

# A Contractor for the Non-Overlapping of Objects Described by Non-Linear Inequalities

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The Non-Overlapping Constraint

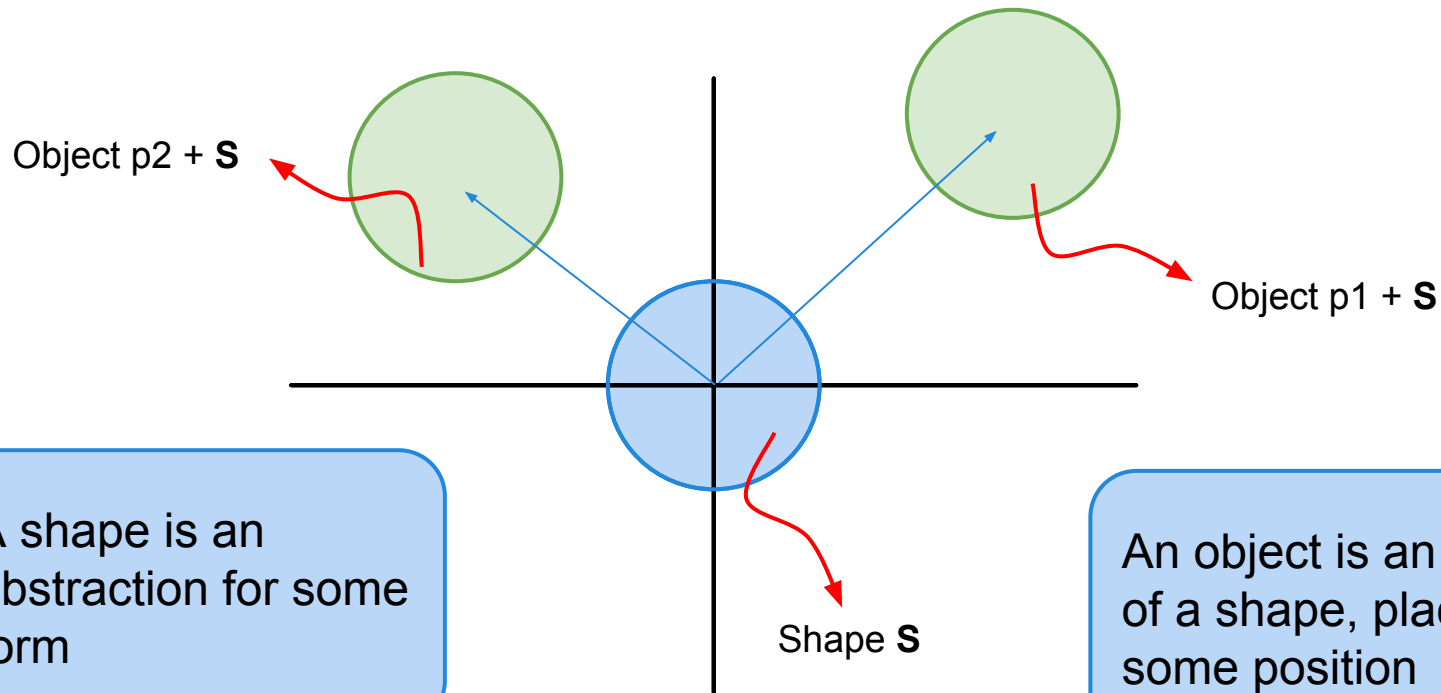
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# Definitions

## About Objects and Shapes

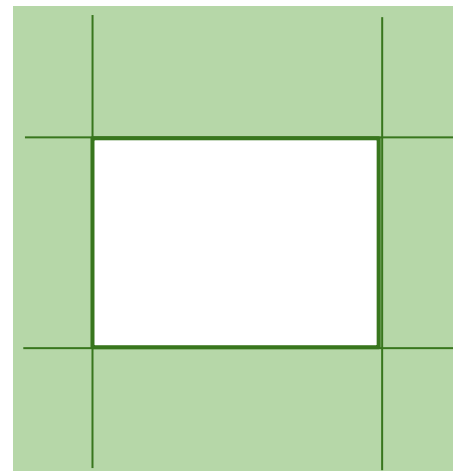
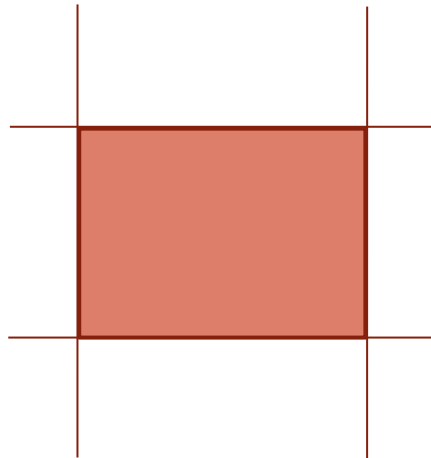
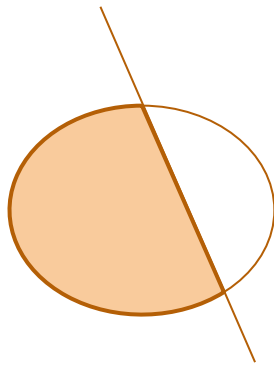
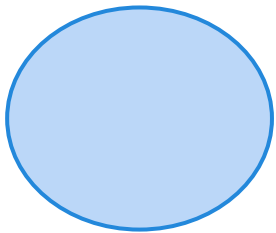


A shape is an abstraction for some form

An object is an instance of a shape, placed at some position

# Definitions

- Our target are the **curved shapes**
- That can be represented by **non-linear inequalities**
- In a general way with **disjunctions of system of non-linear inequalities**



