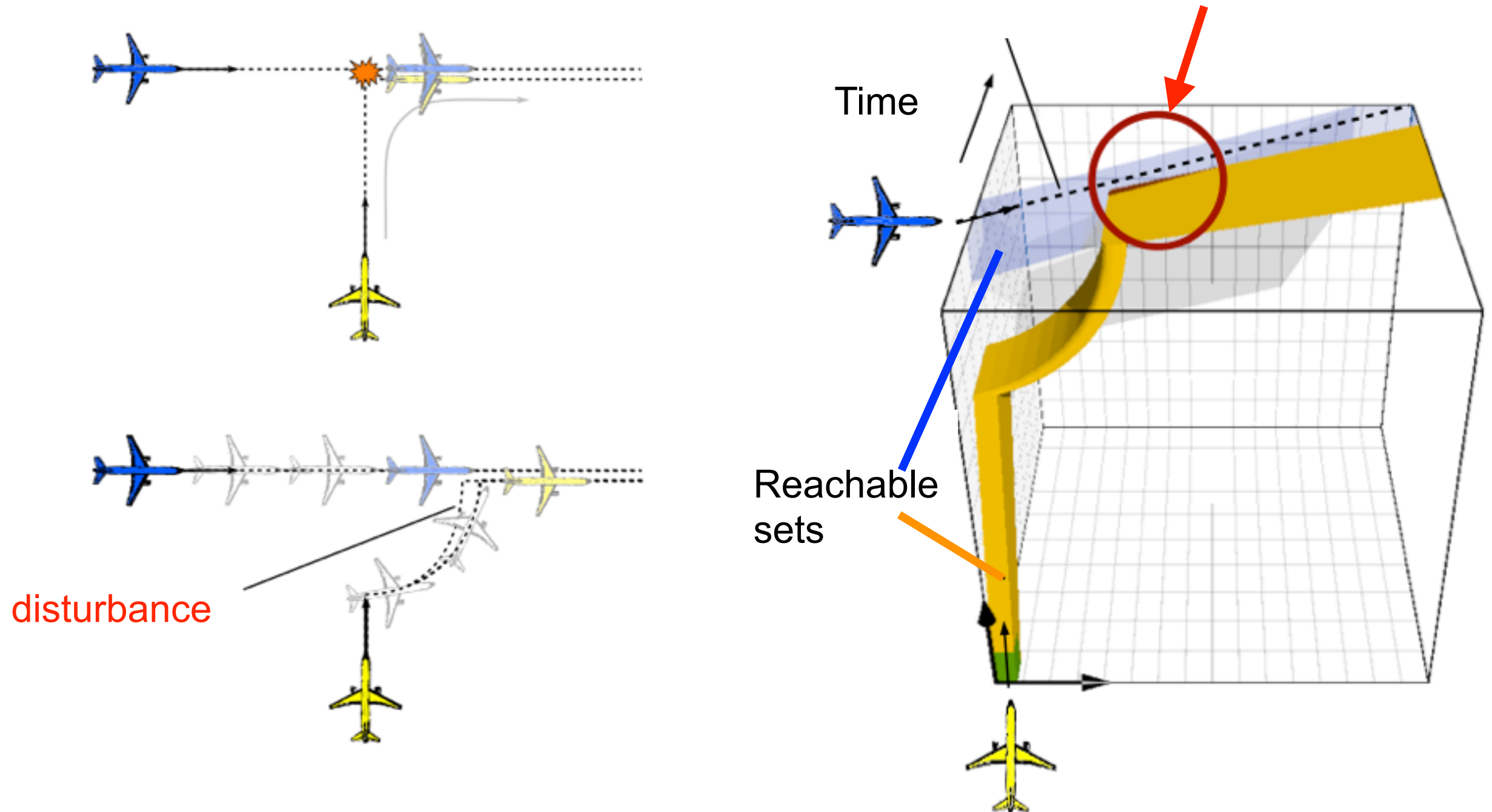
## ■ Verification :

- Reachability of a forbidden area



# Verification of Hybrid Systems

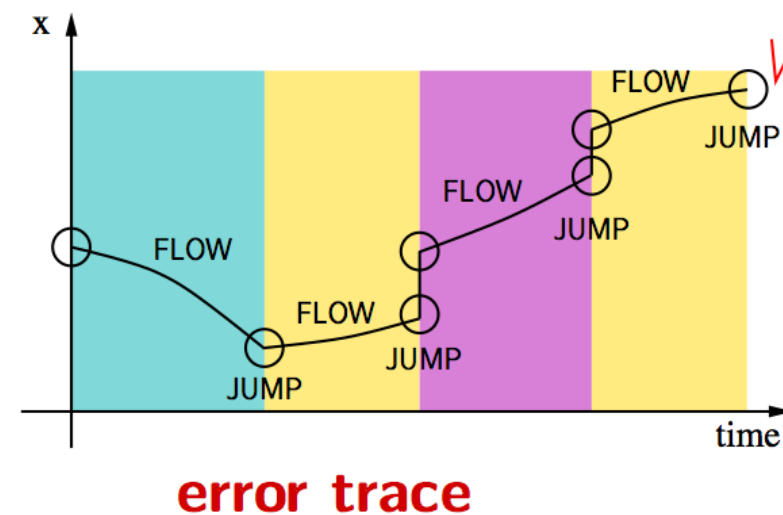
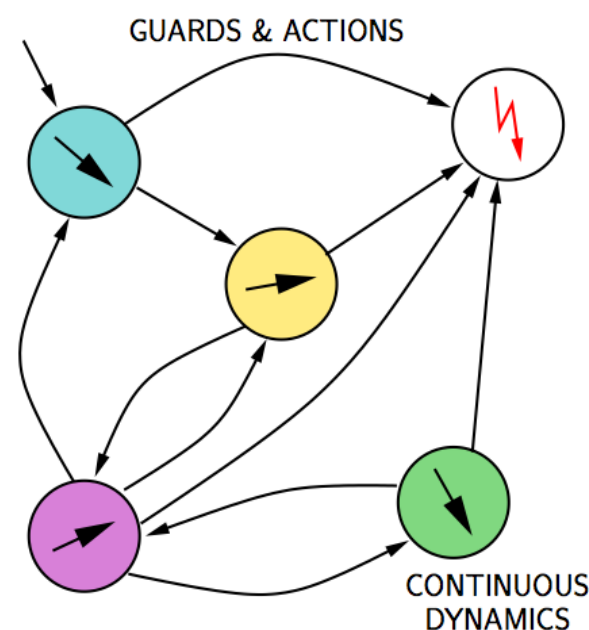
## ● Aircraft traffic control [Tomlin, et al.] Collision possible !



## ■ Bounded Model Checking

- Can the system reach an **unsafe** state within  $k$  (discrete or continuous) transition steps ?
- Check **satisfiability** of a **SAT Mod ODE** formula

$$\Phi_k := \text{init}[0] \wedge \text{trans}[0, 1] \wedge \dots \wedge \text{trans}[k - 1, k] \wedge \text{target}[k]$$

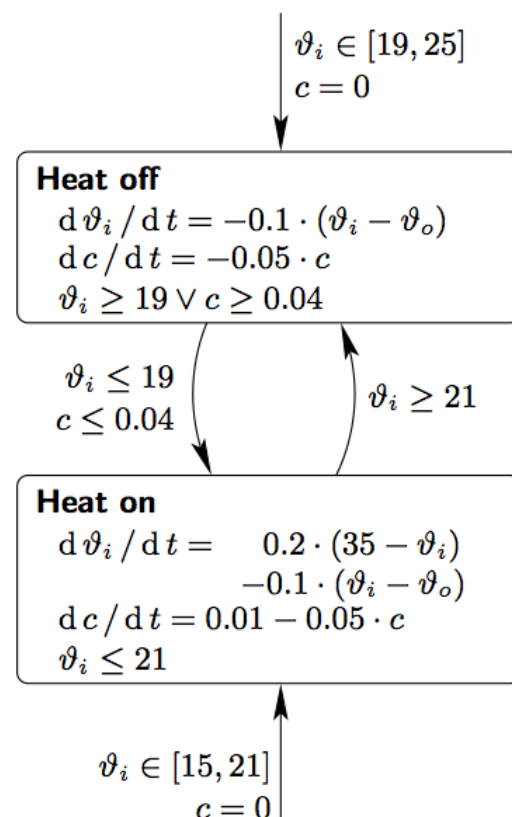




## ■ Bounded Model Checking

- Can the system reach an **unsafe** state within  $k$  (discrete or continuous) transition steps ?
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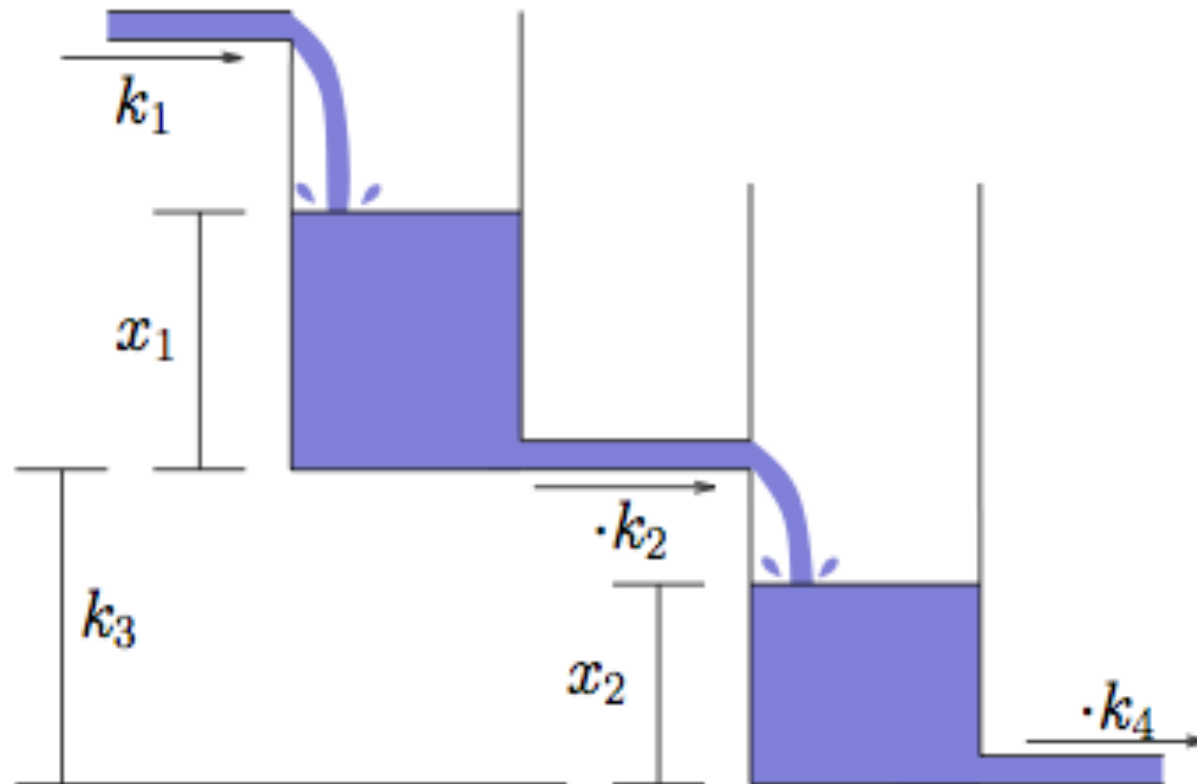


$$\begin{aligned} \text{init} = & -10 \leq \vartheta_o \leq 20 \wedge c = 0 \\ & \wedge \left( \begin{array}{l} 19 \leq \vartheta_i \leq 25 \wedge \neg \text{on} \\ \vee 15 \leq \vartheta_i \leq 21 \wedge \text{on} \end{array} \right) \\ \text{trans} = & \left( \begin{array}{l} \neg \text{on} \wedge \text{on}' \wedge \vartheta_i \leq 19 \wedge c \leq 0.04 \\ \wedge \vartheta_i' = \vartheta_i \wedge \vartheta_o' = \vartheta_o \wedge c' = c \end{array} \right) \\ & \vee \left( \begin{array}{l} \text{on} \wedge \neg \text{on}' \wedge \vartheta_i \geq 21 \\ \wedge \vartheta_i' = \vartheta_i \wedge \vartheta_o' = \vartheta_o \wedge c' = c \end{array} \right) \\ & \vee \left( \begin{array}{l} \neg \text{on} \wedge \neg \text{on}' \\ \wedge \frac{d\vartheta_i}{dt} = -0.1(\vartheta_i - \vartheta_o) \\ \wedge \frac{dc}{dt} = -0.05c \\ \wedge (\vartheta_i' \geq 19 \vee c' \geq 0.04) \wedge \vartheta_o' = \vartheta_o \end{array} \right) \\ & \vee \left( \begin{array}{l} \text{on} \wedge \text{on}' \\ \wedge \frac{d\vartheta_i}{dt} = 0.2 \cdot 35 - 0.3\vartheta_i + 0.1\vartheta_o \\ \wedge \frac{dc}{dt} = 0.01 - 0.05c \\ \wedge \vartheta_i' \leq 21 \wedge \vartheta_o' = \vartheta_o \end{array} \right) \\ \text{target} = & (c > 0.1) \end{aligned}$$

## ■ SAT mod ODE

- *Model:* **init** → definition of variables.  
**trans**[k,k+1] → transition dynamics.
- *Property:* **prop**
- *SAT solvers check the following formulas:*
  - init**  $\wedge$   $\neg$ **prop**
  - init**  $\wedge$  **trans**[0,1]  $\wedge$   $\neg$ **prop**
  - init**  $\wedge$  **trans**[0,1]  $\wedge$  **trans**[1,2]  $\wedge$   $\neg$ **prop**
  - init**  $\wedge$  **trans**[0,1]  $\wedge$  **trans**[1,2]  $\wedge$  **trans**[2,3]  $\wedge$   $\neg$ **prop** ...
- **If one formula is satisfiable → Property is violated !**

## ■ Example : 2-tanks system



For  $x_2 > k_3$ :

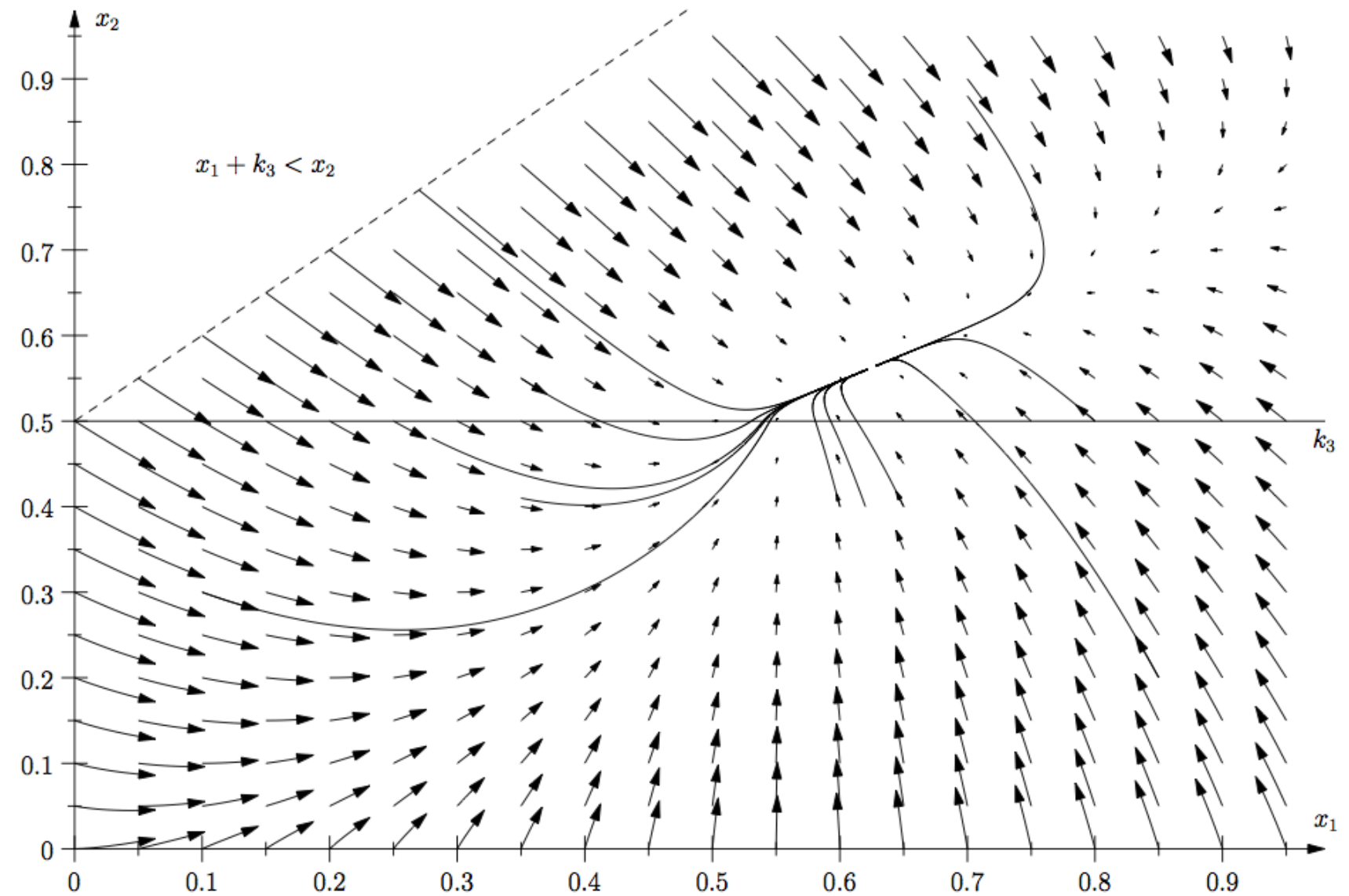
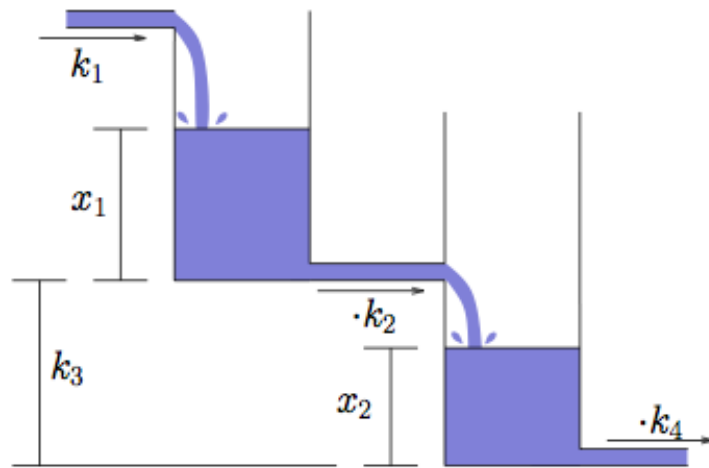
$$\begin{pmatrix} \dot{x}_1 \\ \dot{x}_2 \end{pmatrix} = \begin{pmatrix} k_1 - k_2\sqrt{x_1 - x_2 + k_3} \\ k_2\sqrt{x_1 - x_2 + k_3} - k_4\sqrt{x_2} \end{pmatrix}$$