# Definitions

System<sub>1</sub>

U

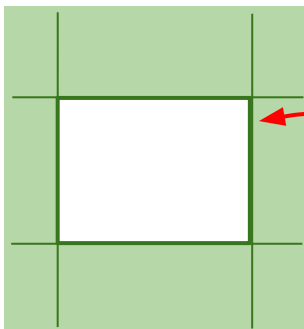
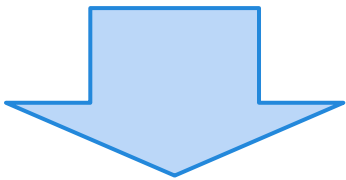
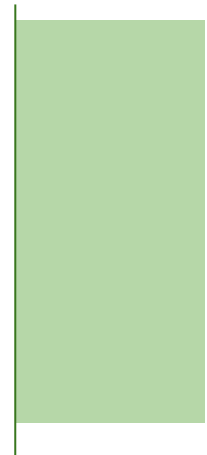
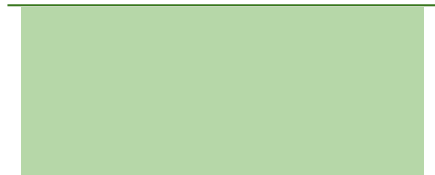
System<sub>2</sub>

U

System<sub>3</sub>

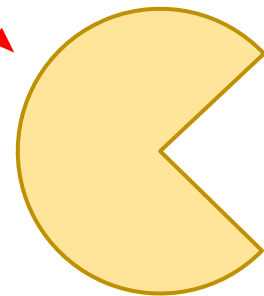
U

System<sub>4</sub>



The resulting shape is not necessarily convex

And this allows to represent the available space for packing as a regular shape (called "enclosing shape")

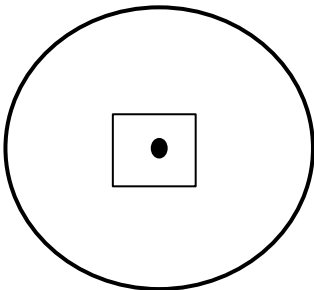


# Definitions

## Moving and Reference Objects

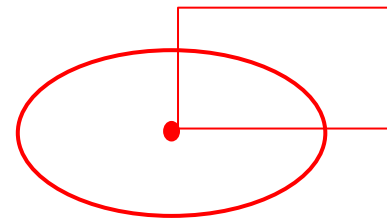
All the objects that we consider must take a **rol** at some moment

Reference Object



It is an object fixed, and it is used as reference for to contract the **origin box** of other object

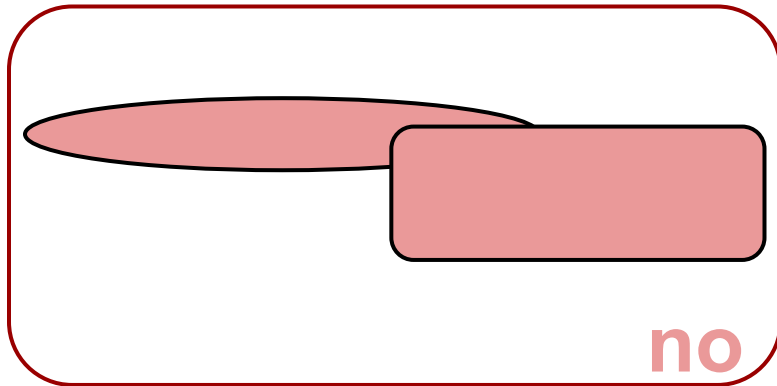
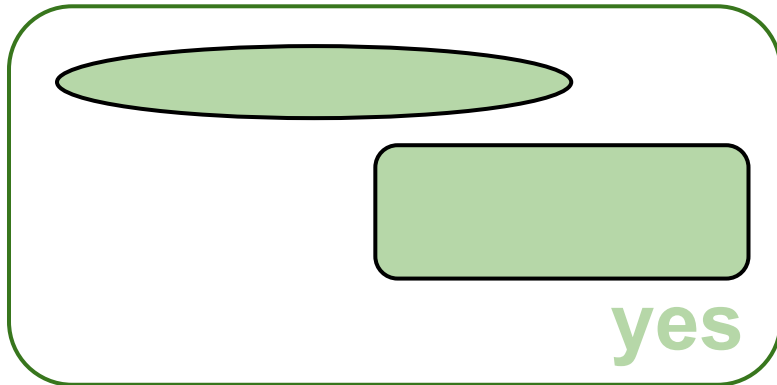
Moving Object



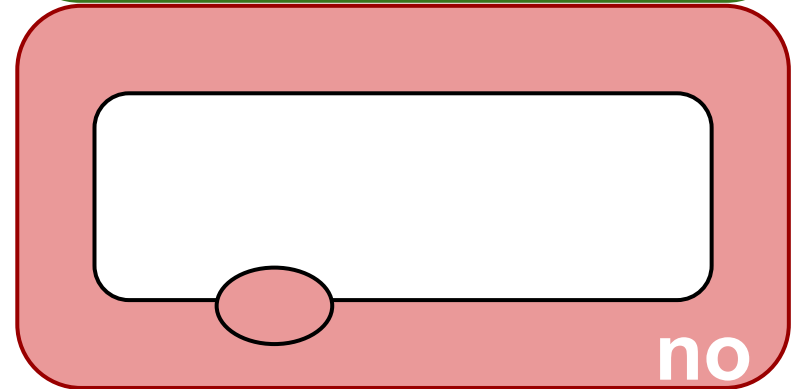
It is an object which **origin box** we want to contract

# The Non-Overlapping Constraint

Between Regular Shapes



Between an Enclosing Shape and a regular Shape



The overlapping with the enclosing shape is handled as with any other shape, since the description of the objects allows a transparent process.

# The Non-Overlapping Constraint

Domain of the  
positions of the  
object **a**

Domain of the  
positions of the  
object **b**

$$\forall p_a \in [\mathbf{O}_a], \forall p_b \in [\mathbf{O}_b],$$

$$(S_a + p_a) \cap (S_b + p_b) = \emptyset$$

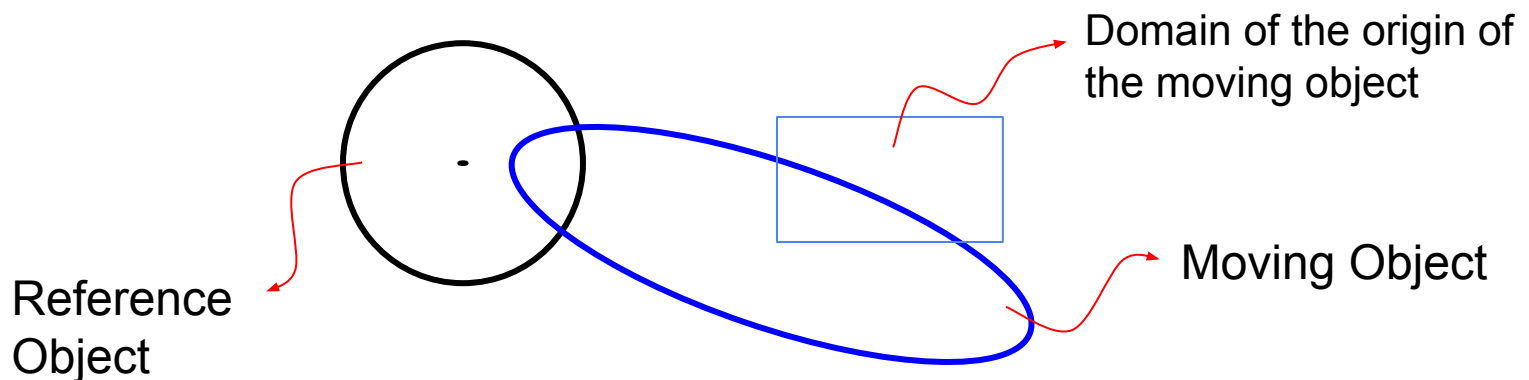
Shape **a**

Shape **b**



# The Contractor

- **Target:** Remove all the parts that violate the non-overlapping constraint
- The contraction operates over 2 objects
  - The reference object
  - The moving object