For  $x_2 \leq k_3$ :

$$\begin{pmatrix} \dot{x}_1 \\ \dot{x}_2 \end{pmatrix} = \begin{pmatrix} k_1 - k_2\sqrt{x_1} \\ k_2\sqrt{x_1} - k_4\sqrt{x_2} \end{pmatrix}$$

$$k_1 = 0.75, k_2 = 1, k_3 = 0.5, k_4 = 1$$

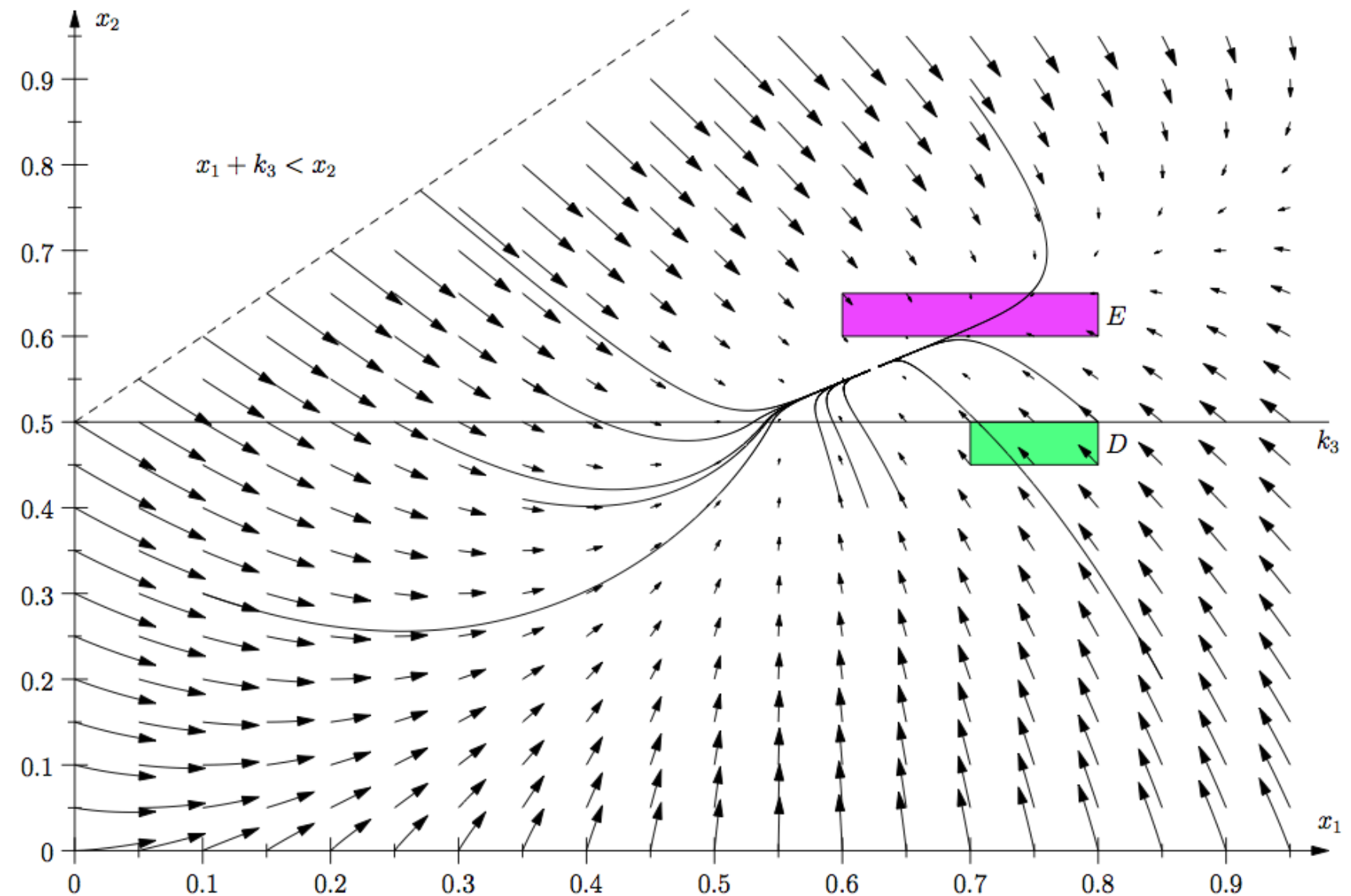
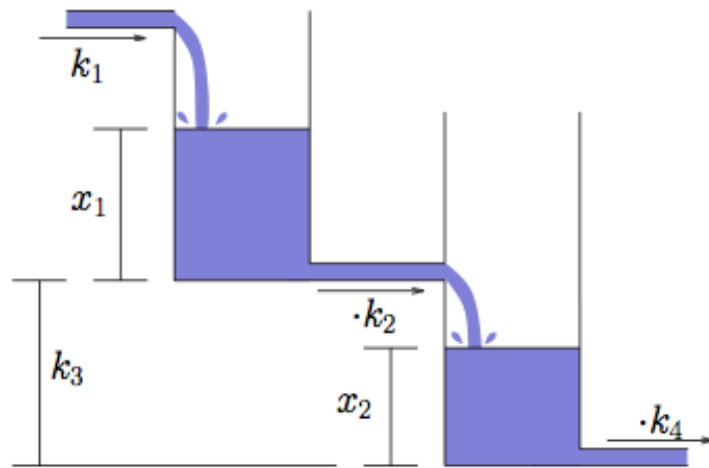
## ■ Example : 2-tanks system



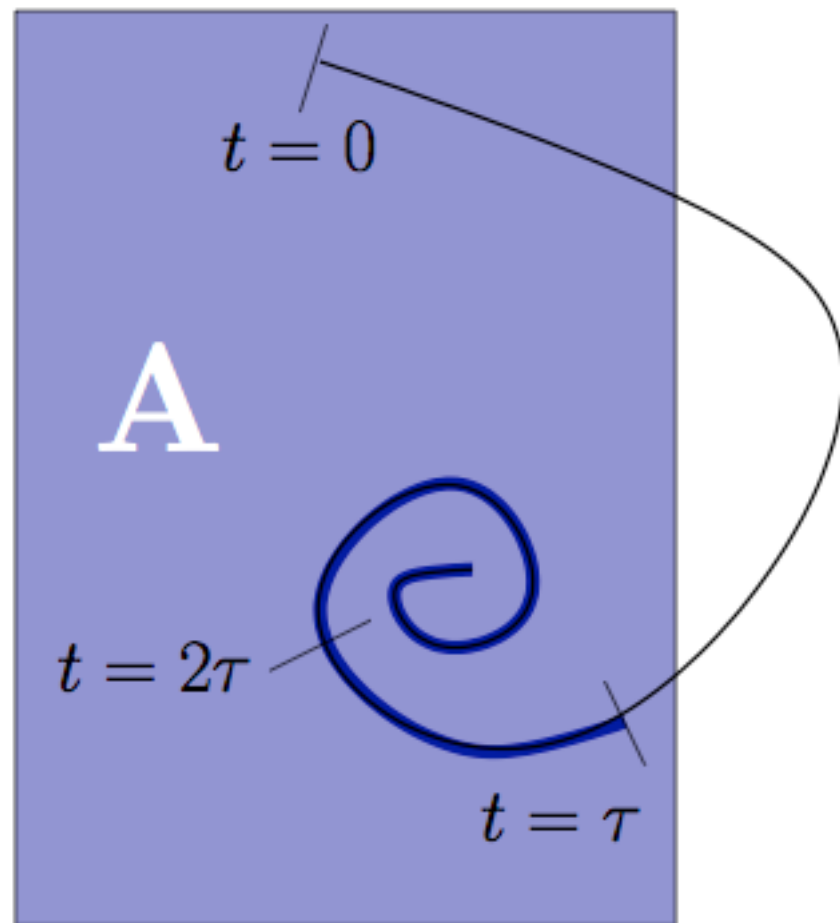


**E non reachable from D.** [Eggers, Ramdani, Nediakov, Fränzle, 2011]

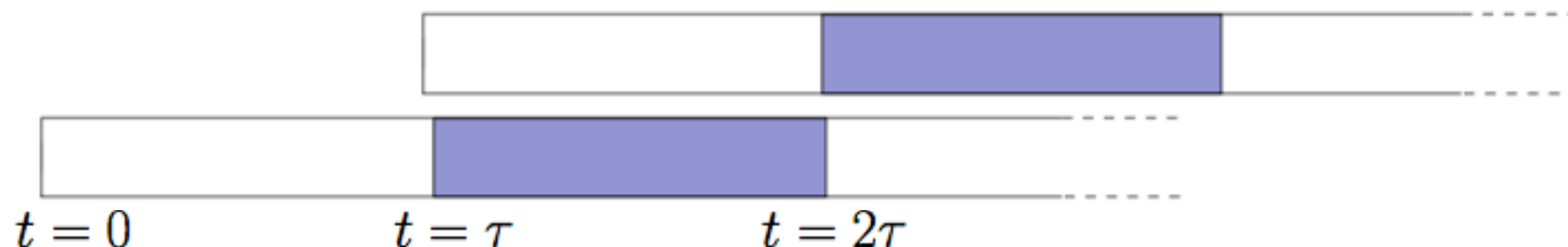
iSAT-ODE: Proof in 260s CPU 2.4 GHz AMD Opteron



[Podelski et Wagner, 2007] [Eggers, Ramdani, Nediaklov, Fränze, 2011]



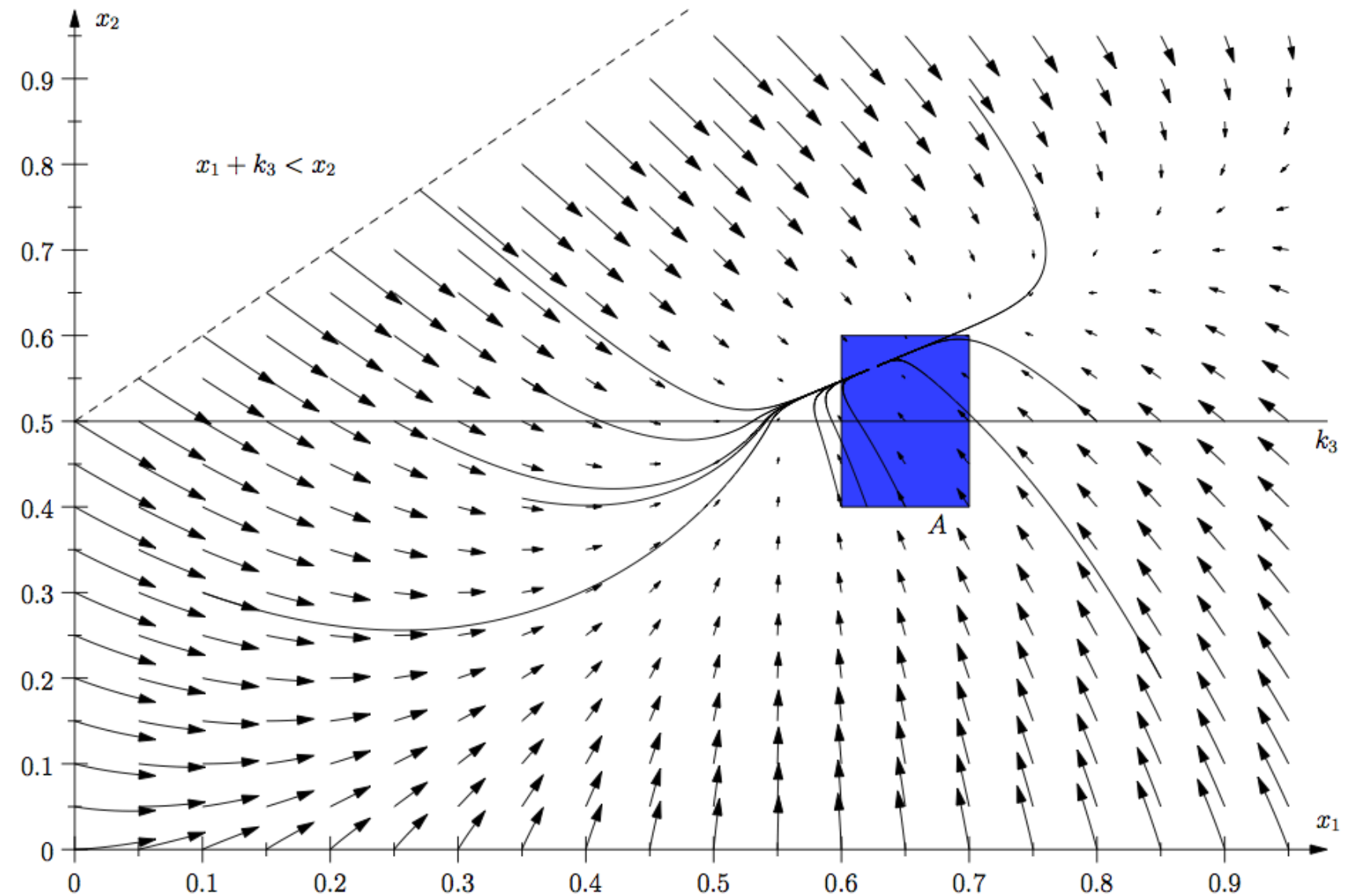
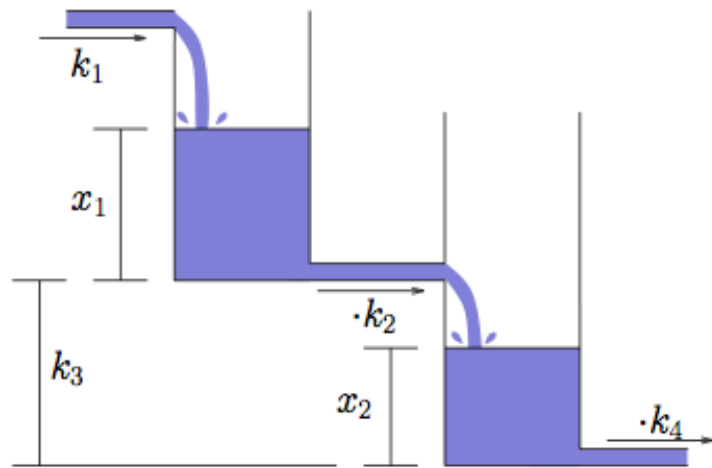
- Proof: a trajectory starting in  $A$ , stays in  $A$  during  $[\tau, 2\tau]$
- SAT mod ODE formula  
Target :  
Non reached at  $2\tau$   
or left  $A$  during  $[\tau, 2\tau]$
- If UNSAT, recurrence,  
time-invariance, infinite time property.



# Region stability

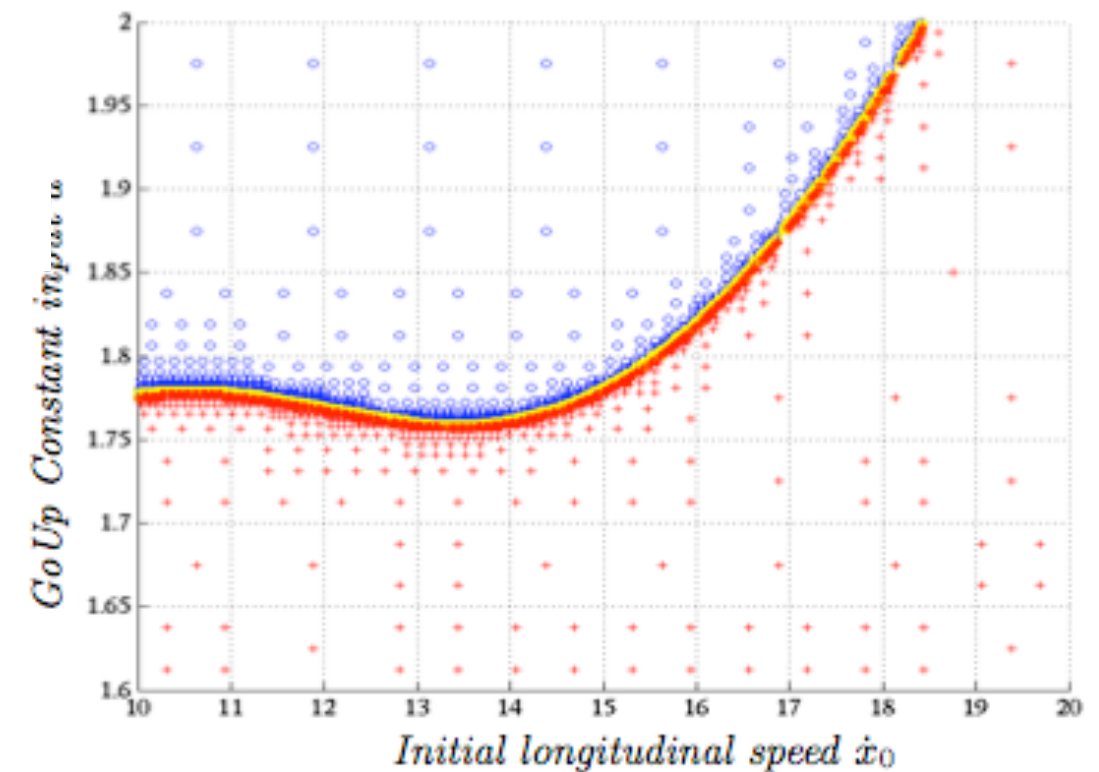
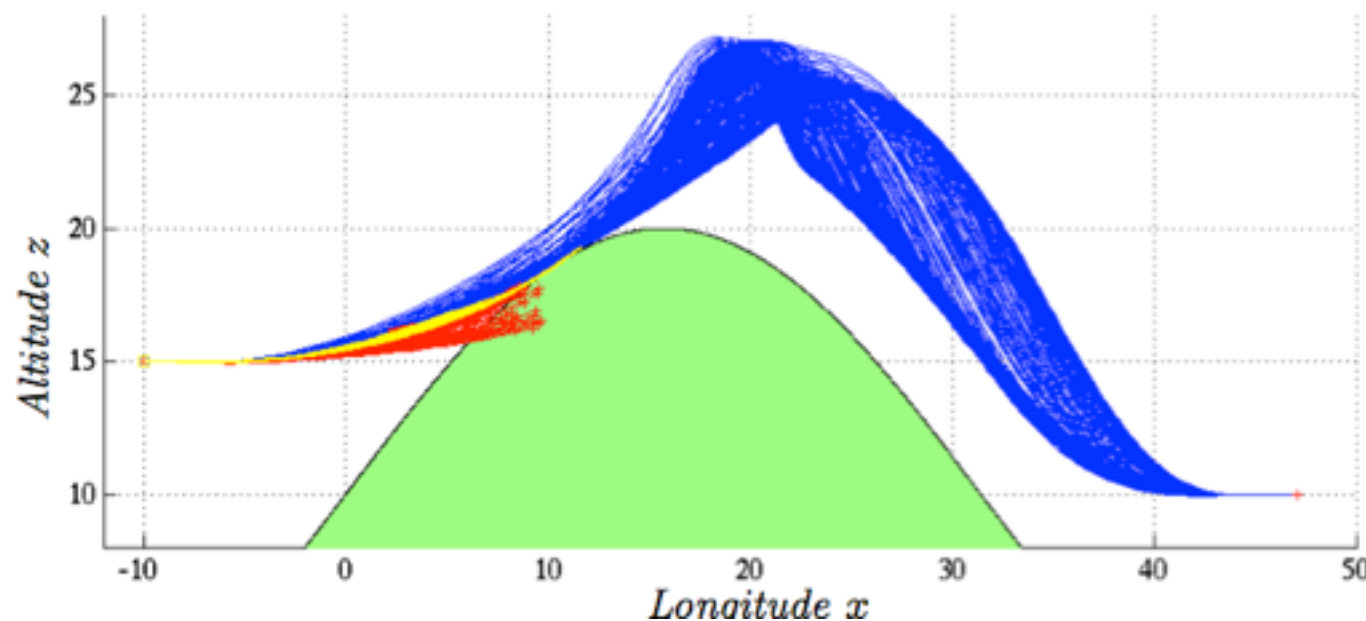
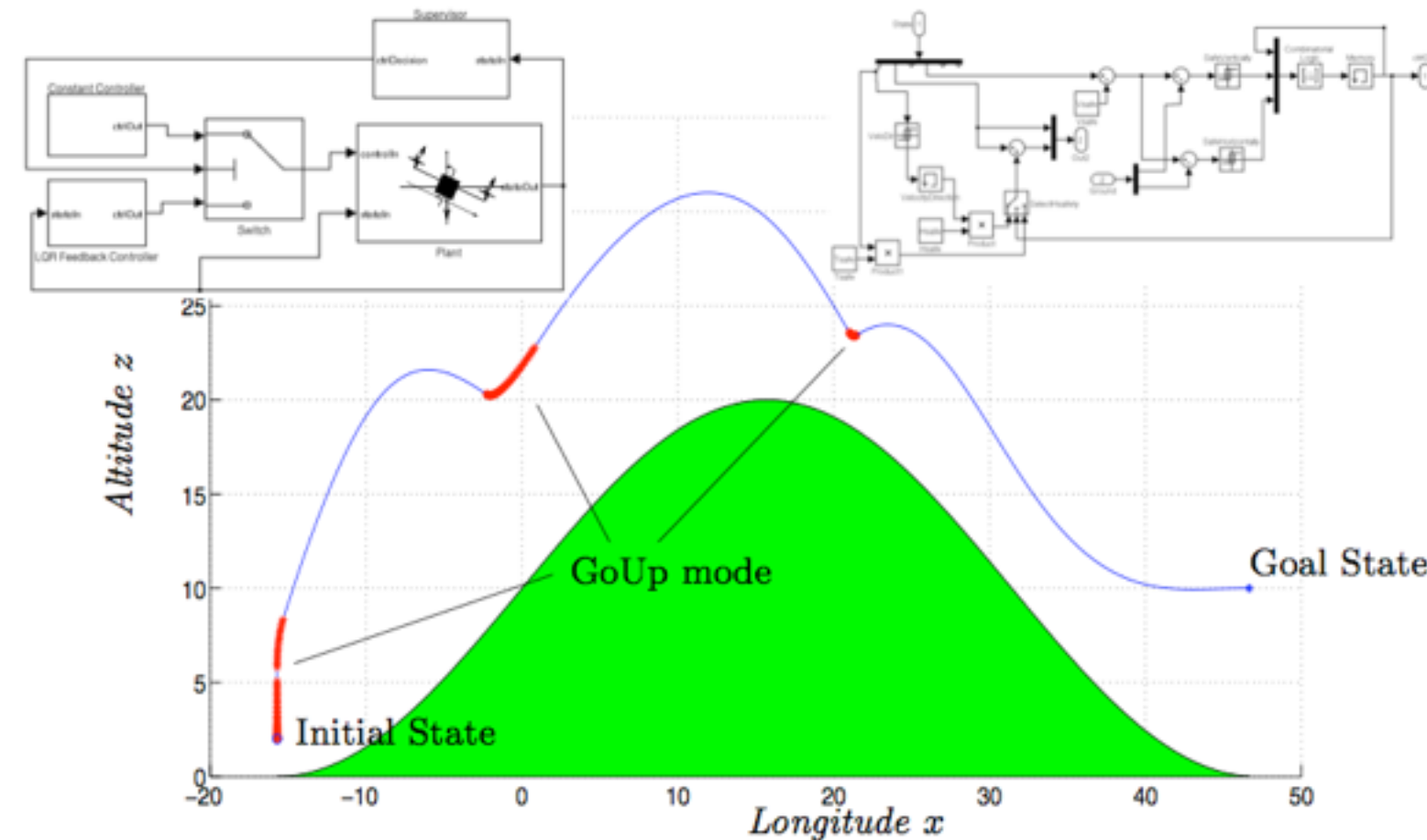
[Eggers, Ramdani, Nedialkov, Fränzle, 2011]

iSAT-ODE: proof in 150s CPU 2.4 GHz AMD Opteron



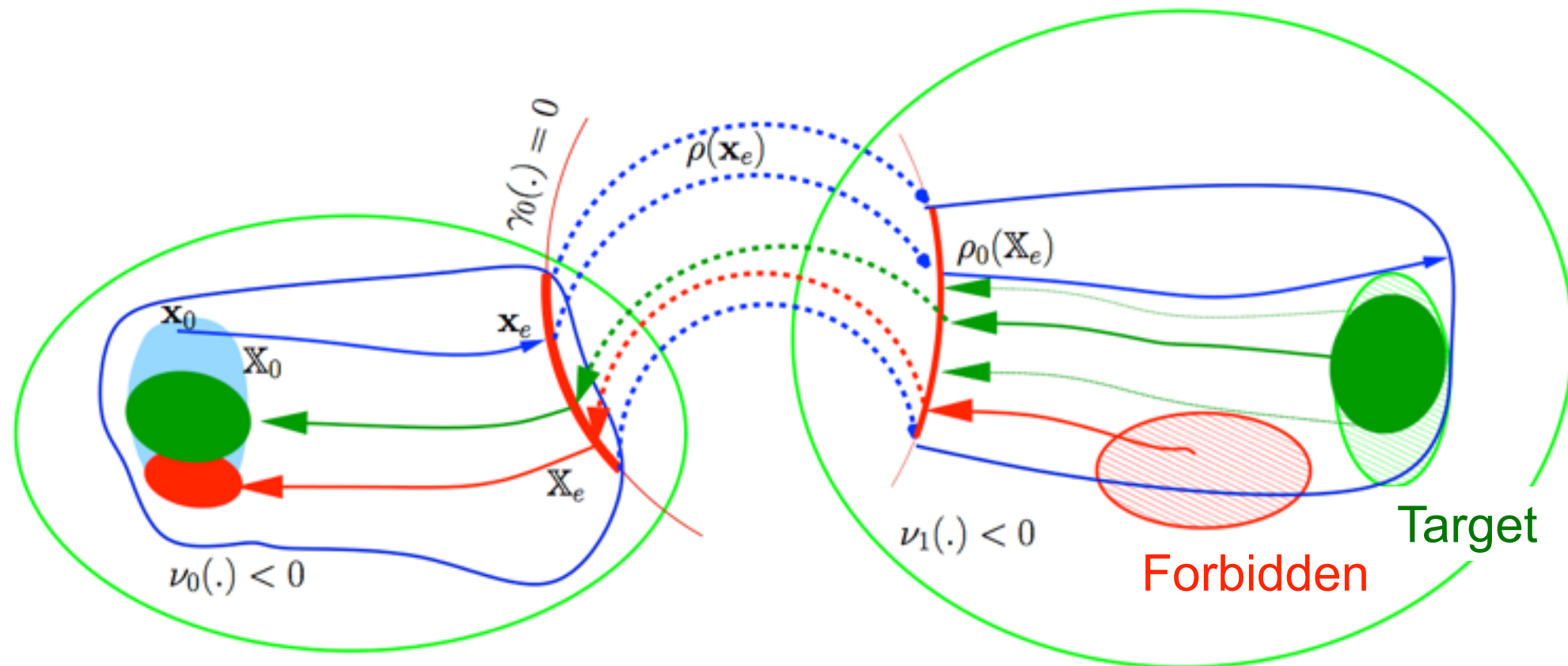
# Hybrid systems synthesis

**A. Donzé, B. Krogh & A. Rajhans.**  
*Parameter synthesis for hybrid systems with an application to simulink models.*  
**HSCC 2009:165-179.**



# Hybrid Cyber-Physical Systems

## ■ Hybrid reachability analysis



- Verification
- Synthesis
- Set-theoretic estimation

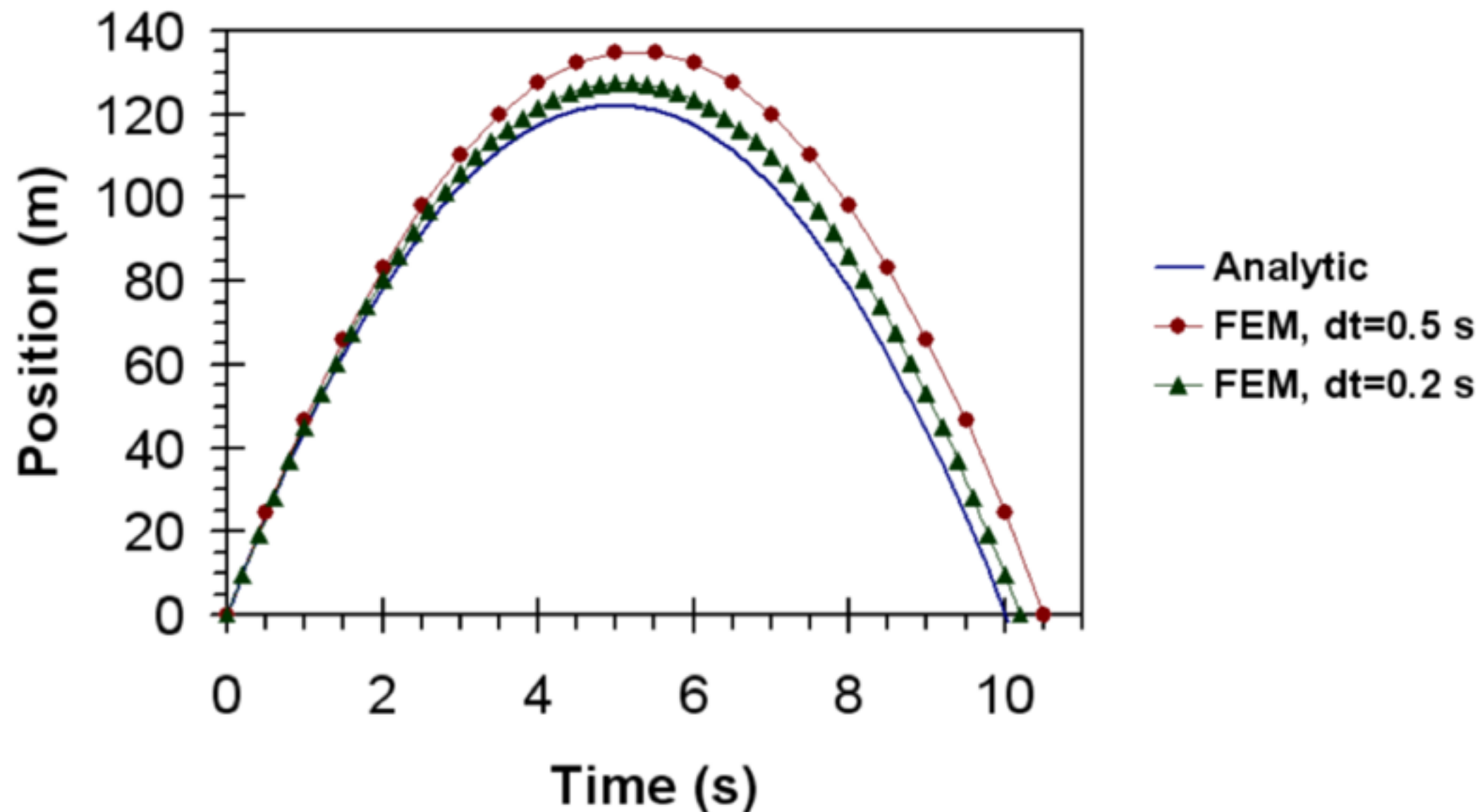


- Hybrid Reachability : Motivations
  - Analysis of complex dynamical systems
  - Reachability-based methods
  
- Nonlinear reachability
  - Interval Taylor methods
  - Bracketing enclosures
  - Software implementation

# Nonlinear Set Integration

## ■ Non-guaranteed integration method ...

$$\dot{\mathbf{x}}(t) = f(\mathbf{x}, \mathbf{p}, t), \quad t_0 \leq t \leq t_N, \quad \mathbf{x}(t_0) \in [\mathbf{x}_0], \quad \mathbf{p} \in [\mathbf{p}]$$



# Nonlinear Set Integration

- **Guaranteed set integration with Taylor methods**
  - (Moore,66) (Lohner,88) (Rihm,94) (Berz,98) (Nedialkov,99)

$$\dot{\mathbf{x}}(t) = f(\mathbf{x}, \mathbf{p}, t), \quad t_0 \leq t \leq t_N, \quad \mathbf{x}(t_0) \in [\mathbf{x}_0], \quad \mathbf{p} \in [\mathbf{p}]$$

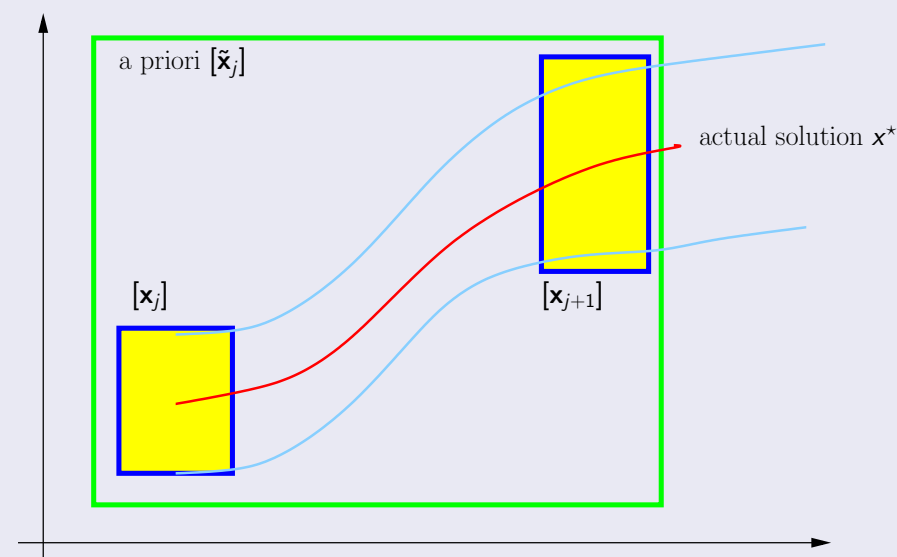
# Nonlinear Set Integration

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Time grid  $\rightarrow t_0 < t_1 < t_2 < \dots < t_N$



- Proof of existence
- Yield *a priori* solution  $[\tilde{\mathbf{x}}_j] : \forall \tau \in [t_j, t_{j+1}] \quad \mathbf{x}(\tau) \in [\tilde{\mathbf{x}}_j]$

# Nonlinear Set Integration

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$$[\mathbf{x}_j] + [0, h]\mathbf{f}([\tilde{\mathbf{x}}_j]) \subseteq [\tilde{\mathbf{x}}_j]$$



# Nonlinear Set Integration

## ■ Guaranteed set integration with Taylor methods

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*a priori enclosure* (entrée :  $[\mathbf{x}_j]$ ,  $h$ ,  $\alpha$  ; sortie :  $[\tilde{\mathbf{x}}_j]$ )

1. Initialisation :  $[\tilde{\mathbf{x}}_j] := [\mathbf{x}_j] + [0, h] \mathbf{f}([\mathbf{x}_j])$  ;
2. tant que  $([\mathbf{x}_j] + [0, h] \mathbf{f}([\tilde{\mathbf{x}}_j]) \not\subset [\tilde{\mathbf{x}}_j])$

$$\begin{aligned} [\tilde{\mathbf{x}}_j] &:= [\tilde{\mathbf{x}}_j] + [-\alpha, \alpha] ||[\tilde{\mathbf{x}}_j]|| \\ h &:= h/2 \end{aligned}$$

fin

*An Effective High-Order Interval Method for Validating Existence and Uniqueness of the Solution of an IVP for an ODE, Nedialko S. Nedialkov, Kenneth R. Jackson, and John D. Pryce, Reliable Computing 7(6) :449 - 465, 2001.*

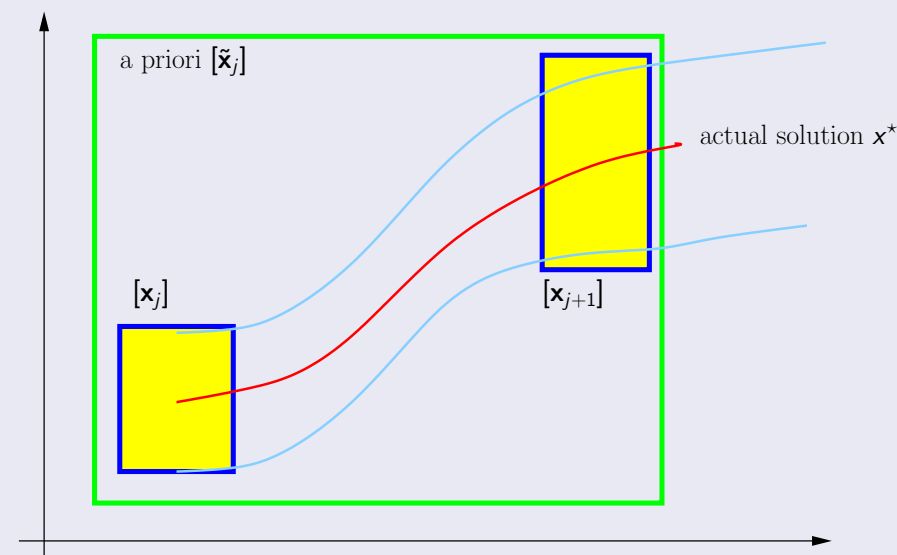
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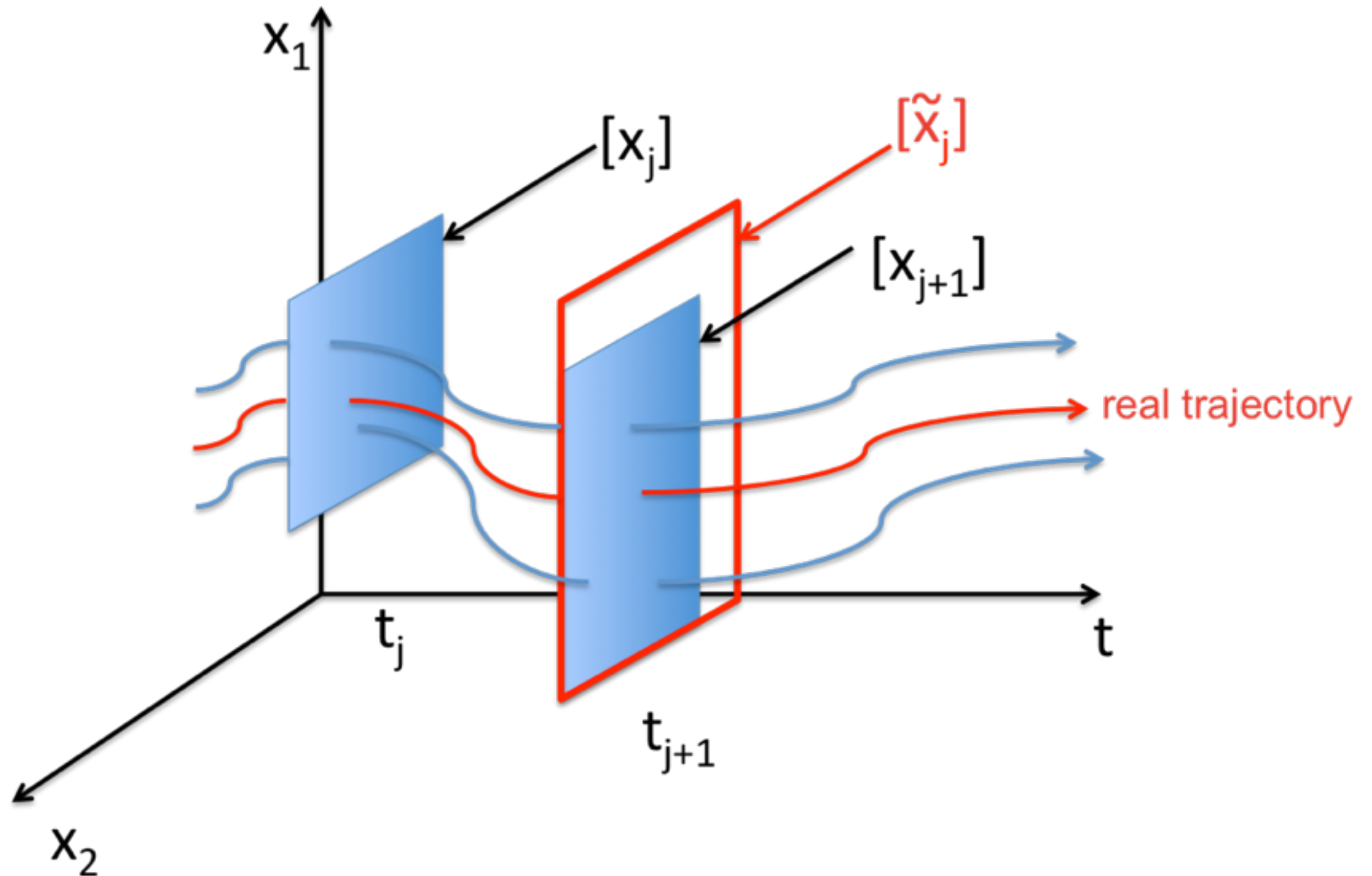
Time grid  $\rightarrow t_0 < t_1 < t_2 < \dots < t_N$



- Compute tight enclosure  $[\mathbf{x}_{j+1}] \ni \mathbf{x}(t_{j+1})$

$$[\mathbf{x}_{j+1}] = [\mathbf{x}_j] + \sum_{i=1}^{k-1} (t_{j+1} - t_j)^i \mathbf{f}^{[i]}([\mathbf{x}_j], [\mathbf{p}]) + (t_{j+1} - t_j)^k \mathbf{f}^{[k]}([\tilde{\mathbf{x}}_j], [\mathbf{p}])$$

# Nonlinear Set Integration



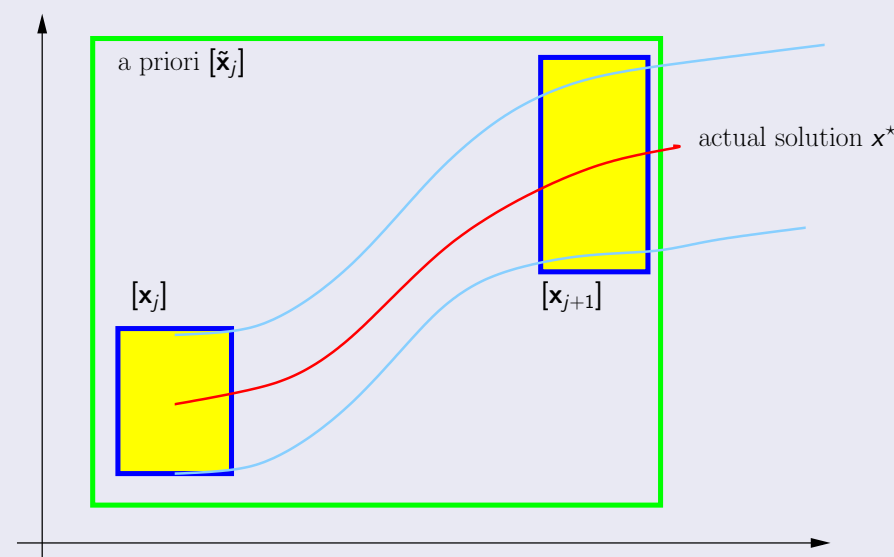
# Nonlinear Set Integration

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Time grid  $\rightarrow t_0 < t_1 < t_2 < \dots < t_N$