# The Contractor

CtcPacking

CtcShadowPropagator

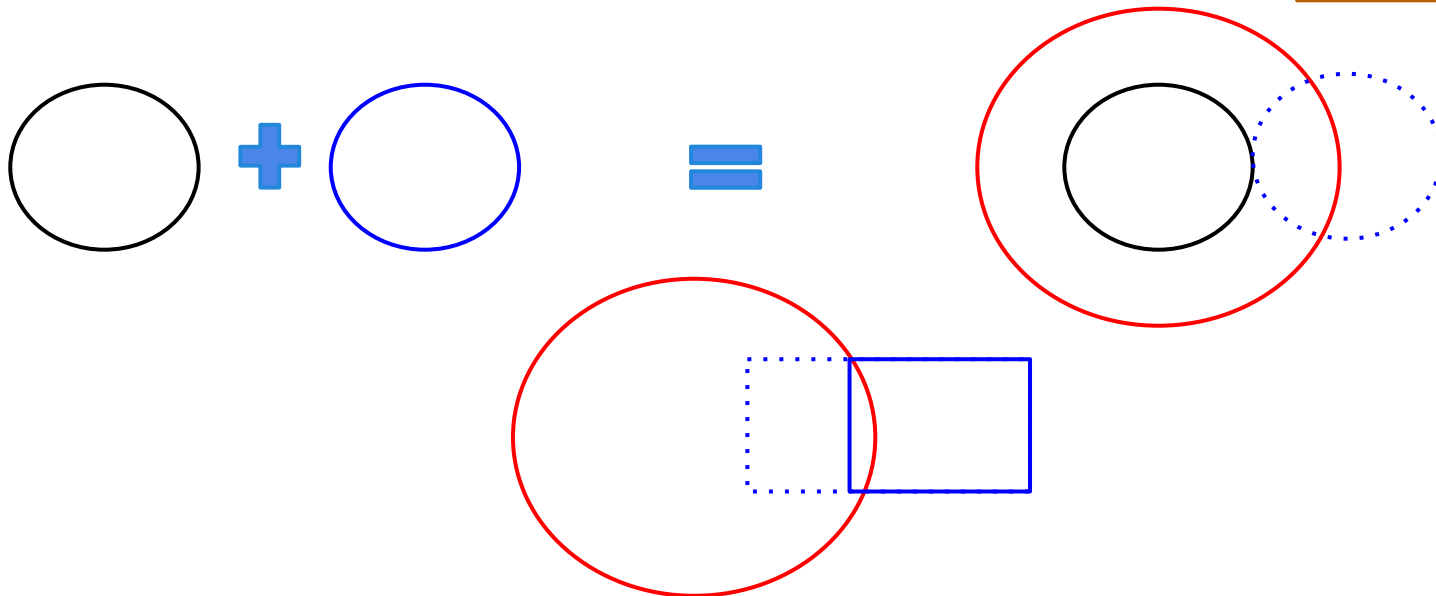
CtcSubpacking

# The Contractor

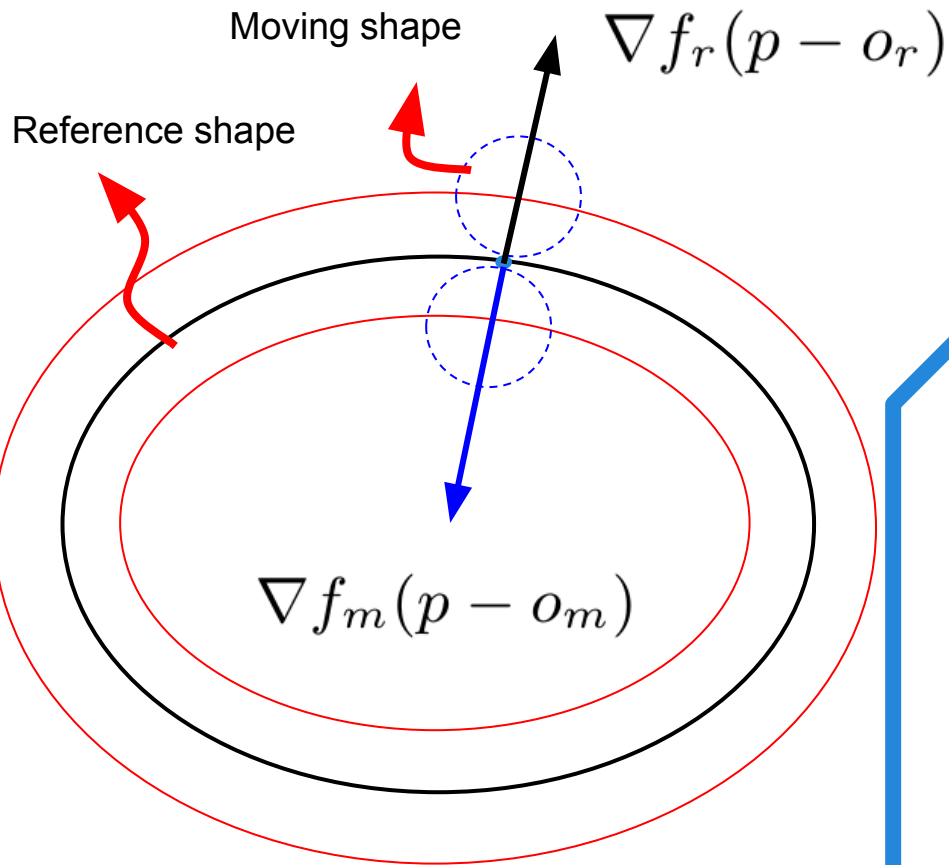
## CtcShadowPropagator

It is important not to lose solutions and perform the contraction in an efficient way

The Shadow



# The Contractor



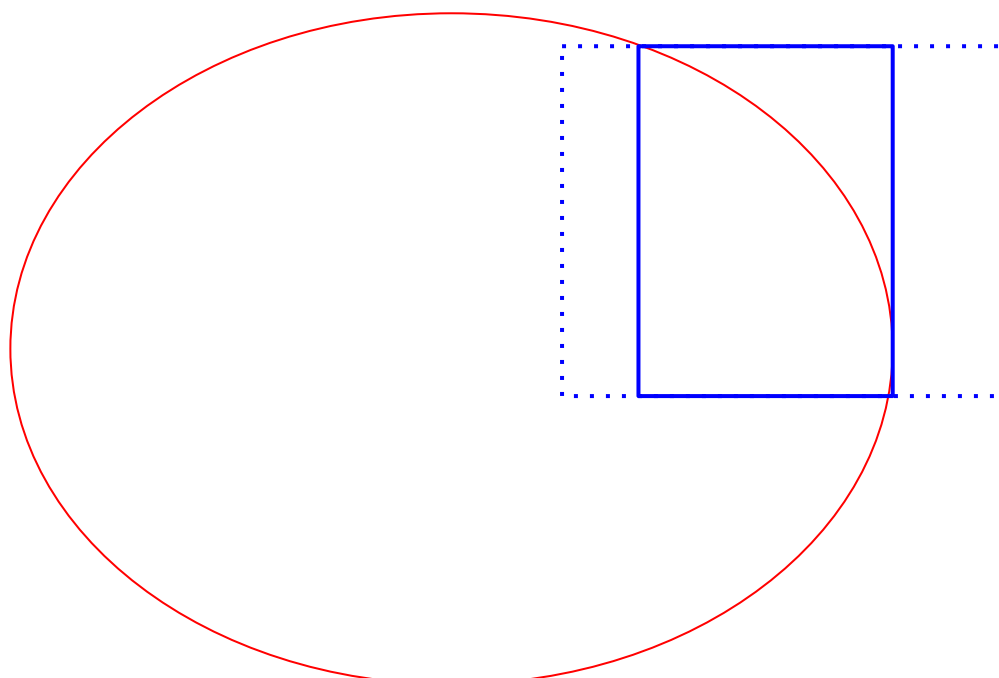
How to calculate the bound of the shadow?

$$f_r(p - o_r) = 0$$

$$f_m(p - o_m) = 0$$

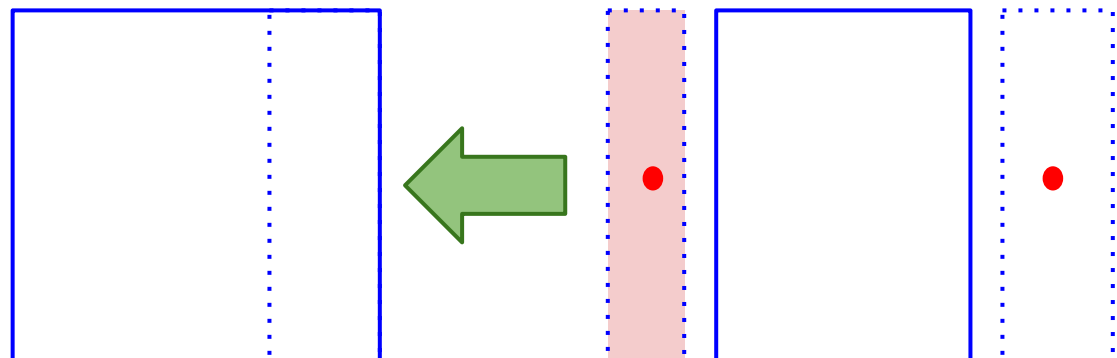
$$\frac{\nabla f_r(p - o_r) \bullet \nabla f_m(p - o_m)}{\|\nabla f_r(p - o_r)\| * \|\nabla f_m(p - o_m)\|} = -1$$

# The Contractor



How to obtain the final contracted box for each pair of objects?

Apply the **lbex** function **diff** with the original box, and check the center of the resulting boxes