- **Analytical solution** for  $[\mathbf{x}](t)$ ,  $t \in [t_j, t_{j+1}]$

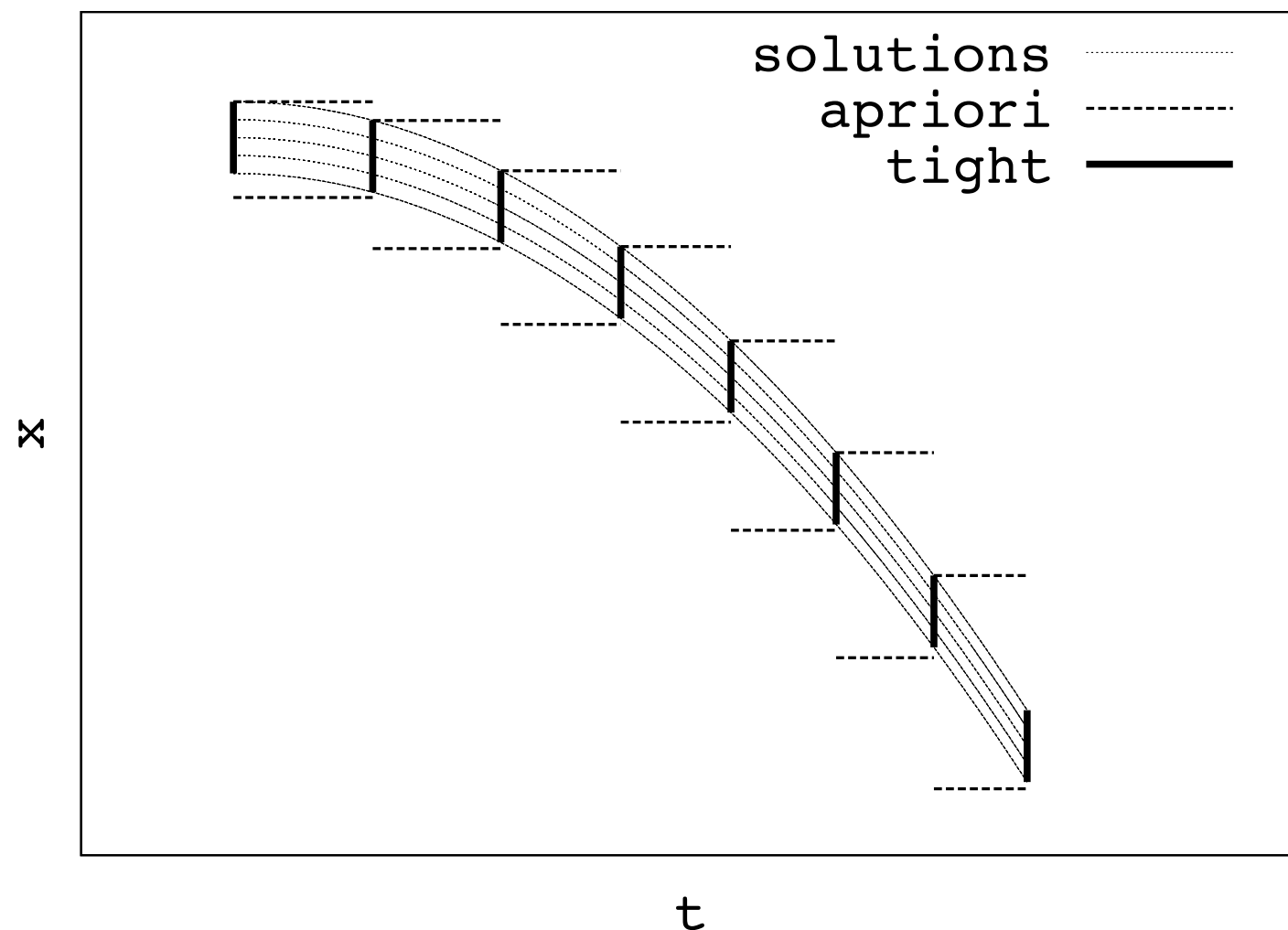
$$[\mathbf{x}](t) = [\mathbf{x}_j] + \sum_{i=1}^{k-1} (t - t_j)^i \mathbf{f}^{[i]}([\mathbf{x}_j], [\mathbf{p}]) + (t - t_j)^k \mathbf{f}^{[k]}([\tilde{\mathbf{x}}_j], [\mathbf{p}])$$

# Nonlinear Set Integration

## ■ Guaranteed set integration with Taylor methods

● (Moore,66) (Lohner,88) (Rihm,94) (Berz,98) (Nedialkov,99)

$$\dot{\mathbf{x}}(t) = f(\mathbf{x}, \mathbf{p}, t), \quad t_0 \leq t \leq t_N, \quad \mathbf{x}(t_0) \in [\mathbf{x}_0], \quad \mathbf{p} \in [\mathbf{p}]$$



# Nonlinear Set Integration

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$$\mathbf{f}^{[1]} = \mathbf{x}^{(1)} = \mathbf{f}$$

$$\mathbf{f}^{[2]} = \frac{1}{2} \mathbf{x}^{(2)} = \frac{1}{2} \frac{d\mathbf{f}}{d\mathbf{x}} \mathbf{f}$$

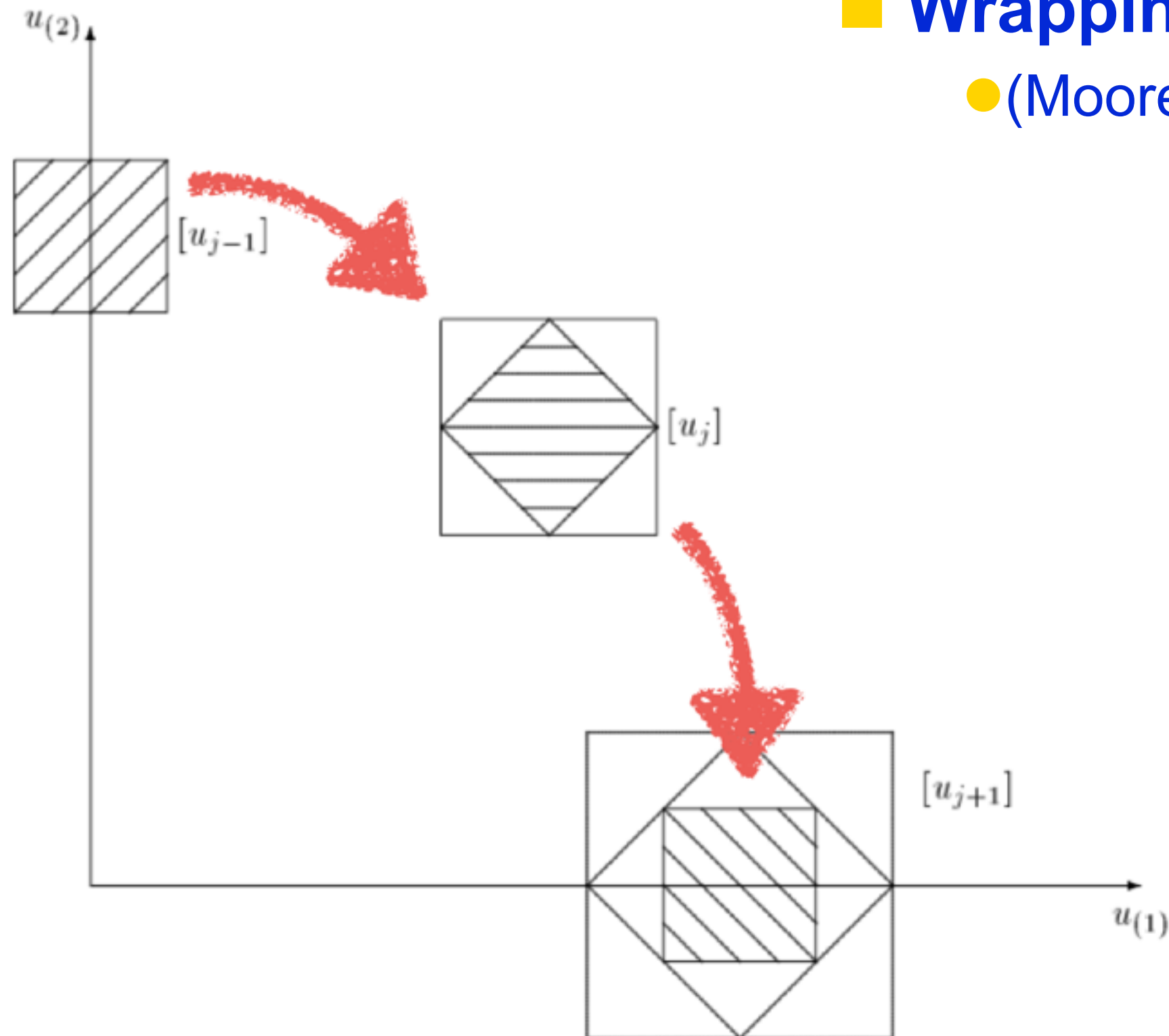
$$\mathbf{f}^{[i]} = \frac{1}{i!} \mathbf{x}^{(i)} = \frac{1}{i} \frac{d\mathbf{f}^{[i-1]}}{d\mathbf{x}} \mathbf{f}, \quad i \geq 2$$



# Nonlinear Set Integration

## ■ Wrapping effect

● (Moore, 66)



# Nonlinear Set Integration

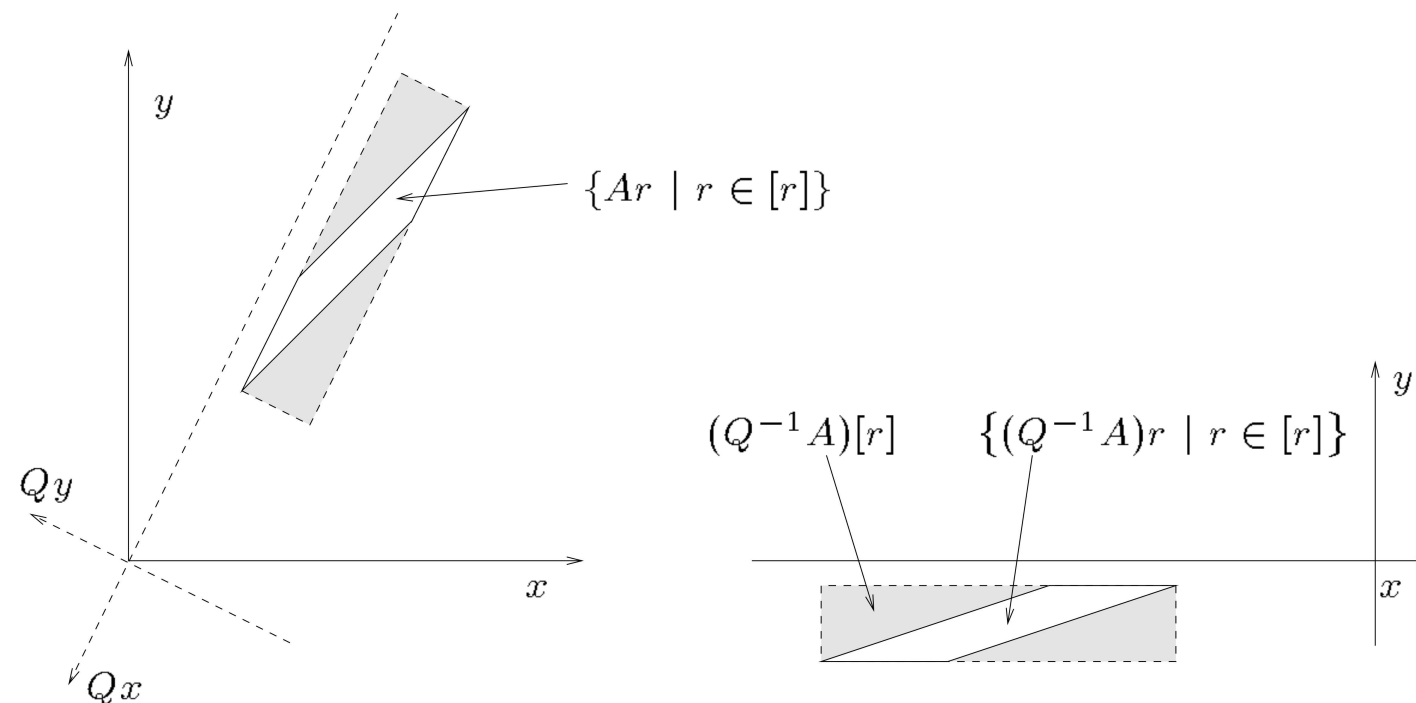
## ■ Guaranteed set integration with Taylor methods

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Mean-value approach

$$[\mathbf{x}](t) \in \{ \mathbf{v}(t) + \mathbf{A}(t)\mathbf{r}(t) \mid \mathbf{v}(t) \in [\mathbf{v}](t), \mathbf{r}(t) \in [\mathbf{r}](t) \}.$$



# Nonlinear Set Integration

## ■ Guaranteed set integration with Taylor methods

- (Moore,66) (Lohner,88) (Rihm,94) (Berz,98) (Nedialkov,99)

### ● Complexity

- *Work per step is of polynomial complexity*

- Computing Taylor coefficients  $\rightarrow o(k^2)$

- Linear algebra  $\rightarrow o(n^3)$

- **In practice** : Obtaining Taylor coefficients ...

- **FADBAD++** ([www.fadbad.com](http://www.fadbad.com))

Flexible Automatic differentiation using templates  
and operator overloading in C++

# VNODE-LP

**An Interval Solver for Initial Value Problems in  
Ordinary Differential Equations**

Ned Nediaklov  
[nediaklov@mcmaster.ca](mailto:nediaklov@mcmaster.ca)

---

VNODE-LP is a C++ package for computing bounds on solutions in IVPs for ODEs. In contrast to traditional ODE solvers, which compute approximate solutions, this solver tries to prove that a unique solution to a problem exists and then computes bounds that contain this solution. Such bounds can be used to help prove a theoretical result, check if a solution satisfies a condition in a safety-critical calculation, or simply to verify the results produced by a traditional ODE solver.

This package is a successor of the [VNODE](#) package of N. Nediaklov. A distinctive feature of the present solver is that it is developed entirely using [Literate Programming](#). As a result, the correctness of VNODE-LP's implementation can be examined easier than the correctness of VNODE: the theory, documentation, and source code are produced from the same [CWEB](#) files.

download

## ■ Comparison theorems for differential inequalities

- Müller's existence theorem (1936)

$$\text{If } \left\{ \begin{array}{l} \forall t \in [t_0, t_N], \forall \mathbf{x} \in \mathbb{D}, \forall \mathbf{p} \in \mathbb{P} \\ \forall i \min_{x_i = \omega_i} f_i(\mathbf{x}, \mathbf{p}, t) \geq D^\pm \omega_i(t) \\ \forall i \max_{x_i = \Omega_i} f_i(\mathbf{x}, \mathbf{p}, t) \leq D^\pm \Omega_i(t) \\ \omega(t_0) \leq \mathbf{x}(t_0) \leq \Omega(t_0) \end{array} \right. \Rightarrow \left\{ \begin{array}{l} \text{solution exists} \\ \forall t \in [t_0, t_N], \\ \omega(t) \leq \mathbf{x}(t) \leq \Omega(t) \end{array} \right.$$

- **Bracketing systems**

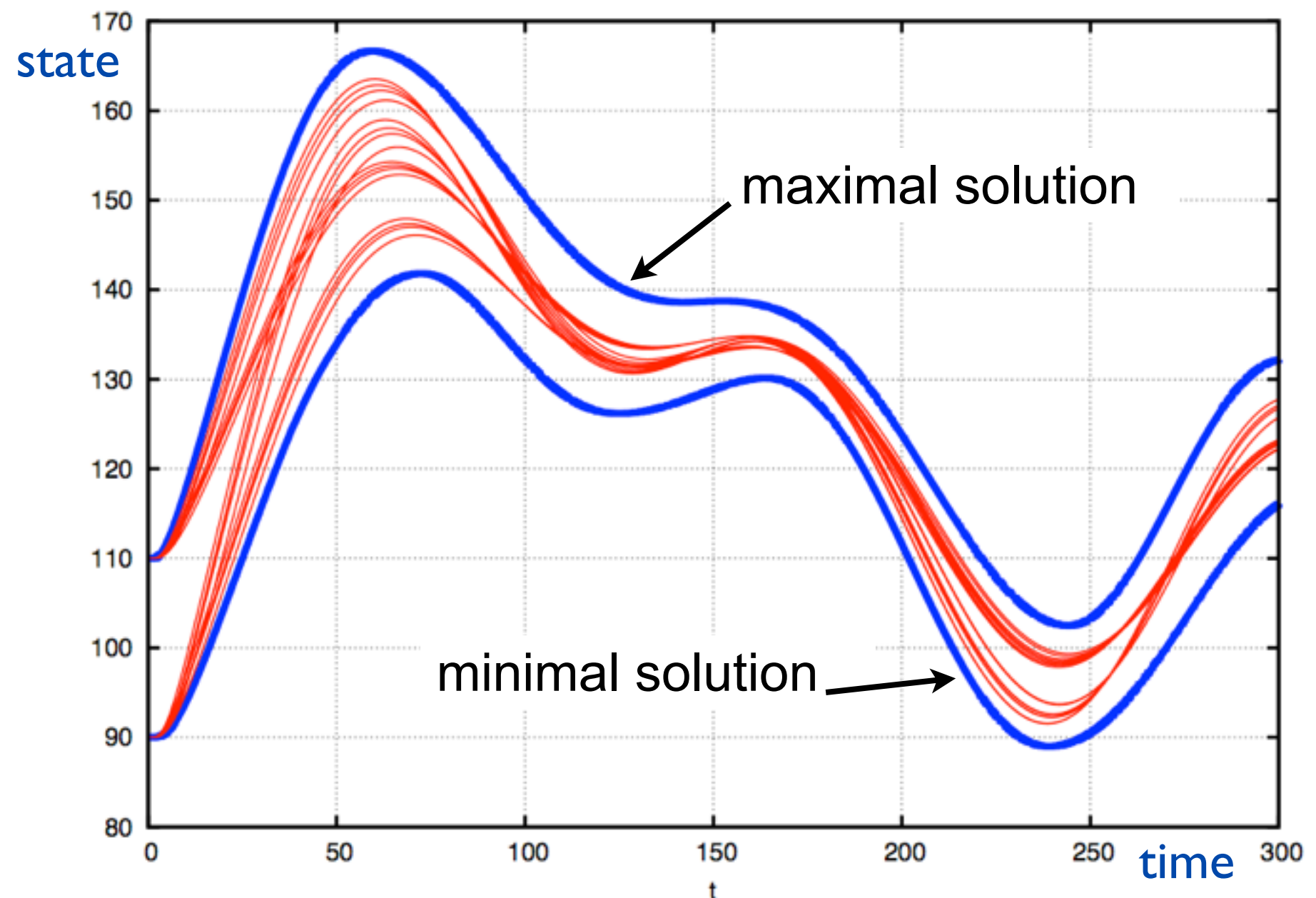
- (Ramdani, et al., IEEE Trans. Automatic Control 2009)



# Nonlinear Set Integration

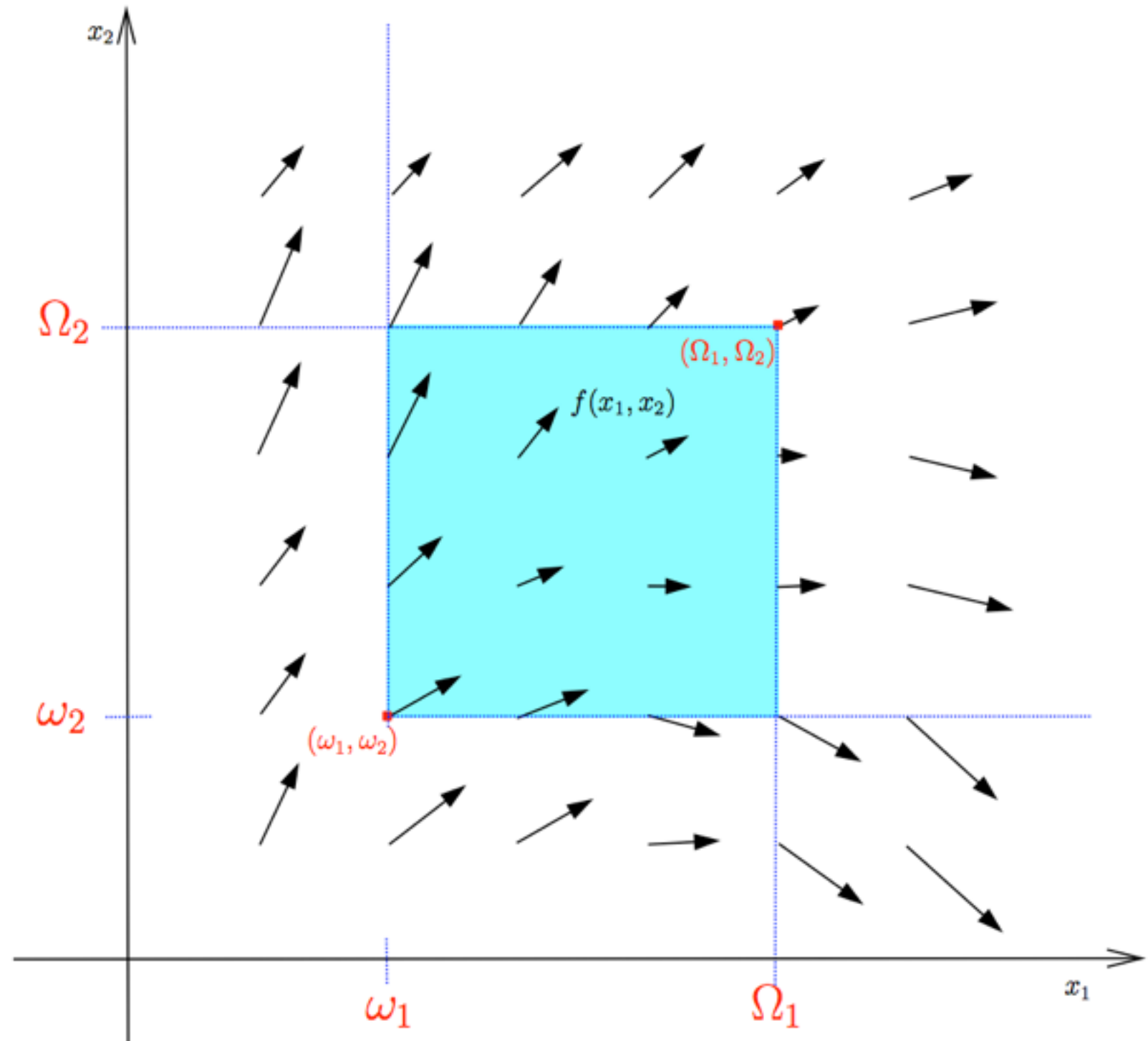
## ■ Comparison theorems for differential inequalities

### ● Bracketing systems

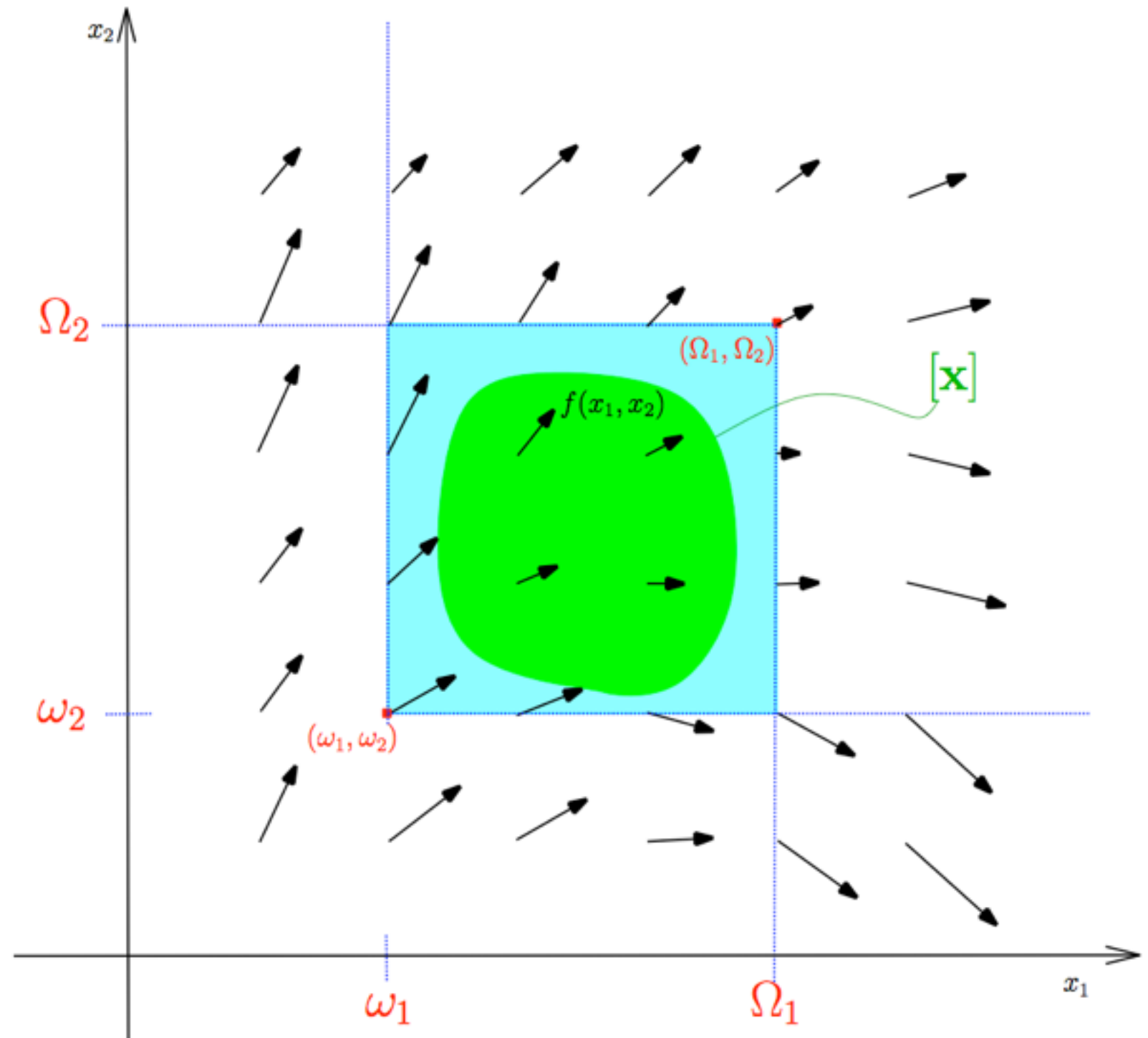




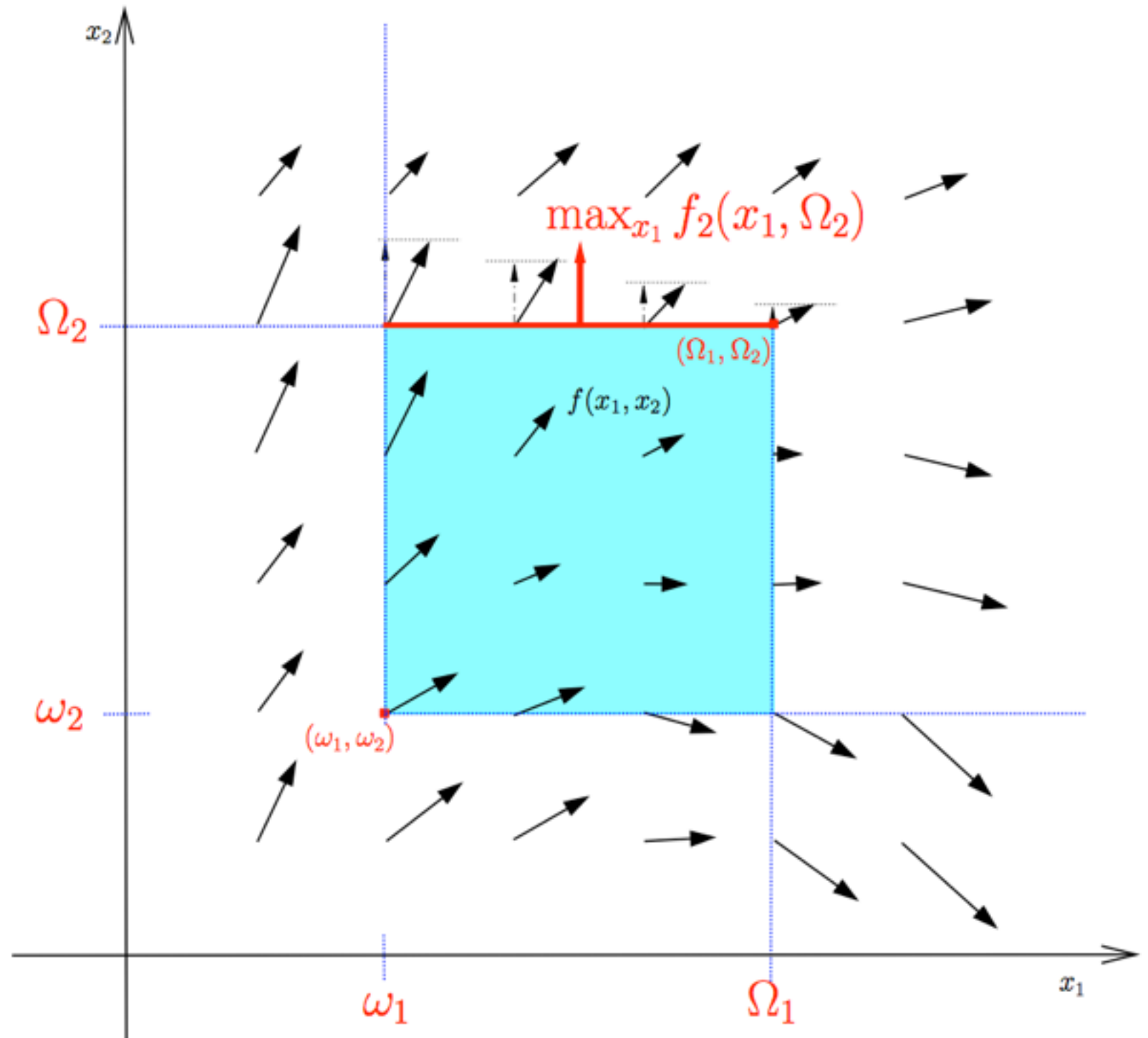
# Müller's theorem



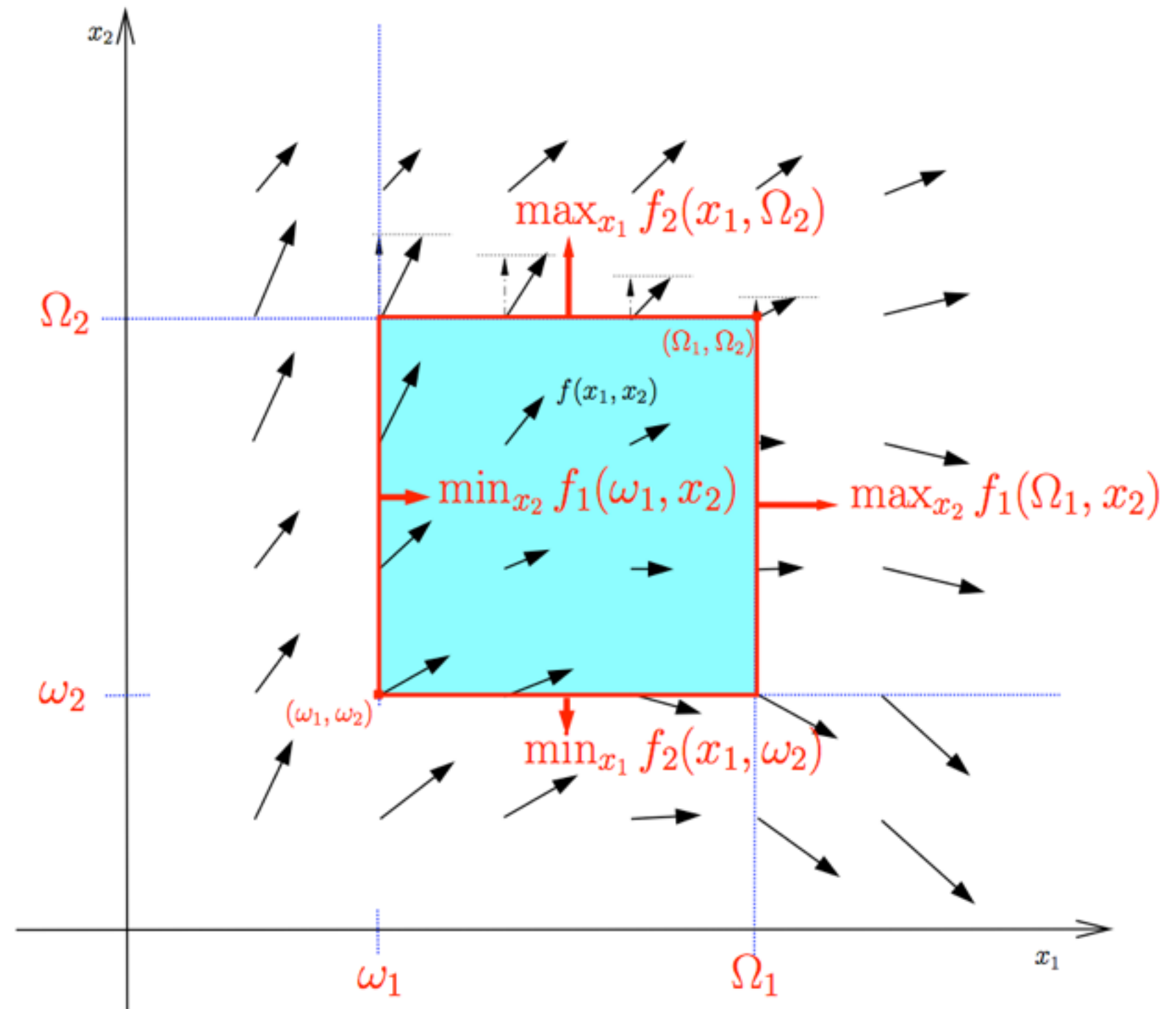
# Müller's theorem



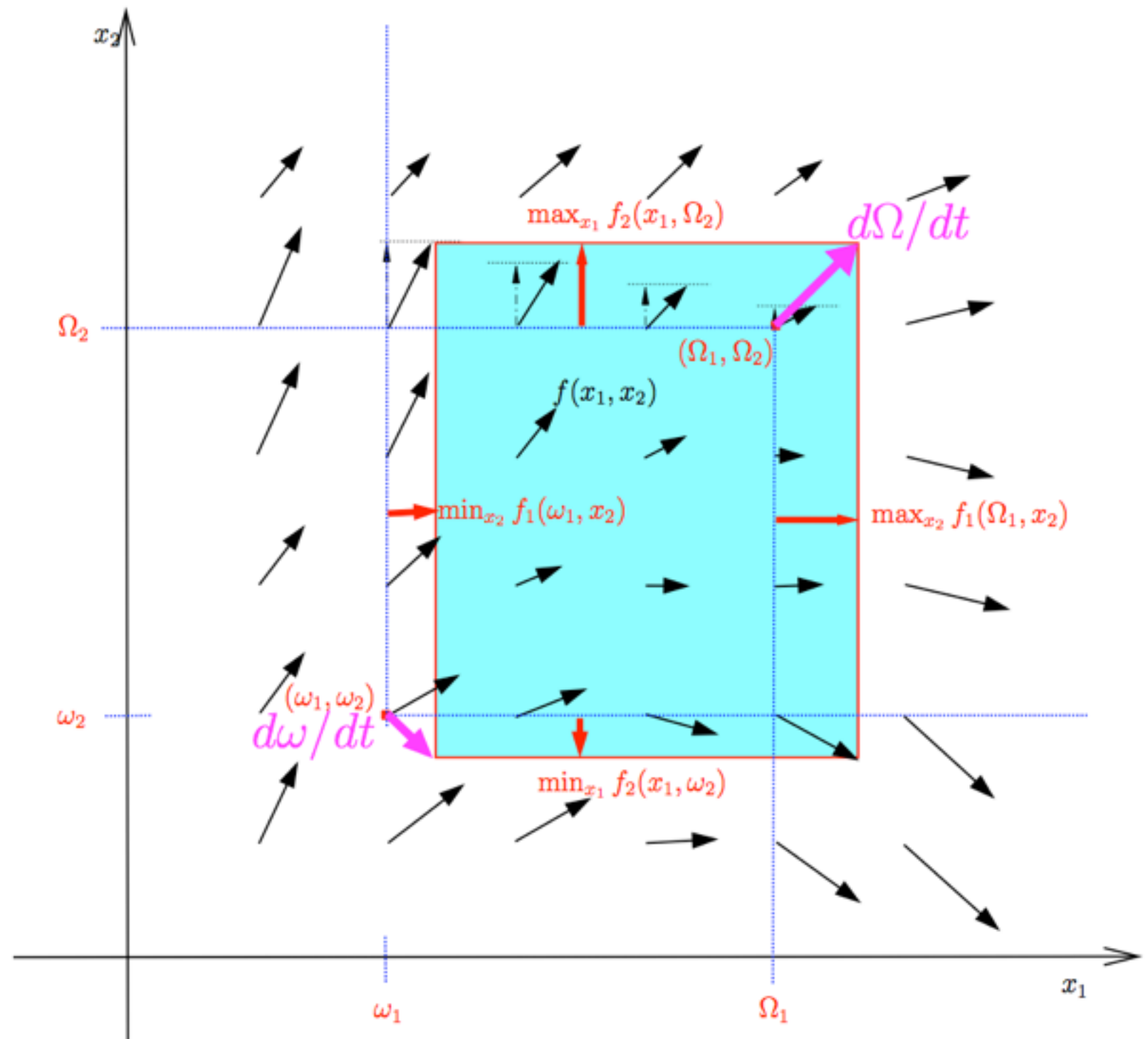
# Müller's theorem



# Müller's theorem



# Müller's theorem



## ■ Bracketing systems

### ● Dynamics of ...

$$\begin{cases} \dot{x}_1 = f_1(x_1, x_2, p, t), & x_1(t_0) \in [\underline{x}_{1,0}, \bar{x}_{1,0}] \subset \mathbb{R}, \\ \dot{x}_2 = f_2(x_1, x_2, p, t), & x_2(t_0) \in [\underline{x}_{2,0}, \bar{x}_{2,0}] \subset \mathbb{R}, \end{cases} \quad p \in [\underline{p}, \bar{p}] \quad t \geq t_0$$

If  $\forall t \geq t_0, \forall \mathbf{x}(t) \in [\omega(t), \Omega(t)] \subset \mathbb{R}^2, \forall p \in [\underline{p}, \bar{p}]$ ,

$$\frac{\partial f_1}{\partial x_2} > 0 \quad \wedge \quad \frac{\partial f_1}{\partial p} > 0$$

then  $f_1(\omega_1, \omega_2, \underline{p}) \leq f_1(\omega_1, x_2, p, t)$  and  $f_1(\Omega_1, x_2, p, t) \leq f_1(\Omega_1, \Omega_2, \bar{p})$   
 $\dot{\omega}_1(t) \equiv f_1(\omega_1, \omega_2, \underline{p})$  and  $f_1(\Omega_1, \Omega_2, \bar{p}) \equiv \dot{\Omega}_1(t)$



## ■ Bracketing systems

### ● Dynamics of ...

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If  $\forall t \geq t_0, \forall \mathbf{x}(t) \in [\omega(t), \Omega(t)] \subset \mathbb{R}^2, \forall p \in [\underline{p}, \bar{p}]$ ,

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$$\dot{\omega}_1(t) \equiv f_1(\omega_1, \omega_2, \underline{p})$$

$$\text{and } f_1(\Omega_1, \Omega_2, \bar{p}) \equiv \dot{\Omega}_1(t)$$

## ■ Comparison theorems for differential inequalities

- Müller's existence theorem (1936)

$$\text{If } \left\{ \begin{array}{l} \forall t \in [t_0, t_N], \forall \mathbf{x} \in \mathbb{D}, \forall \mathbf{p} \in \mathbb{P} \\ \forall i \min_{x_i = \omega_i} f_i(\mathbf{x}, \mathbf{p}, t) \geq D^\pm \omega_i(t) \\ \forall i \max_{x_i = \Omega_i} f_i(\mathbf{x}, \mathbf{p}, t) \leq D^\pm \Omega_i(t) \\ \omega(t_0) \leq \mathbf{x}(t_0) \leq \Omega(t_0) \end{array} \right. \Rightarrow \left\{ \begin{array}{l} \text{solution exists} \\ \forall t \in [t_0, t_N], \\ \omega(t) \leq \mathbf{x}(t) \leq \Omega(t) \end{array} \right.$$

- **Bracketing systems : coupled EDOs**

$$\Rightarrow \left\{ \begin{array}{l} \dot{\omega}(t) = \underline{f}(\omega, \Omega, \underline{\mathbf{p}}, \bar{\mathbf{p}}, t), \quad \omega(t_0) = \underline{\mathbf{x}}_0 \\ \dot{\Omega}(t) = \bar{f}(\omega, \Omega, \underline{\mathbf{p}}, \bar{\mathbf{p}}, t), \quad \Omega(t_0) = \bar{\mathbf{x}}_0 \end{array} \right.$$

## ■ Bracketing systems

- Example : Mitogen- Activated Protein Kinase (Sontag, 2005)

## ■ Bracketing systems

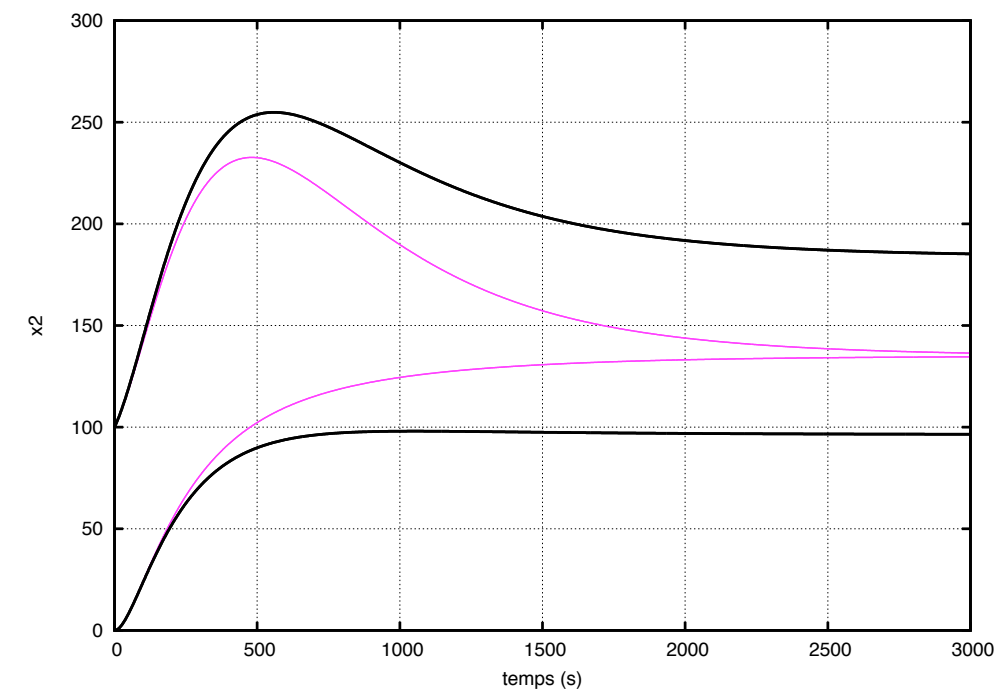
- Example : Mitogen- Activated Protein Kinase (Sontag, 2005)

$$\left\{ \begin{array}{lcl} \dot{x}_1 & = & -\frac{v_2 x_1}{k_2 + x_1} + v_0 u + v_1 \\ \dot{x}_2 & = & \frac{v_6 (y_{tot} - x_2 - x_3)}{k_6 + (y_{tot} - x_2 - x_3)} - \frac{v_3 x_1 x_2}{k_3 + x_2} \\ \dot{x}_3 & = & \frac{v_4 x_1 (y_{tot} - x_2 - x_3)}{k_4 + (y_{tot} - x_2 - x_3)} - \frac{v_5 x_3}{k_5 + x_3} \\ \dot{x}_4 & = & \frac{v_{10} (z_{tot} - x_4 - x_5)}{k_{10} + (z_{tot} - x_4 - x_5)} - \frac{v_7 x_3 x_4}{k_7 + x_4} \\ \dot{x}_5 & = & \frac{v_8 x_3 (z_{tot} - x_4 - x_5)}{k_8 + (z_{tot} - x_4 - x_5)} - \frac{v_9 x_5}{k_9 + x_5} \\ u & = & g x_5 \end{array} \right.$$