# The Contractor

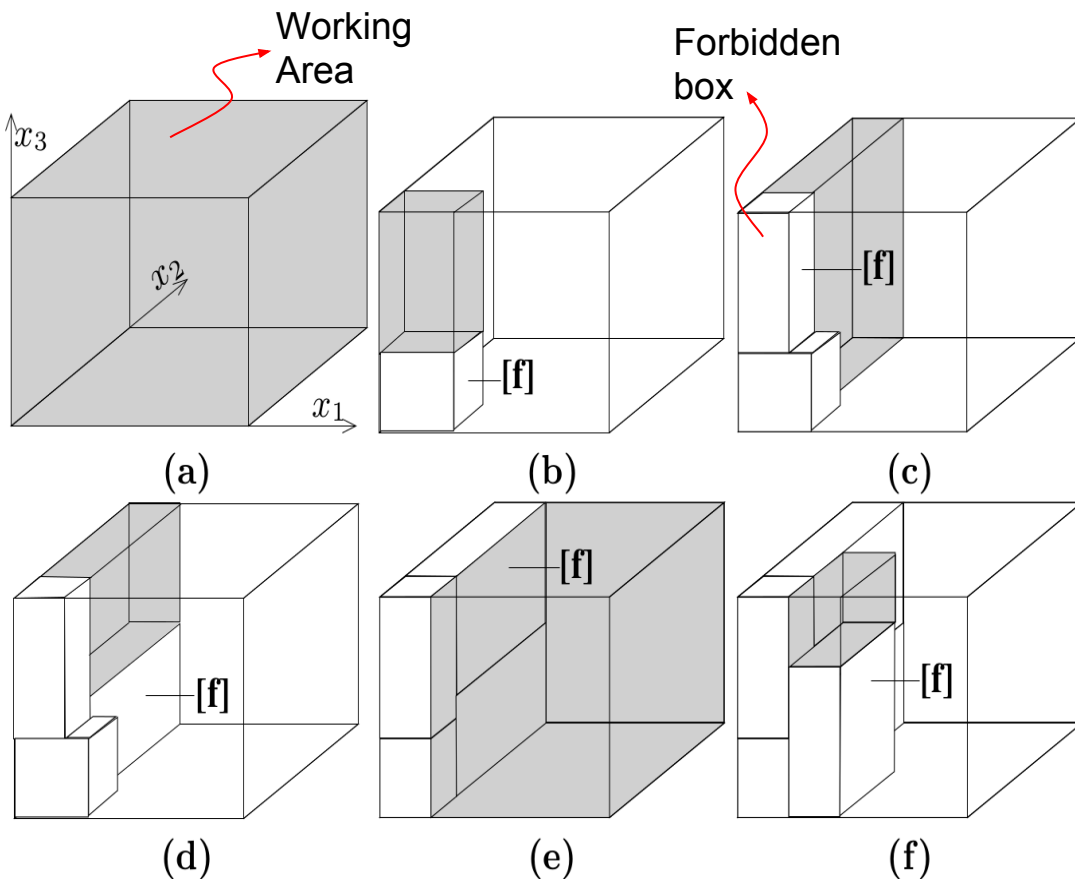
## CtcSubpacking

The origin box of each moving object is contracted with **Sweep**.



# The Contractor

## The Sweep algorithm



The objective is to find all the **forbidden boxes** in a **working area**, such that it will be possible **prune** a section of the box

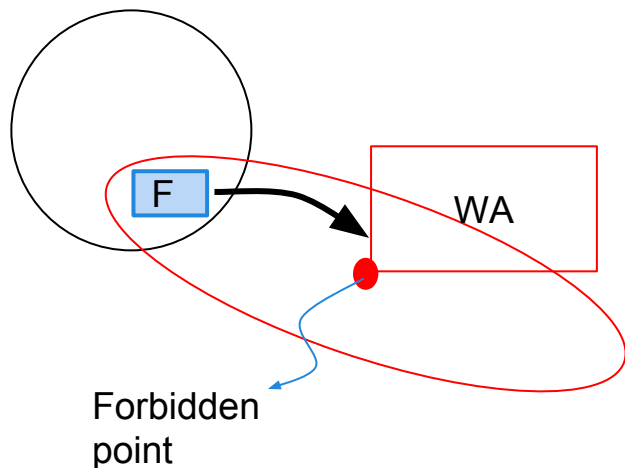
This forbidden boxes are obtained with an **inflator**, that considers one constraint at each time

# The Contractor

## The Inflator

Our main objective it is to find the biggest box in the intersection of the object

And respect the dimensions of the **Working Area (WA)**

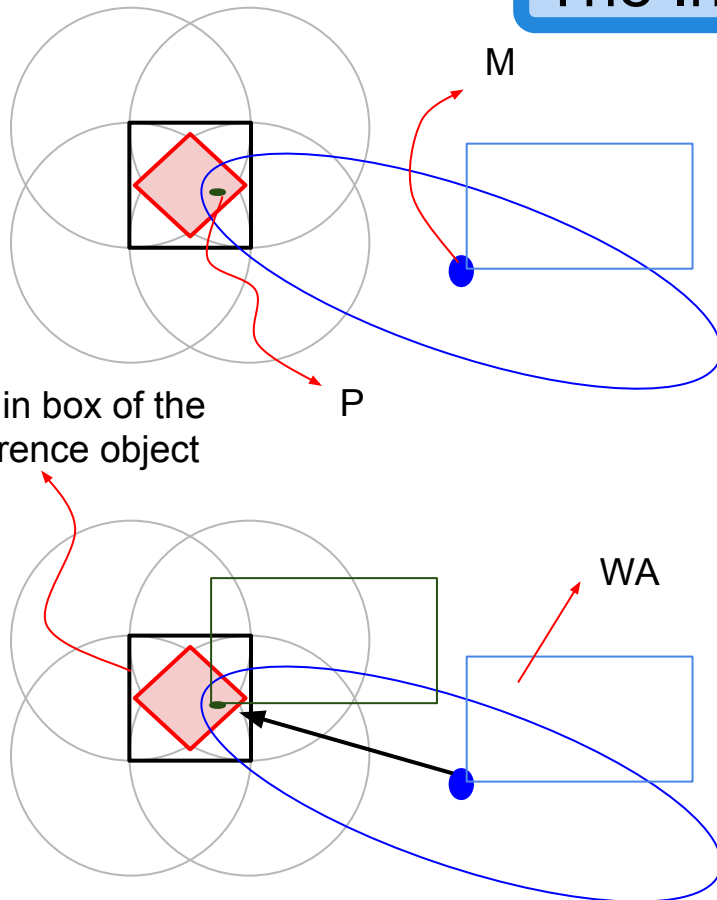


If exists some box **F**, it is possible **translate** it to the **forbidden point**