## ■ Bracketing systems

- Example : Mitogen- Activated Protein Kinase (Sontag, 2005)

$$\left\{ \begin{array}{lcl}
 \dot{x}_1 & = & -\frac{\bar{v}_2 x_1}{k_2 + x_1} + v_0 u + v_1 \\
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 \dot{x}_5 & = & \frac{v_8 x_3 (z_{tot} - \bar{x}_4 - x_5)}{k_8 + (z_{tot} - \bar{x}_4 - x_5)} - \frac{\bar{v}_9 x_5}{k_9 + x_5} \\
 \dot{\bar{x}}_1 & = & -\frac{\bar{v}_2 \bar{x}_1}{k_2 + \bar{x}_1} + \bar{v}_0 \bar{u} + \bar{v}_1 \\
 \dot{\bar{x}}_2 & = & \frac{\bar{v}_6 (\bar{y}_{tot} - \bar{x}_2 - \bar{x}_3)}{k_6 + (\bar{y}_{tot} - \bar{x}_2 - \bar{x}_3)} - \frac{\bar{v}_3 \bar{x}_1 \bar{x}_2}{k_3 + \bar{x}_2} \\
 \dot{\bar{x}}_3 & = & \frac{\bar{v}_4 \bar{x}_1 (\bar{y}_{tot} - \bar{x}_2 - \bar{x}_3)}{k_4 + (\bar{y}_{tot} - \bar{x}_2 - \bar{x}_3)} - \frac{\bar{v}_5 \bar{x}_3}{k_5 + \bar{x}_3} \\
 \dot{\bar{x}}_4 & = & \frac{\bar{v}_{10} (\bar{z}_{tot} - \bar{x}_4 - \bar{x}_5)}{k_{10} + (\bar{z}_{tot} - \bar{x}_4 - \bar{x}_5)} - \frac{\bar{v}_7 \bar{x}_3 \bar{x}_4}{k_7 + \bar{x}_4} \\
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 \frac{u}{\bar{u}} & = & \frac{g x_5}{g \bar{x}_5}
 \end{array} \right.$$

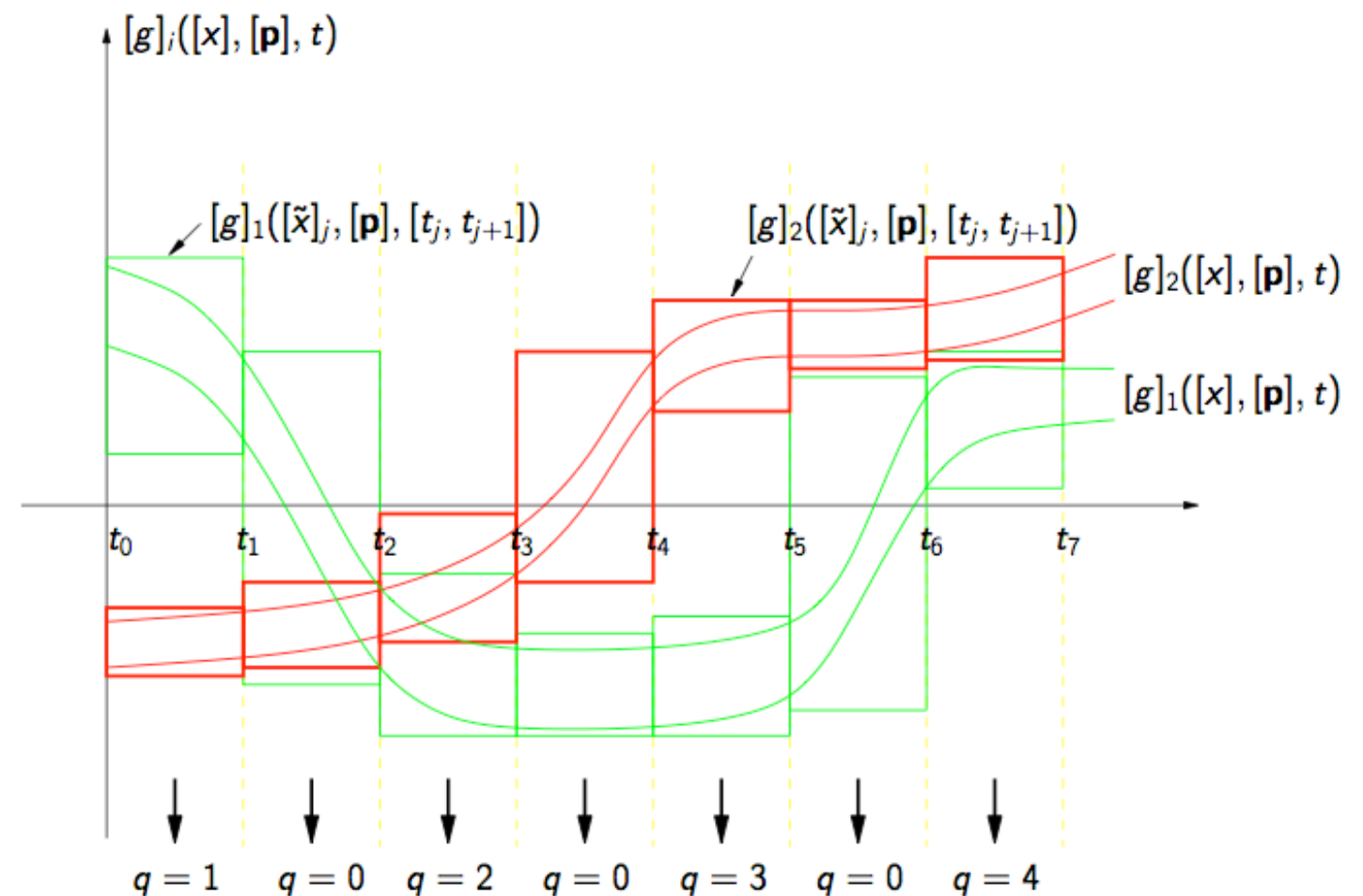


## ■ Nonlinear hybridization

- (Ramdani, et al., IEEE Trans. Automatic Control 2009)

$$\dot{x} = f(x, p_1, p_2, t) \quad x(t_0) \in [\underline{x}_0, \bar{x}_0] \subset \mathbb{R}, \quad p_i \in [\underline{p}_i, \bar{p}_i]$$

$$g_i(\cdot) = \frac{\partial f}{\partial p_i}(\cdot)$$



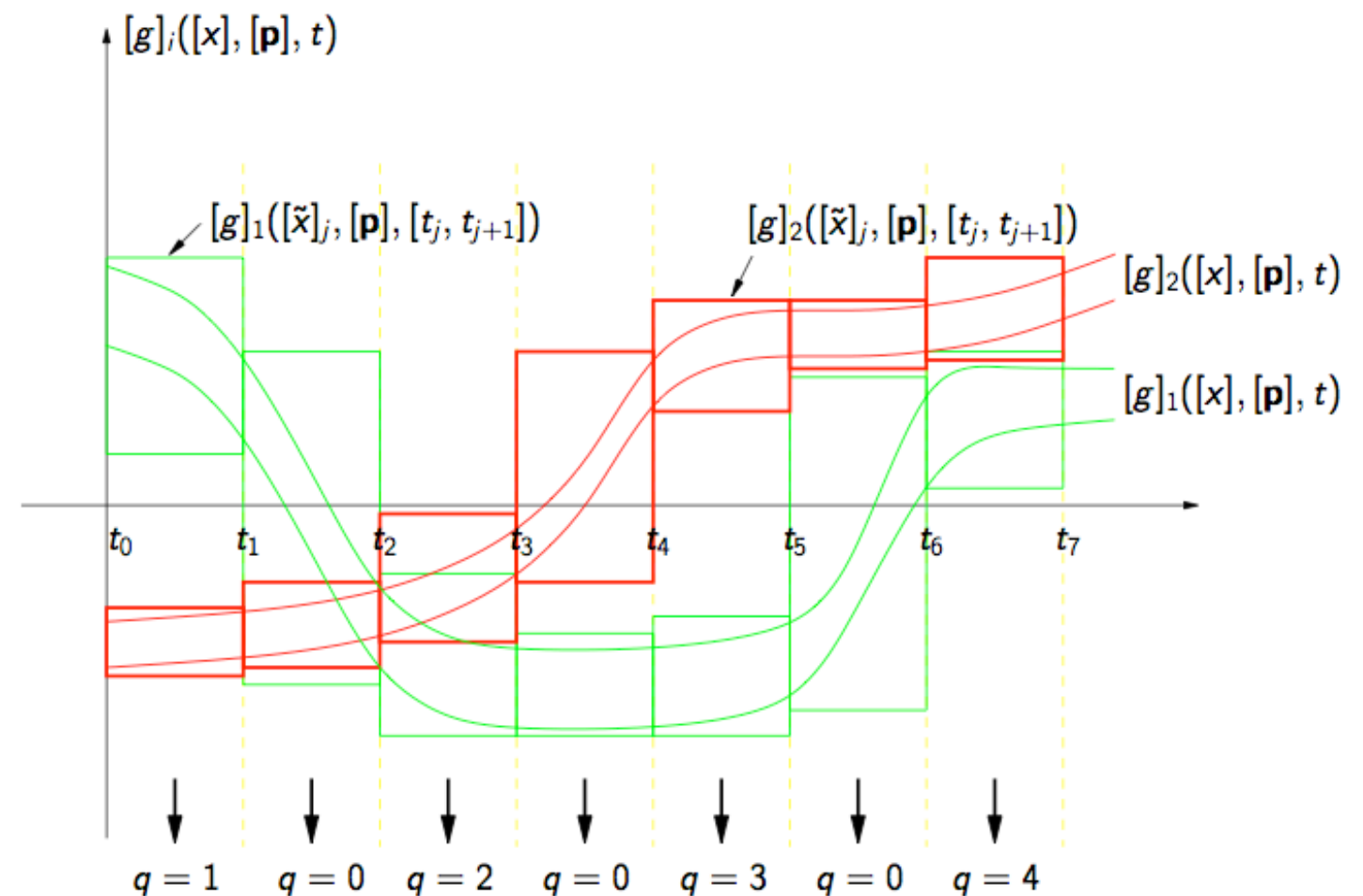


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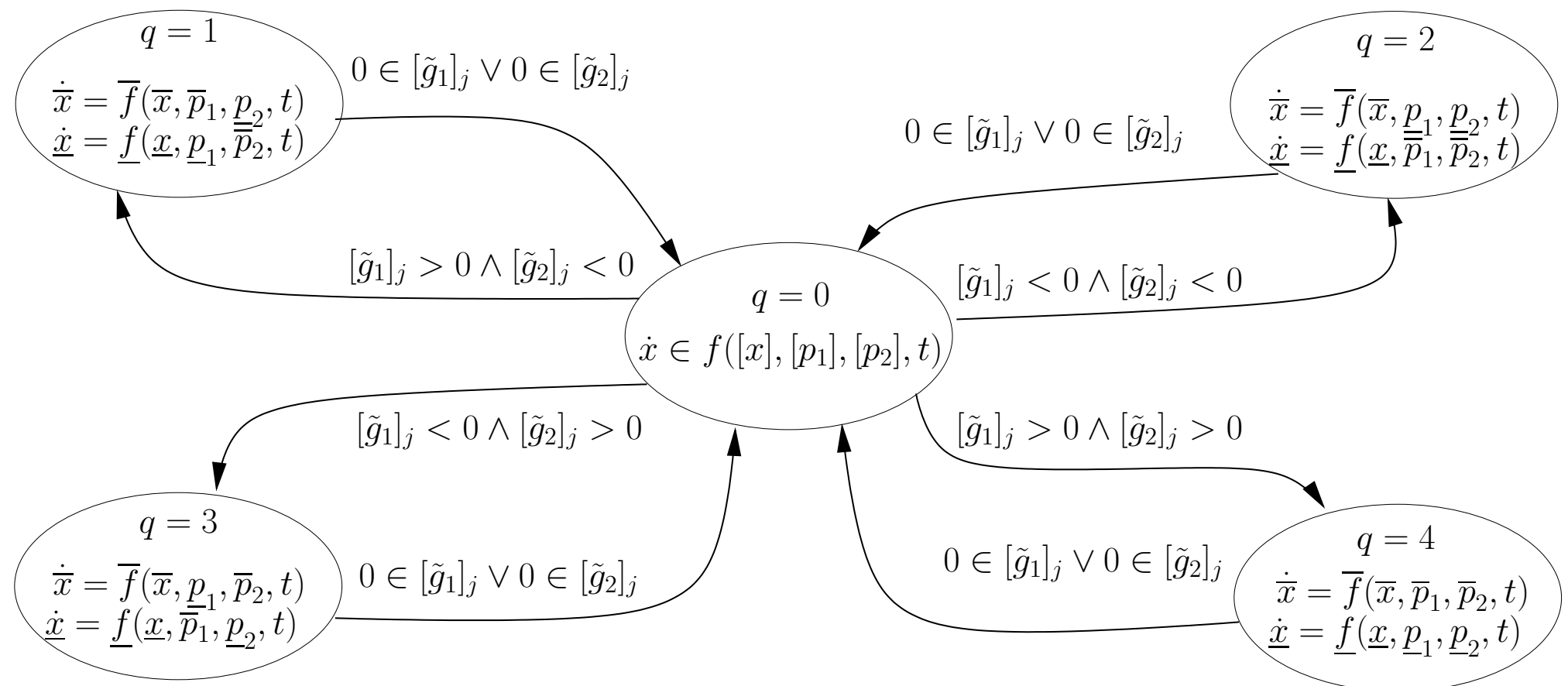
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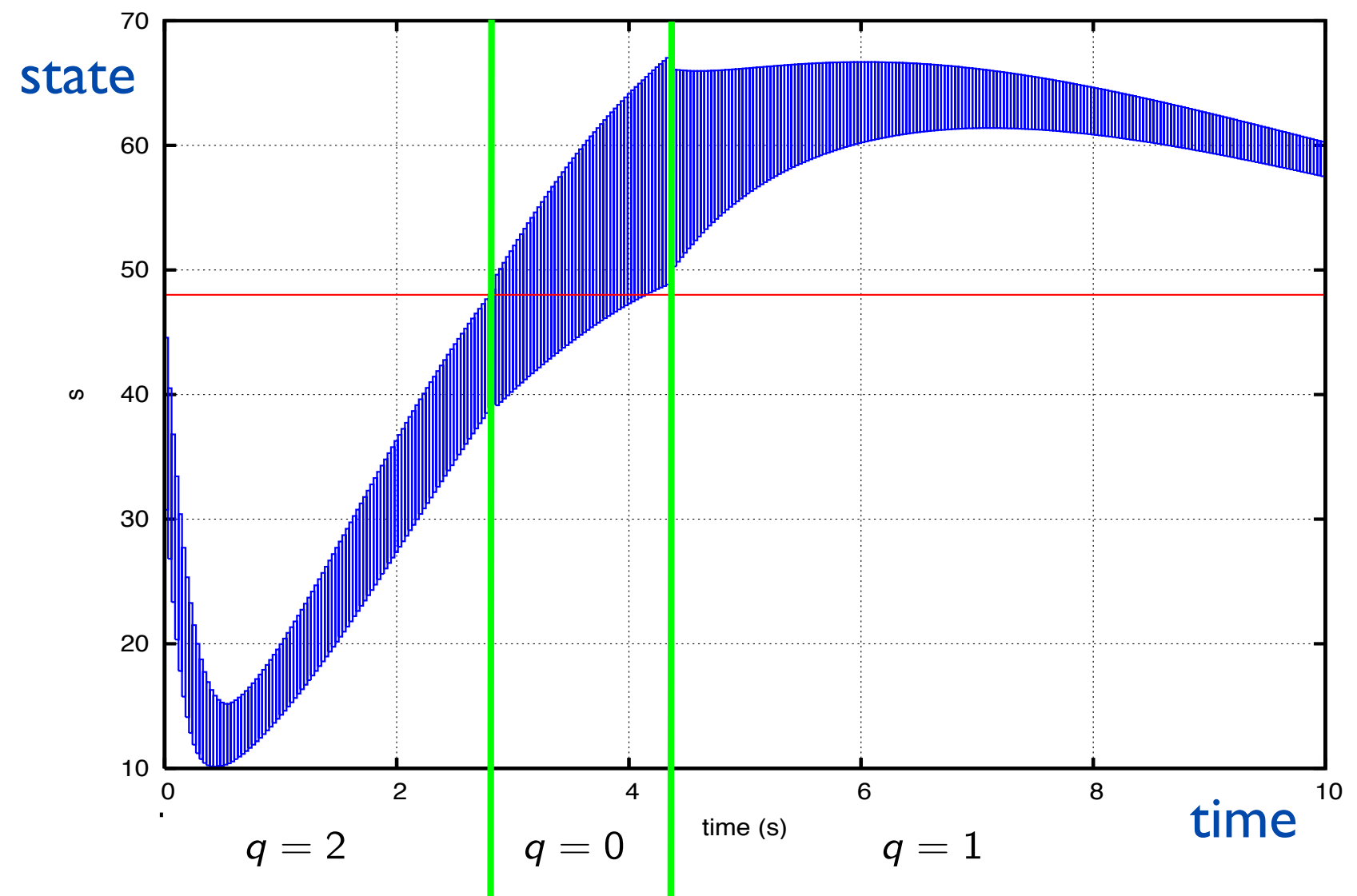
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# Nonlinear Set Integration

## ■ Nonlinear hybridization

- (Ramdani, et al., IEEE Trans. Automatic Control 2009)



## ■ Monotone order-preserving systems

- Müller, Kamke, Krasnoselskii, Hirsch, Smith, Angeli and Sontag.
- Preserve ordering on initial conditions.

$$\mathbf{x}(t_0) \prec \mathbf{y}(t_0) \Rightarrow \forall \mathbf{t} \geq \mathbf{t}_0 \quad \mathbf{x}(\mathbf{t}) \prec \mathbf{y}(\mathbf{t}) \quad \prec \in \{<, \leq, \geq, >\}$$

## ■ Monotone order-preserving systems

- Graphical test : monotone wrt orthant cones (Kunze & Siegel, 1999)

if  $\exists \mathbf{D} = \text{diag}[(-1)^{\varepsilon_1}, \dots, (-1)^{\varepsilon_n}], \varepsilon_i \in \{0, 1\}$

s.t  $\mathbf{x}(t, \mathbf{x}_0, t_0)$  and  $\mathbf{y}(t, \mathbf{y}_0, t_0)$  satisfy

$$\mathbf{D}\mathbf{y}_0 \geq \mathbf{D}\mathbf{x}_0 \Rightarrow \mathbf{D}\mathbf{y}(t, \mathbf{y}_0, t_0) \geq \mathbf{D}\mathbf{x}(t, \mathbf{x}_0, t_0) \quad \forall t \geq t_0.$$

## ■ Monotone order-preserving systems

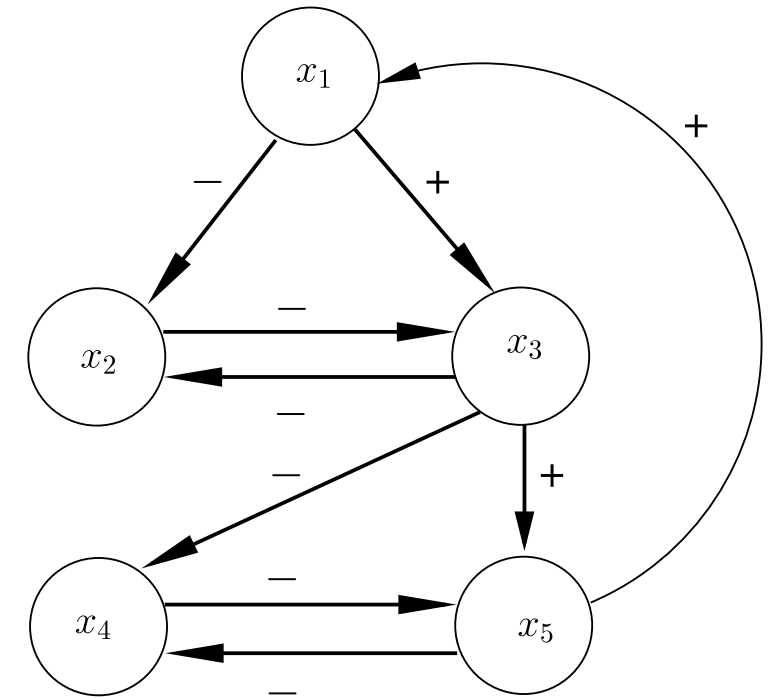
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$$\mathbf{D}\mathbf{y}_0 \geq \mathbf{D}\mathbf{x}_0 \Rightarrow \mathbf{D}\mathbf{y}(t, \mathbf{y}_0, t_0) \geq \mathbf{D}\mathbf{x}(t, \mathbf{x}_0, t_0) \quad \forall t \geq t_0.$$

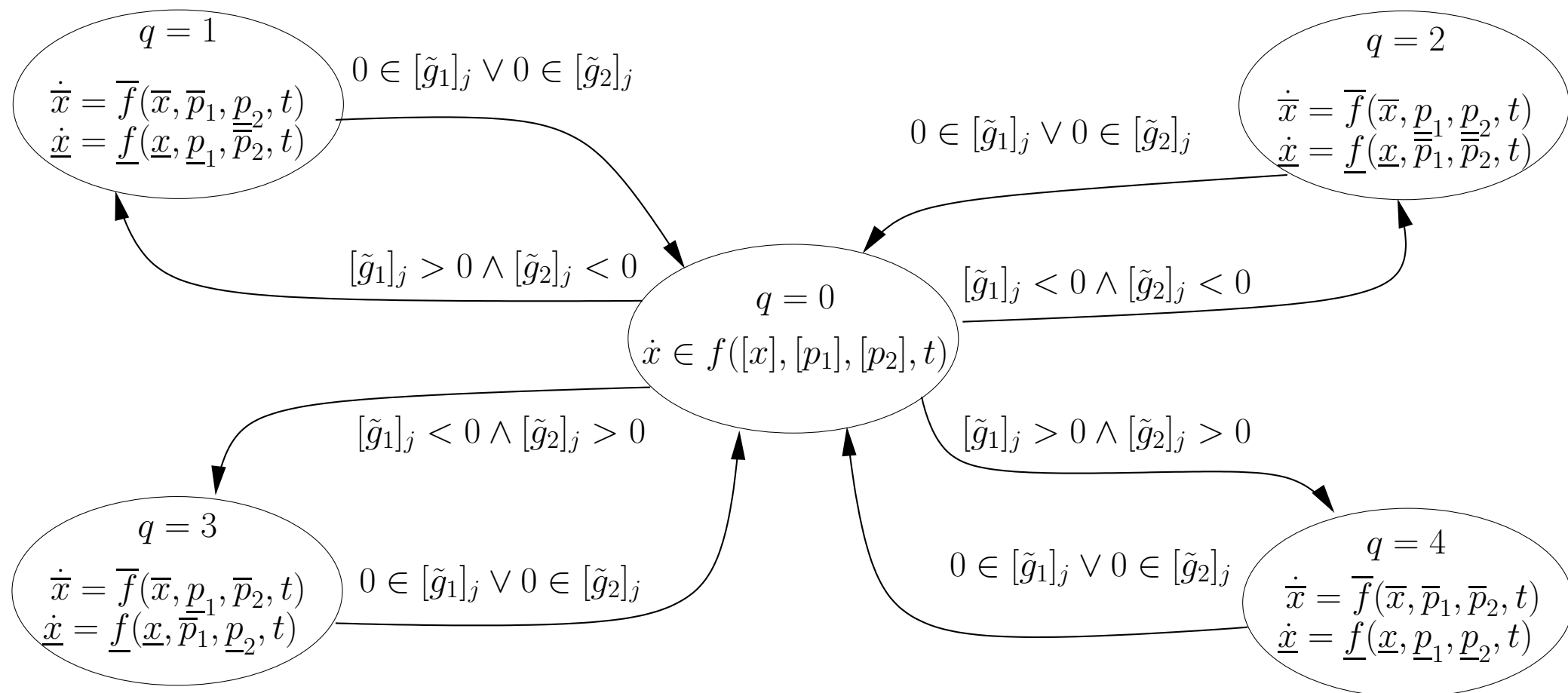
$$\begin{cases} \dot{x}_1 = -(v_2 x_1)/(k_2 + x_1) + v_0 g x_5 + v_1 \\ \dot{x}_2 = (v_6 (y_{tot} - x_2 - x_3))/(k_6 + (y_{tot} - x_2 - x_3)) - (v_3 x_1 x_2)/(k_3 + x_2) \\ \dot{x}_3 = (v_4 x_1 (y_{tot} - x_2 - x_3))/(k_4 + (y_{tot} - x_2 - x_3)) - (v_5 x_3)/(k_5 + x_3) \\ \dot{x}_4 = (v_{10} (z_{tot} - x_4 - x_5))/(k_{10} + (z_{tot} - x_4 - x_5)) - (v_7 x_3 x_4)/(k_7 + x_4) \\ \dot{x}_5 = (v_8 x_3 (z_{tot} - x_4 - x_5))/(k_8 + (z_{tot} - x_4 - x_5)) - (v_9 x_5)/(k_9 + x_5) \end{cases}$$





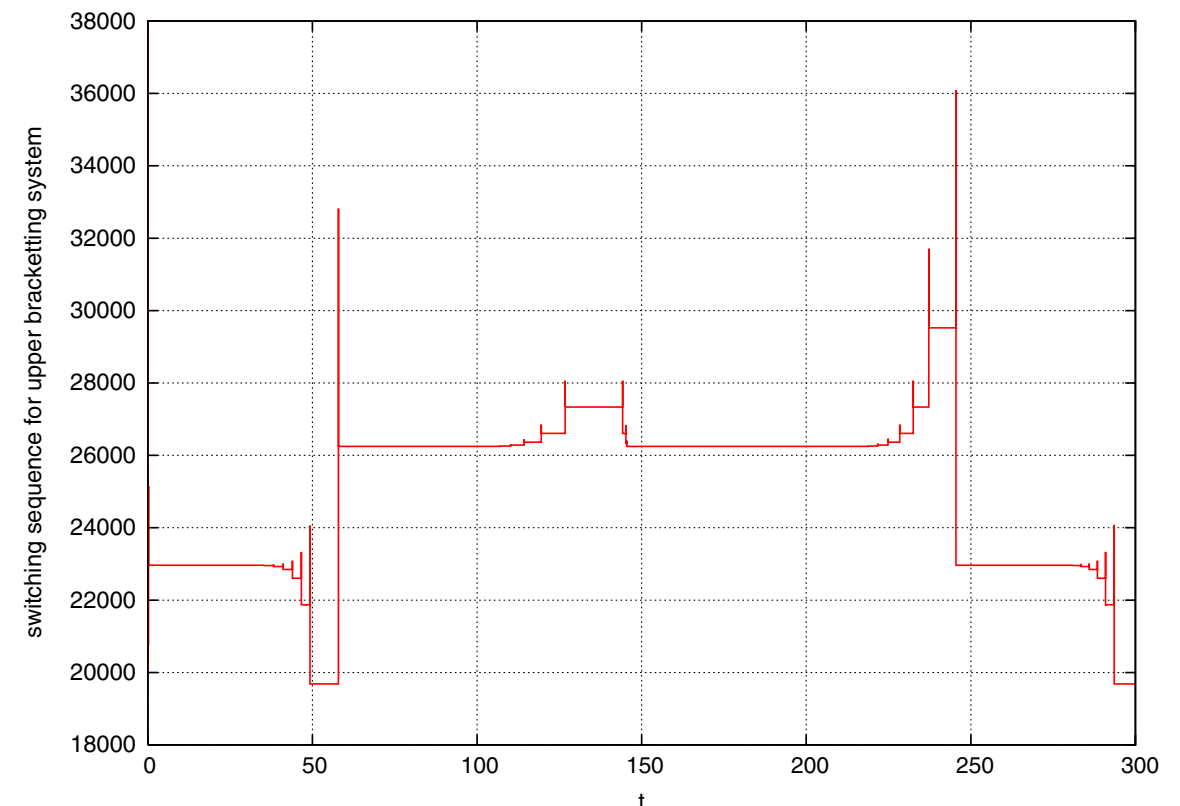
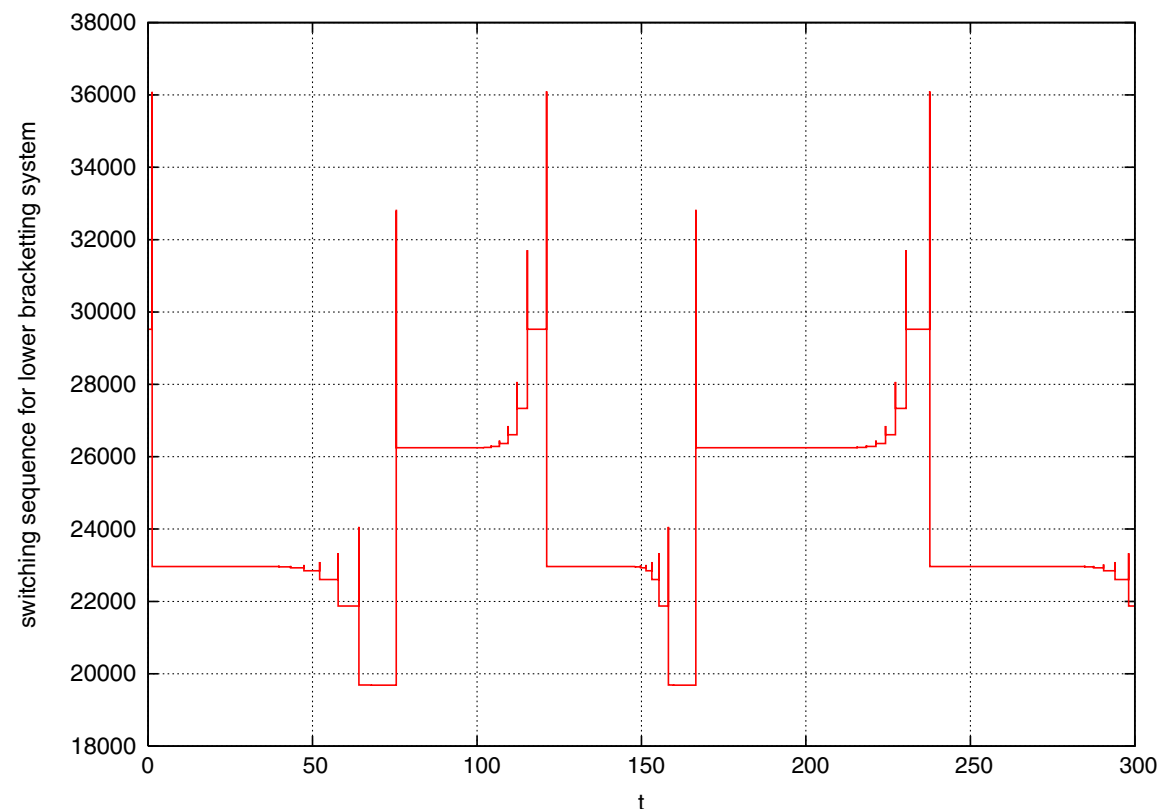
## ■ Nonlinear hybridization & Monotone systems

- (Ramdani, et al., Nonlinear Analysis Hybrid Systems 2010)



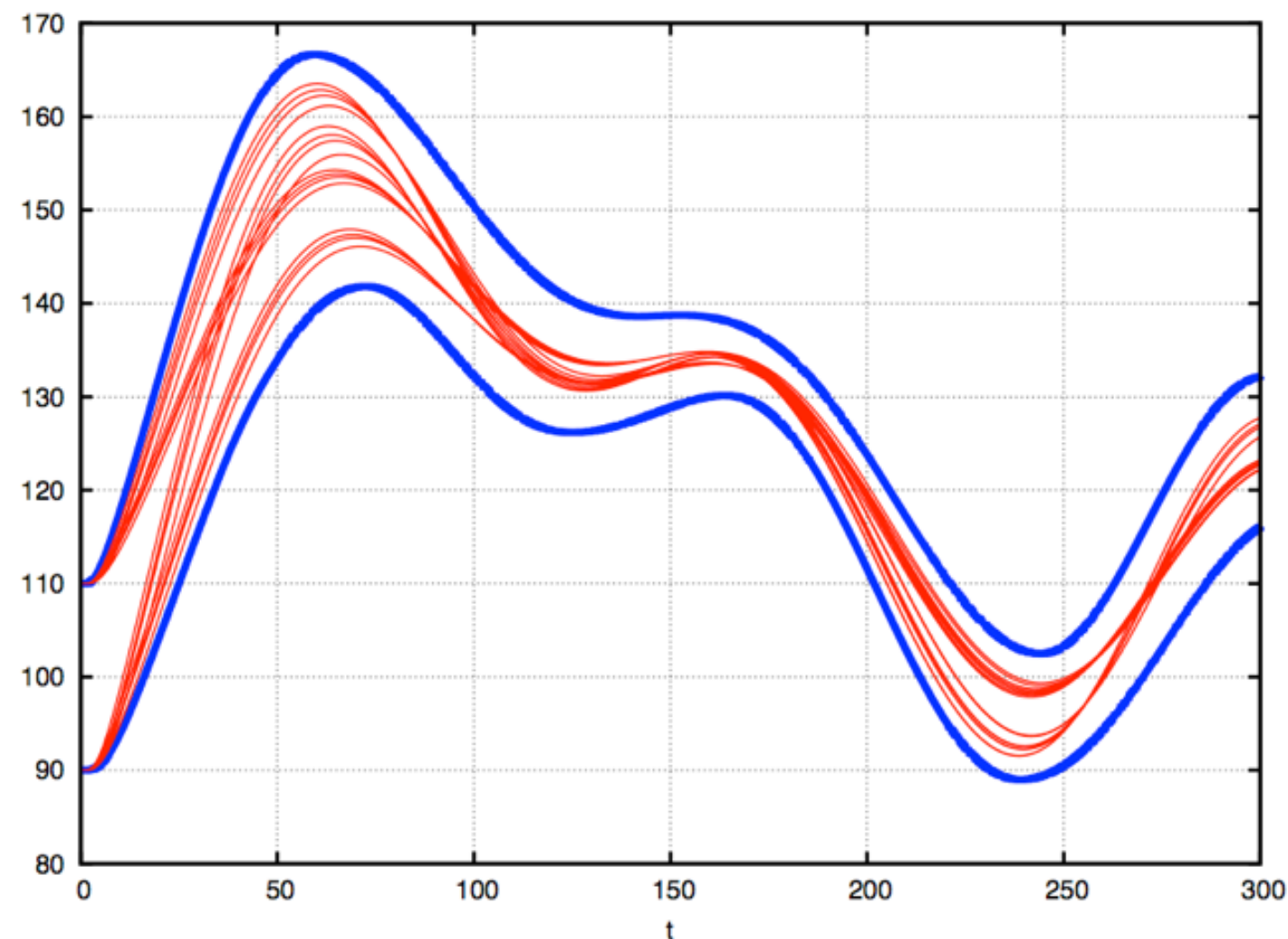
## ■ Nonlinear hybridization & Monotone systems

- (Ramdani, et al., Nonlinear Analysis Hybrid Systems 2010)
- **Non-coupled** bracketing systems



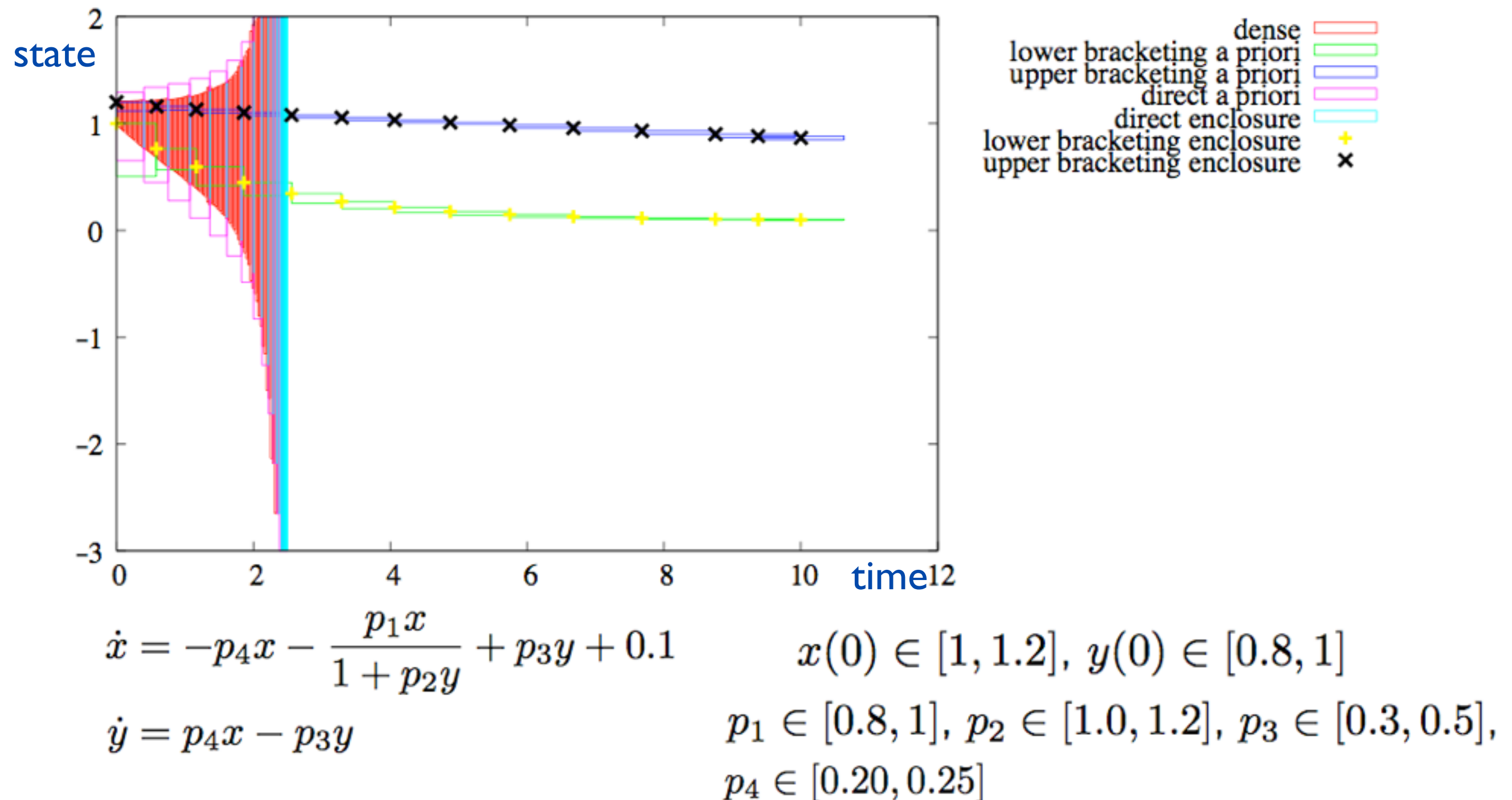
## ■ Nonlinear hybridization & Monotone systems

- (Ramdani, et al., Nonlinear Analysis Hybrid Systems 2010)
- Non-coupled bracketing systems



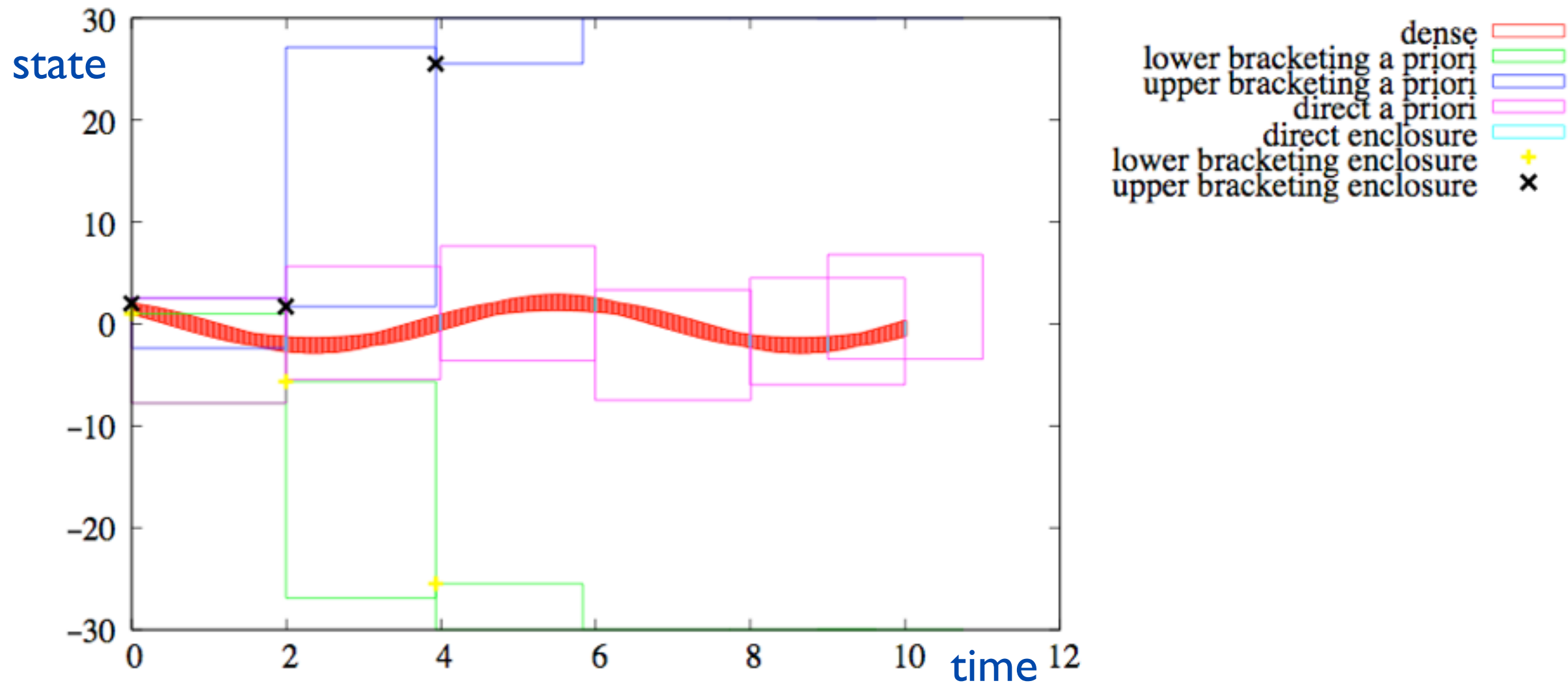
# Nonlinear Set Integration

## Interval Taylor methods vs Bracketing systems



# Nonlinear Set Integration

## Interval Taylor methods vs Bracketing systems



$$\dot{x} = y, \dot{y} = -x,$$

$$x(0), y(0) \in [1, 2]$$



## IOLAVABE: iSAT-ODE Layer Around VNODE-LP and Bracketing Enclosures

### About

The IOLAVABE library encapsulates the part of the iSAT-ODE tool that handles the generation of ODE enclosures using VNODE-LP and bracketing systems.