# The Contractor

## The Inflator

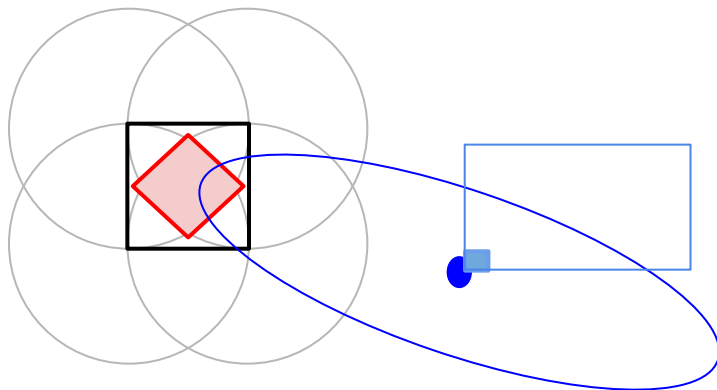
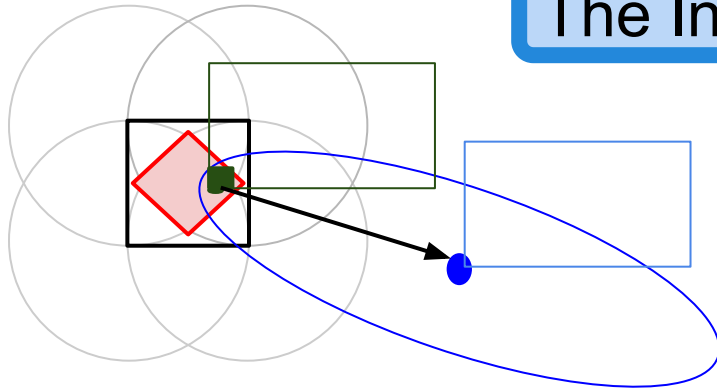


- Center the moving object at  $M$
- Search a point  $P$  that belongs to the intersection of the two objects
- Translate the  $WA$  to  $P$

To be explained further

# The Contractor

## The Inflator



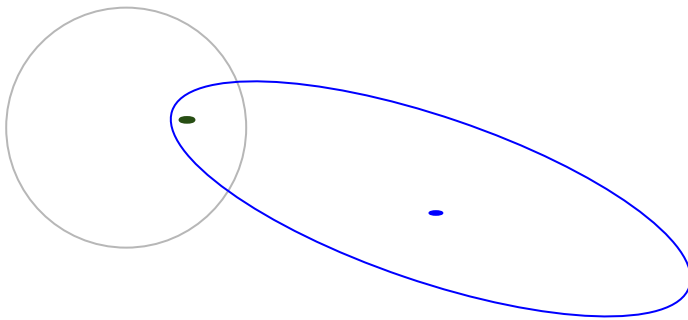
- *Inflate  $P$*  to an inner box of the reference shape

To be explained further

- Translate this box back to the forbidden point

# The Contractor

How to obtain the point **P**?

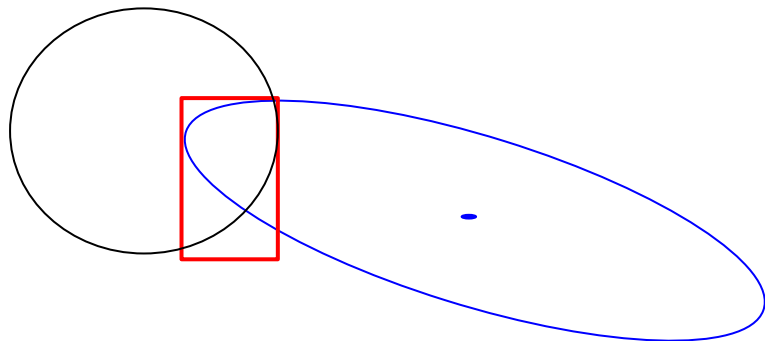
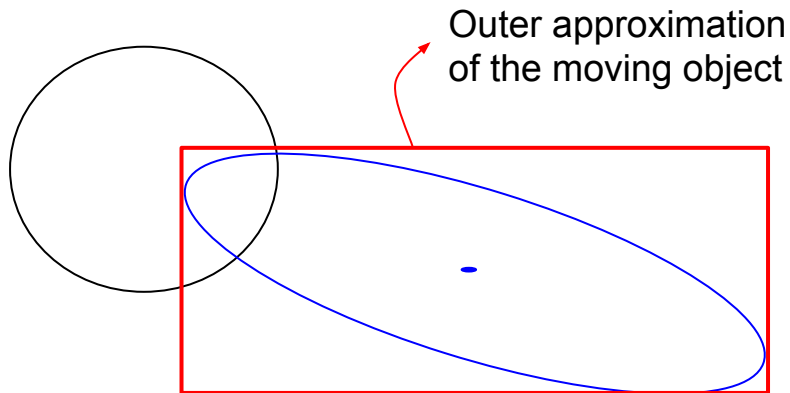


Our objective is find a point that belongs to the moving object and to the **compulsory part** of the reference object

Where always will be possible to find the reference object, given an origin box

# The Contractor

How to obtain the point **P**?



In the beginning, it is an unbounded box

**Contract** some box **C** w.r.t. the moving object

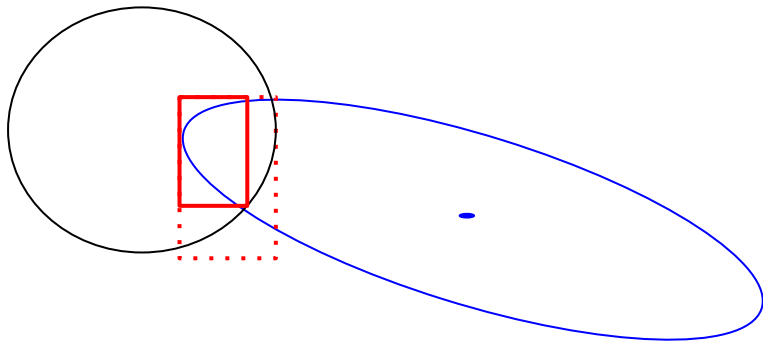
$$o_m + S_m \subseteq C$$

**Contract** the box **C** w.r.t. the reference object, centered in the middle point

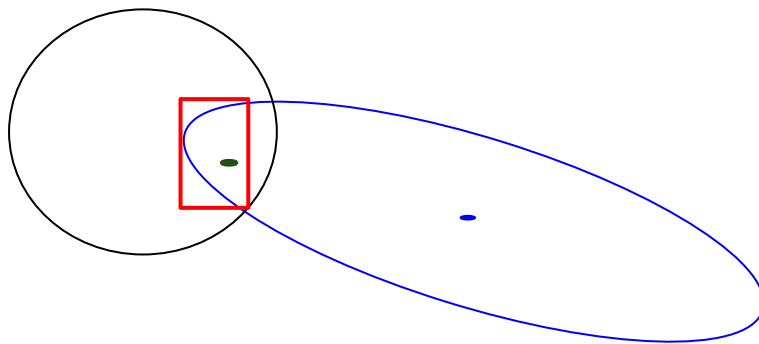
$$o_r + S_r \subseteq C$$

# The Contractor

How to obtain the point **P**?



**Inflate C** w.r.t. the reference object and the middle point of the box **C**



**Pick** randomly a point in the resulting box and check if it belongs to both objects

Using interval evaluation

Else, **Bisect**

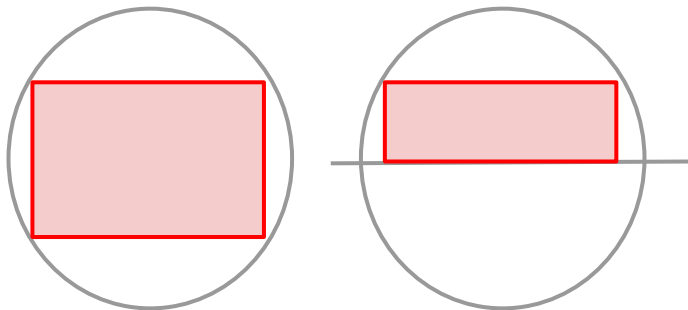
The contracted box

# The Contractor

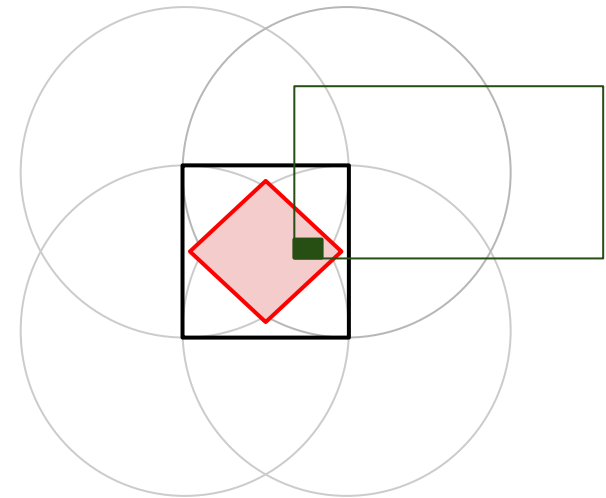
Inflate the point  $\mathbf{P}$

The point  $\mathbf{P}$  is inflated using the box translated and the compulsory part

The inflators of  $\text{Ibex}$  can inflate inside inequalities



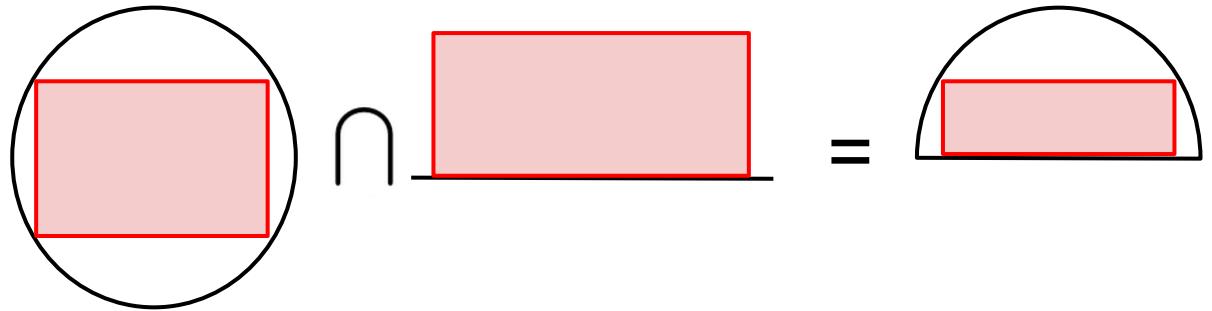
It is possible to inflate inside the same box all the inequalities that are part of the shape



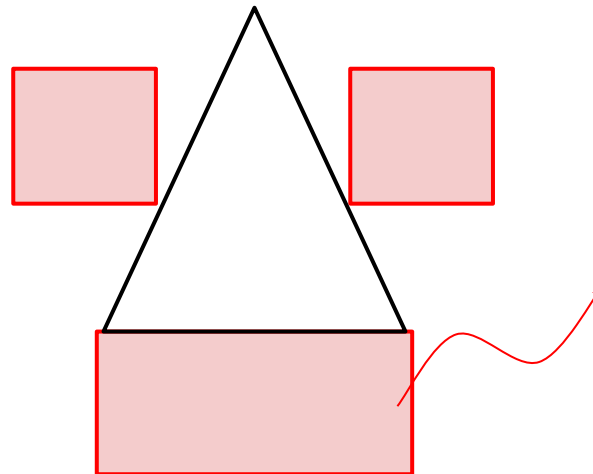
# The Contractor

## Inflate the point **P**

For this we have defined an **intersection** of inner inflators ...



And a **union** of inner inflators

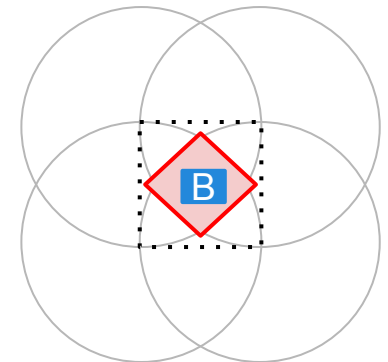


Choose the box with the greatest perimeter

# The Contractor

Hence, it is possible to construct a box **B** inside the **compulsory part** of the reference object, i.e.

$$B \subseteq \left( \bigcap_{p_r \in [O_r]} \{p_r + S_r\} \right)$$





# The Contractor: overlapping shape

- The overlapping shape it is the region where the reference shape will always collide with the moving shape.
- Can be calculated exactly with polytopes, but we are using curved objects.