**IOLAVABE is made available here solely for scientific research.**

Detailed licensing information can be found in the LICENSE file inside the source code archive. IOLAVABE depends on and the archive file contains modified versions of [VNODE-LP](#) (itself including a copy of [FADBAD++](#)) and of [filib++](#). The unmodified versions can be found in the bundled archive as well. Please note the licensing information shipped with these and all indirectly or directly used libraries as well (you will find pointers to the respective terms of use in the INSTALL or LICENSE file or in your system's package management system).

Installation instructions are to be found in the INSTALL file, and a list of changes with respect to earlier releases can be found in the changelog file.

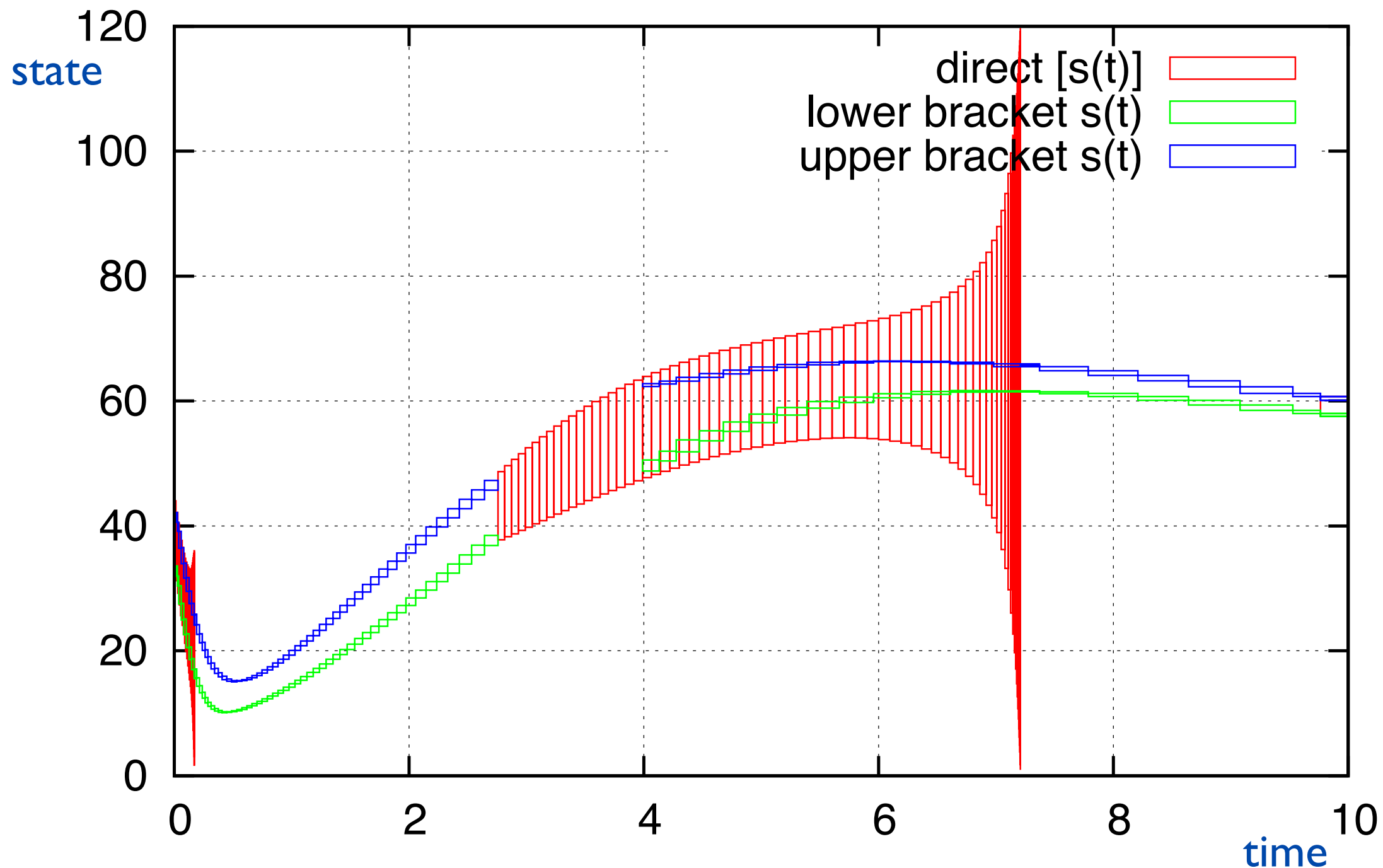
Contact the author: [Andreas Eggers](#)

<https://seshome.informatik.uni-oldenburg.de/eggers/iolavabe.php>



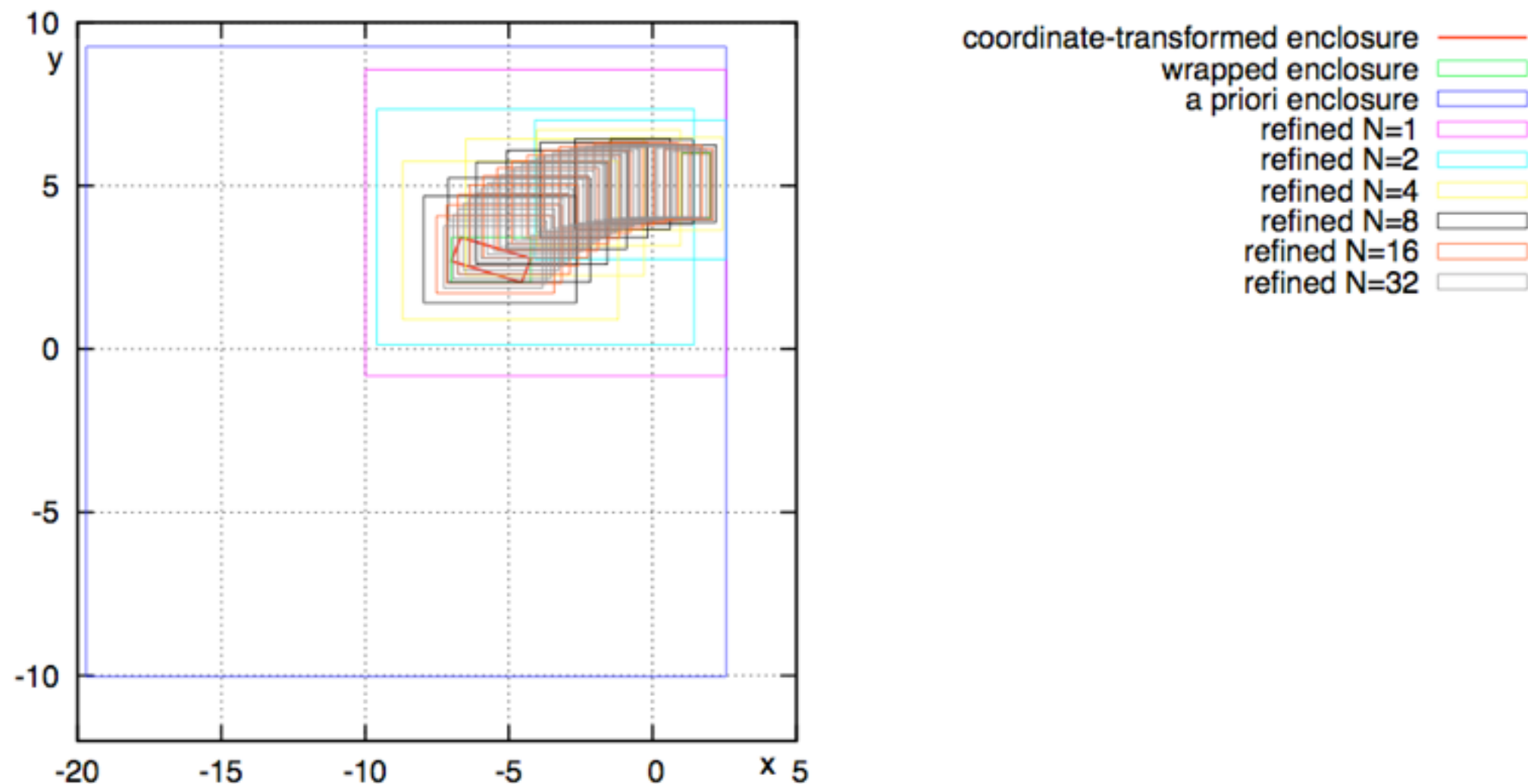
- generates **on-the-fly** hybrid bracketing systems, i.e. tries to re-start bracketing system when monotonicity changes
- uses subordinate local optimization to compute **signs of partial derivatives** on subranges to improve bracketing

## ■ Typical results: Taylor methods vs Bracketing systems



- harmonizes bracketing and direct enclosures,  
i.e. **synchronizes time step**,
- often **intersects enclosures** and reinitializes methods

- stores Taylor coefficients to **recompute «refined»** enclosures at intermediate steps.

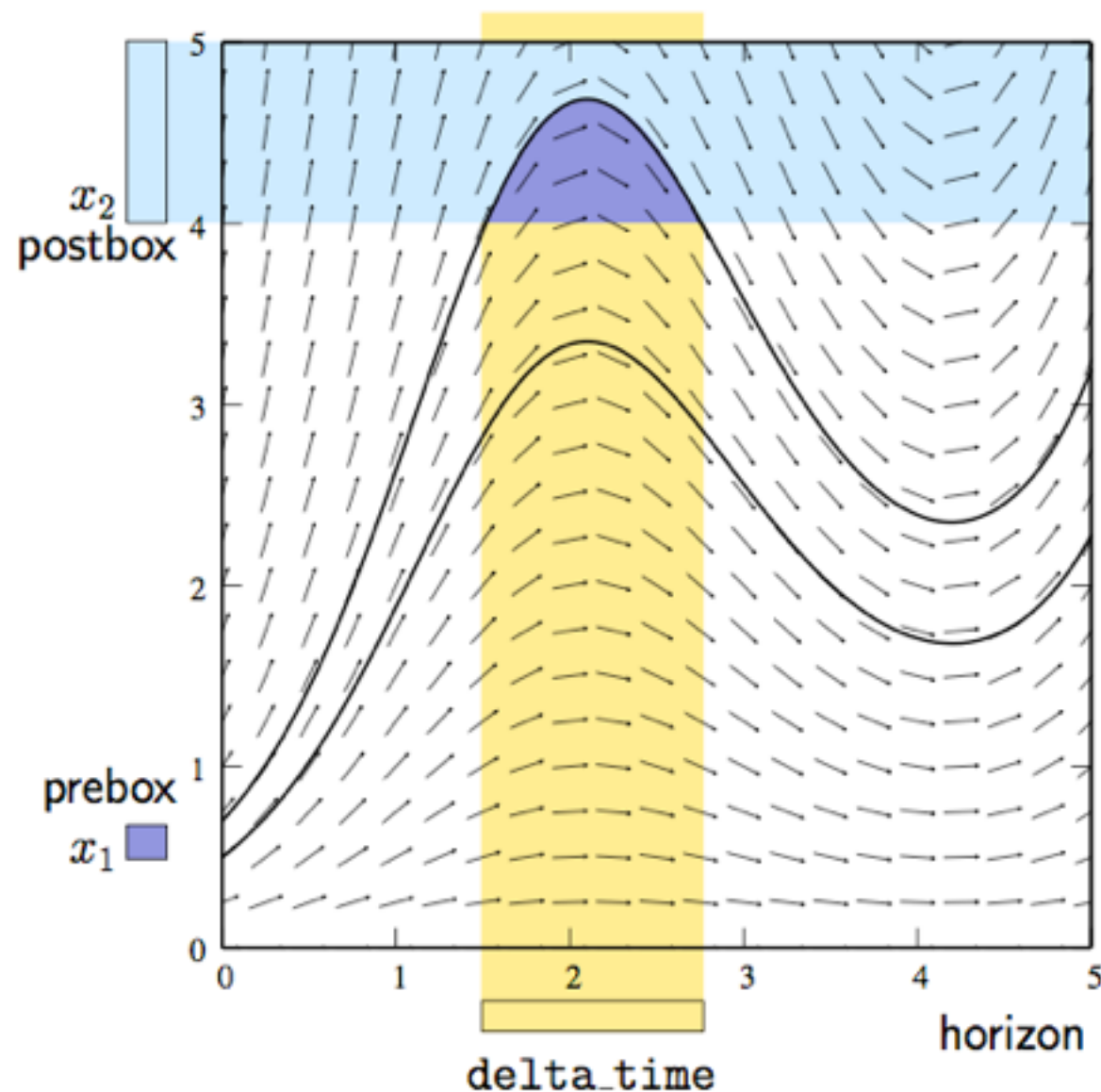


$$\dot{x} = -y, \quad \dot{y} = 0.6 \cdot x, \quad x_0 \in [1, 2], \quad y_0 \in [4, 6], \quad t_1 = 1.6$$

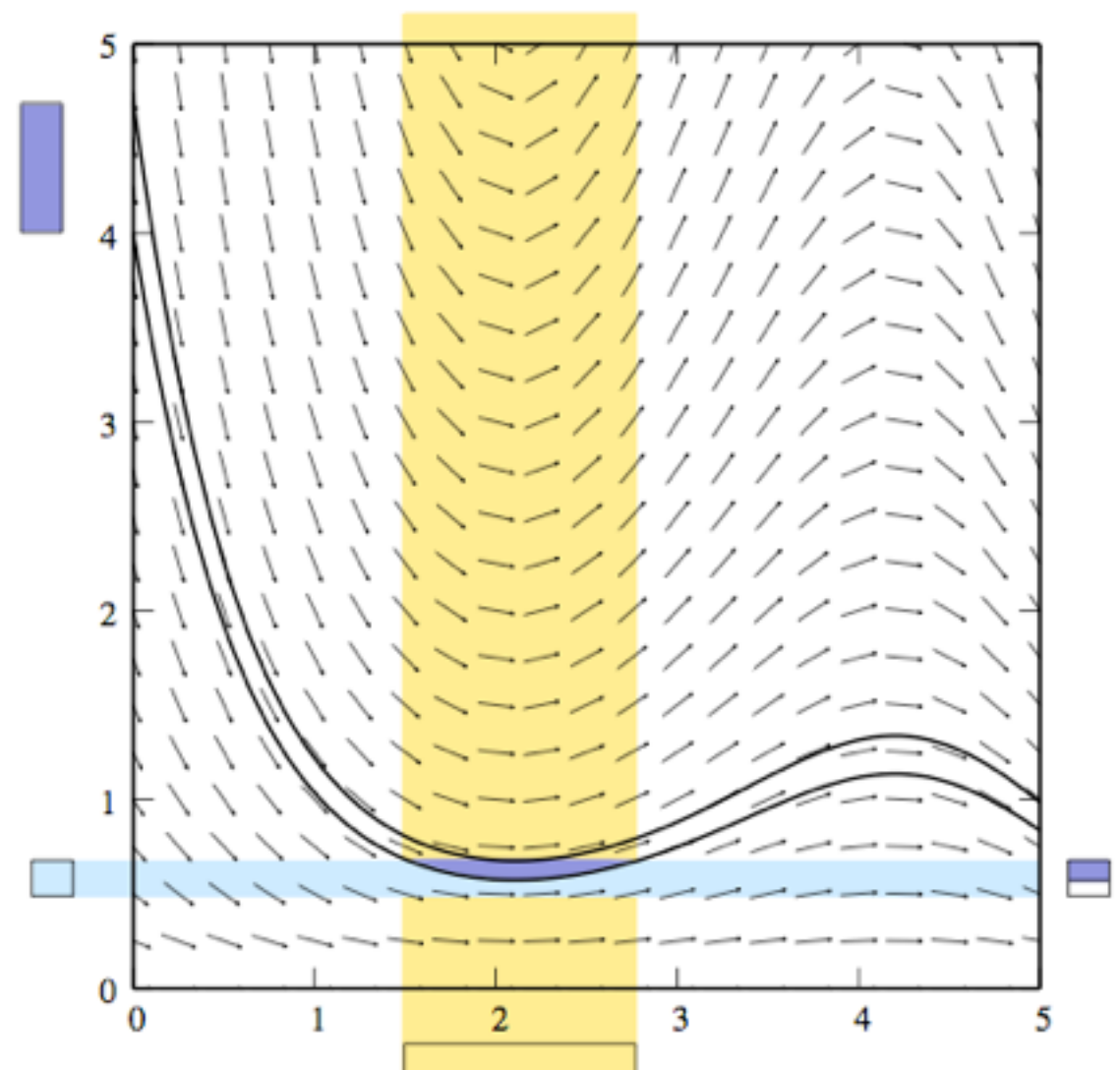
- detects independent group of ODEs
- detects when flow invariants are being left

- can **contract** pre- & post-box using forward and backward deductions

forward propagation



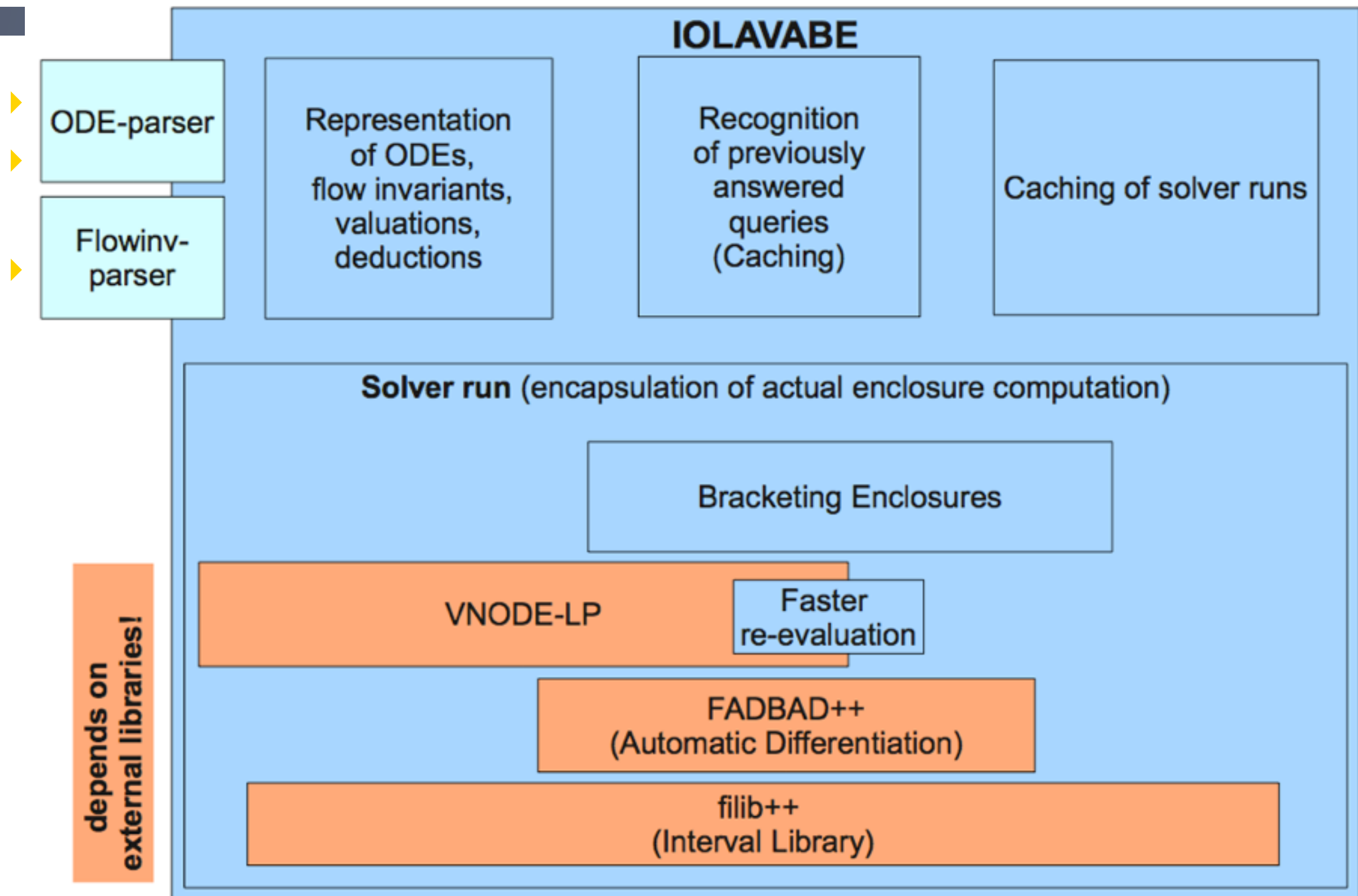
backward propagation





- algorithm's parameters are exposed to the outside
- parsers for ODEs and flow invariants offer string interface

# IOLAVABE Architecture Sketch



- IOLAVABE :  
the iSAT-ODE layer around VNODE-LP and  
bracketing enclosures
- gives a high-level interface for generating enclosures  
of ODE constraints
- Source code available for not-for-profit civilian  
scientific research : **try it !**

- Hybrid Reachability : Motivations
  - Analysis of complex dynamical systems
  - Reachability-based methods
  
- Nonlinear reachability
  - Interval Taylor methods
  - Bracketing enclosures
  - Software implementation

**All papers on**  
**<http://lune.bourges.univ-orleans.fr/ramdani>**

**Thank you !**

**Questions ?**