We calculate something different

$$\textit{Overlapping}_{r,m} = \bigcap_{x \in [\mathbf{X}]} (x \oplus S_r \oplus -S_m)$$

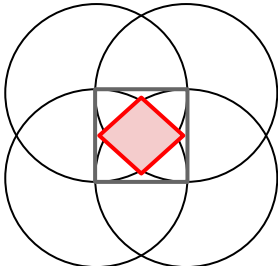
# The Contractor: our forbidden box

The forbidden box that we obtain belongs to this expression

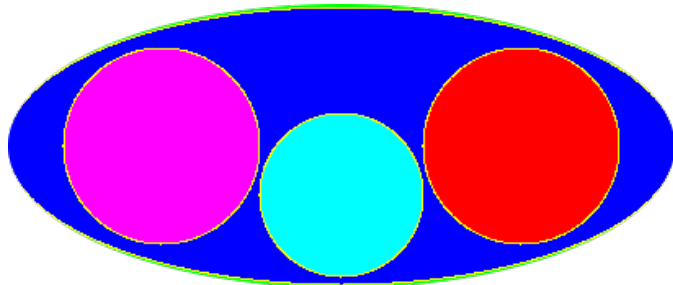
$$B \subseteq \left( \bigcap_{p_r \in [O_r]} \{p_r + [iO]_r\} \right) - (p - o_m)$$

Is an inner box of the reference shape

The intersection point



# Experiments



We consider 2 of the objects as reference and 1 as moving

Reference/Moving	$2^{-M}$	$2^{-M+1}$	...
$2^{-R}$		Diameter of the box of the moving object	
$2^{-R+1}$			
...	Diameter of the box of each reference object		

The contractor that were used to the experiments:

- CtcShadowPropag (S)
- CtcPacking without CtcShadowPropag (P)
- Ctcpacking with CtcShadowPropag (S+P)

Experiments in progress

# Summary

It is possible describe a great number of shapes with disjunction of conjunctions of non linear inequalities

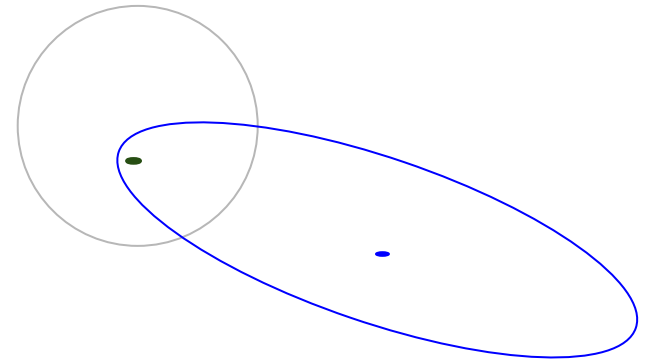
It is possible find the edges of the shadow shape, which give us a first economic contraction

To use an inflator that searchs the boxes that belongs to the intersection of the objects, allows us to obtain an approximation of the optimal forbidden region



# Satisfiability Check

The satisfiability check works similar to the point selector, but only evaluate if the point exists



To find this point, was used a branch & bound algorithm

**Contract** some box with the moving shape

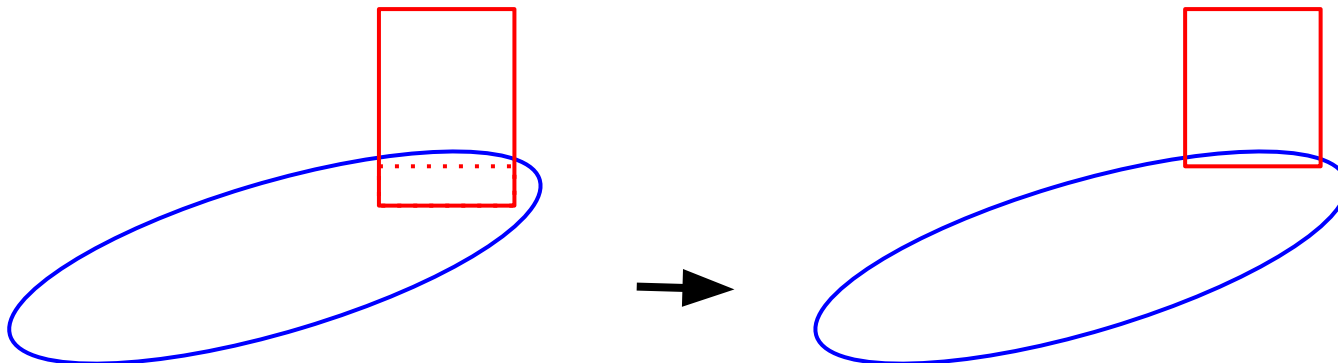
**Contract** the reference shape centered in the some point

**Evaluate** the intersection in a random point of the resulting box

Else, **Bisect**

# Sweep

- The target of the sweep algorithm is to prune all the parts of some domain that violate some constraint
- To perform this, we must “saturate” one of the dimensions with forbidden boxes
- This forbidden boxes must consider some forbidden point and the working area



# Shadow Calculation

$$\frac{\nabla f_r(x - o_r, y - o_r) \bullet \nabla f_m(x - o_m, y - o_m)}{\|\nabla f_r(x - o_r, y - o_r)\| * \|\nabla f_m(x - o_m, y - o_m)\|} = \cos(\theta)$$

$$\frac{\nabla f_r(x - o_r, y - o_r) \bullet \nabla f_m(x - o_m, y - o_m)}{\|\nabla f_r(x - o_r, y - o_r)\| * \|\nabla f_m(x - o_m, y - o_m)\|} = -1$$



$\nabla f_r(x, y)$

$$\nabla f_r(x - o_r, y - o_r) \bullet \nabla f_m(x - o_m, y - o_m) +$$

$$\|\nabla f_r(x - o_r, y - o_r)\| * \|\nabla f_m(x - o_m, y - o_m)\| = 0$$

$\nabla f_m(x - m_x, y - m_y)$